

Hyper-V memory internals. Guest OS memory access

 hvinternals.blogspot.com/2019/09/hyper-v-memory-internals-guest-os-memory-access.html

Software, used in article (operation systems have August 2019 patches):

Windows 10, build 1903 x64

Windows Server 2019

Windows Server 2016

WinDBG Preview

Visual Studio 2019

Process Hacker

PyKd plugin for WinDBG

Testing lab works on Intel-based PC. Therefore, Intel specific Hyper-V terms: hvix64.exe, vmcall instruction, etc will be used in article context.

Terms and definitions:

- WDAG – Windows Defender Application Guard;
- Full VM (virtual machine) – virtual server, which was created in Hyper-V manager. Differs from WDAG container, Windows Sandbox, docker in Hyper-V isolation mode;
- Root OS – operation system, where server part of Hyper-V is installed;
- Guest OS – operation system, which works in Hyper-V emulation context, uses virtual devices, which is presented by Hyper-V infrastructure. It can be Full VM and Hyper-V containers;
- TLFS – Hypervisor Top-Level Functional Specification 5.0;
- GPA (guest physical address) – Guest OS physical memory address;
- SPA (system physical address) – Root OS physical memory address;
- Hypercall – hypervisor service function, which is called by vmcall execution with specifying hypercall number;
- PFN – page frame number.

Source of hvmm.sys driver on github.com:

<https://github.com/gerharto1/LiveCloudKd/tree/master/hvmm>

Python-script for GPAR and MBlock objects parsing

<https://github.com/gerharto1/Hyper-V-Internals/blob/master/ParsePrtnStructure.py>

Intro

Long time ago I didn't write anything in my blogpost. It doesn't mean, that I stopped Hyper-V research. Since Microsoft issued WDAG in Windows 10, build 1803, I started investigate it, but got much problems. First, it was impossible to attach to container, because it doesn't support it. WDAG is isolated environment, and bededit options for debugging can't be configured. More then, every configuration option is reset after rebooting. Sysinternals LiveKD supports Hyper-V attaching, but compatibility was broken in latest OS versions, more then, guest OS memory reading hypercall HvReadGpa, which is used by LiveKd, is not compatible with containers.

It was stalemate, but it turned out, that Matt Suiche (@msuiche), founder of Comae Technologies, shared LiveCloudKd source code for me (many thanks to him!). That program allows attach WinDBG to guest OS, using vid.dll API for reading guest OS memory. But next problem is vid.dll execution blocked by Microsoft: functions from vid.dll can be executed only from vmwp.exe process context, otherwise it will be blocked by vid.sys driver, which compared _EPROCESS object of function's usermode caller process with parent vmwp.exe _EPROCESS. Additionally, some of original LiveCloudKd techniques stopped working in Windows 10. I had to update it too.

Working on adaptation of LiveCloudKd can help me understand Hyper-V guest memory internals better. Soon Matt shared sources on github (<https://github.com/comaeio/LiveCloudKd>).

In 2017, Andrea Allievi made Hyper-V memory management architecture presentation (www.andrea-allievi.com/files/Recon_2017_Montreal_HyperV_public.pptx).

Good work, but details were described quite abstractly, it was hard to match information from presentation to real vid.sys code. I believe it was because at the moment of presentation, Hyper-V symbols information has not yet been published.

Btw, thanks to Andrea to pointing me to some names of vid.sys structures.

Additionally, need say thanks to Microsoft company, which decided to publish symbols for many Hyper-V modules

(<https://docs.microsoft.com/en-us/virtualization/community/team-blog/2018/20180425-hyper-v-symbols-for-debugging>). Without them it was hard to analyze memory-managed vid functions.

First, I planned wrote article about Hyper-V containers, but I made research log above 150 pages (6 from 9 font), but still don't understand whole working scheme. After that I decided to make a list of Hyper-V container components (then, it was extended to all Hyper-V components cheat sheet – no much files were need to add. Containers and Hyper-V has very similar components base).



What categories of hypercalls can be distinguished from this calling statistics? Partition creation, configuring its properties, creating virtual processors and virtual ports (use to send signals, messages), setting interceptions, and various hypercalls for memory management.

See to winhvr.sys!WinHvMapGpaPagesFromMbpArrayScanLargePages function. Rdx contains page number, rsi - size (in pages).

When we start Windows Server 2019 with 1500 Mb of RAM, we got:

1st call rdx=0000000000000000 rsi=000000000005d000

2nd call rdx=0000000000f8000 rsi=0000000000008000

3rd call rdx=0000000000fff800 rsi=0000000000008000

When we start Windows Server 2019 with 2300 Mb of RAM, we got:

1st call: rdx=0000000000000000 rsi=000000000008f000

2nd call: rdx=0000000000f8000 rsi=0000000000008000

3rd call: rdx=0000000000fff800 rsi=000000000000024a

Call stack:

1st call

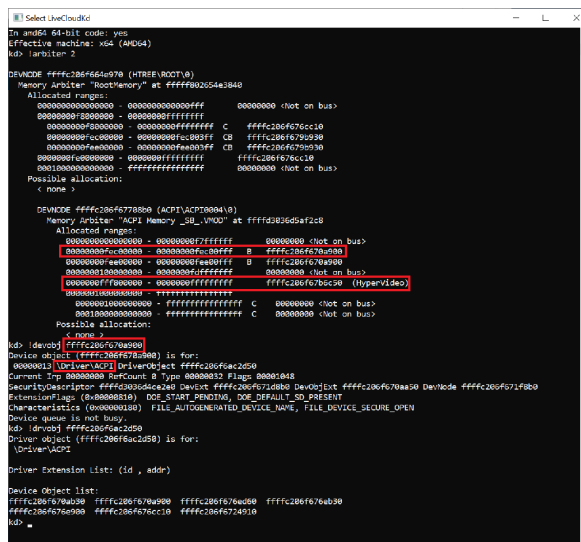
2nd and 3rd calls

```

00 winhvr!WinHvMapGpaPagesFromMbpArrayScanLargePages # Call Site
01 Vid!VsmmHvpMapGpasFromMbpArray 00 winhvr!WinHvMapGpaPagesFromMbpArrayScanLargePages
02 Vid!VsmmHvpMapGpasFromMemoryBlockRange 01 Vid!VsmmHvpMapGpasFromMbpArray
03 Vid!VsmmHvMapGpasFromMemoryBlock 02 Vid!VsmmHvpMapGpasFromMemoryBlockRange
04 Vid!VsmmAdjustGpaSpaceForMemoryBlockRange 03 Vid!VsmmHvMapGpasFromMemoryBlock
05 Vid!VsmmCreateMemoryBlockGpaRange 04 Vid!VsmmAdjustGpaSpaceForMemoryBlockRange
06 Vid!VidloControlPartition 05 Vid!VsmmCreateMemoryBlockGpaRange
07 Vid!VidloControlDispatch 06 Vid!VidloControlPartition
08 Vid!VidloControlPreProcess 07 Vid!VidloControlDispatch
.....WDF Calls..... 08 Vid!VidloControlPreProcess
0d nt!lofCallDriver .....WDF Calls.....
0e nt!lopSynchronousServiceTail 0d nt!lofCallDriver
0f nt!lopXxxControlFile 0e nt!lopSynchronousServiceTail
10 nt!NtDeviceIoControlFile 0f nt!lopXxxControlFile
11 nt!KiSystemServiceCopyEnd 10 nt!NtDeviceIoControlFile
12 ntdll!NtDeviceIoControlFile 11 nt!KiSystemServiceCopyEnd
13 vid_7ffb4de2000!VidCreateMemoryBlockGpaRange 12 ntdll!NtDeviceIoControlFile
14 vmwpl!GpaRangeMbBacked::Initialize 13 vid_7ffb4de2000!VidCreateMemoryBlockGpaRange
15 vmwpl!MemoryManager::CreateGpaRangeInternal 14 vmwpl!MemoryManager::CreateMemoryBlockGpaRange
16 vmwpl!MemoryManager::CreateMemoryBlock 15 vmwpl!VmbComGpaRange::VmbComGpaRange
17 vmwpl!MemoryManager::CreateRamMemoryBlocks 16
18 vmwpl!MemoryManager::CreateRam vmwpl!Vml::VmComMultiInstanceObject<VmbComGpaRange>::CreateInstance
19 vmwpl!VirtualMachine::ConstructGuestRam 17 vmwpl!Vml::CreateComObject<VmbComGpaRange,IMemoryManager
1a vmwpl!WorkerTaskStarting::RunCleanStartSteps 18 vmwpl!VmbComMemoryBlock::CreateGpaRange
1b vmwpl!WorkerTaskStarting::RunTask 19 vmuidevices!VideoSynthDevice::SetupVramGpaRange
1c vmuidevices!VideoSynthDevice::SynthVidOnVramLocation
1d vmuidevices!VideoSynthDevice::OnMessageReceived
1e vmuidevices!VMBusPipeIO::OnReadCompletion
1f vmuidevices!VMBusPipeIO::ProcessCompletionList
1g vmuidevices!VMBusPipeIO::HandleCompletions
1h vmuidevices!VMBusPipeIO::OnCompletion

```

The last memory block is mapped memory of video adapter. A one-page-size block is used for an ACPI devices.



Among other things driver hvmm.sys is needed to remove vmwp.exe protection, that prevent dll injection to that process. That driver works with partition handle with Prtn-signature (VM_PROCESS_CONTEXT), but there is second type, that supporting by vid.sys - EXO-partitions. EXO-partitions can be created using WinHv Platform API Library (<https://docs.microsoft.com/en-us/virtualization/api/hypervisor-platform/hypervisor-platform>), which allows third-party developers to make their virtualization solutions compatible with Hyper-V and run it simultaneously with native Hyper-V VMs. Currently VirtualBox, Qemu, Bochs (f.e. in applepie implementation) have this supporting. VMware, one year after the appearance of these APIs in Windows 1803, finally added support to its VMware Workstation product too. Probably, a new assembly of VMware will be released after the release of Windows 10, build 1909 (19H2).

However, it is still possible to use the vid.dll interface without a driver in Windows Server 2016 and earlier. API execution lock is missing in vid.sys in that OS, and driver hvmm.sys is not needed in that environment. But WDAG and Windows Sandbox containers are presenting in Windows 10 only, where API is locked.

What structures will be needed to work with Guest OS memory? I tried to visualize them in a diagram. In the future, while reading the article, it should become clearly, how they are using.


```

3: kd> dc FFFFC08AE03E000 L30
ffffce08`ae03e000 6e747250 00000000 00000000 00000000 Prtn.....
ffffce08`ae03e010 0480001c 00000000 00000000 00000000 .....
ffffce08`ae03e020 00000001 00000000 00000000 00000000 .....
ffffce08`ae03e030 00000000 00000000 00000000 00000000 .....
ffffce08`ae03e040 00000000 00000000 00000000 00000000 .....
ffffce08`ae03e050 00000000 00000000 004a0048 00000000 .....H.I.....
ffffce08`ae03e060 adb93230 fffffc08 02000014 00000000 02.....
ffffce08`ae03e070 ae03e078 fffffc08 00690057 0032006e x.....W.i.n.2.
ffffce08`ae03e080 00310030 002d0039 00320030 00000000 0.1.9.-.0.2....

```

```

if (objVmPartition->FsContext != NULL)
{
    pPartitionHandle = (PVM_PROCESS_CONTEXT)((PCHAR)objVmPartition->FsContext - 1);
    switch (pPartitionHandle->VmType)
    {
        case VidVmTypeDockerHyperVContainerUserName:
        case VidVmTypeDockerHyperVContainerGUID:
        case VidVmTypeContainer:
            Ret = VidGetContainerMemoryBlock(pPartitionHandle, pBuffer, len, GPA);
            break;
        case VidVmTypeFullWin10VM:
        case VidVmTypeFullWinSrvVMSecure:
        case VidVmTypeFullWinSrvVM:
            Ret = VidGetFullVmMemoryBlock(pPartitionHandle, pBuffer, len, GPA);
            break;
        default:
            break;
    }
}

```

The first 0x278 bytes contain section signature, the name and its identifier. The size of structure is not small (0x3EFO for Windows Server 2019) and it is different for different operating systems. The exact size of partition handle can be found in vid.sys!VidCreatePartition (by the amount of memory allocated for it). We will not need it in driver.

When we get partition handle type (VmType), we can perform one of two procedures for memory blocks reading. There are actually quite a lot of possible VmType values, and moreover, they differ for different versions of operating systems. For example, VmType for Full VM in Windows 10 and Windows Server 2019 have different values. Not all of them have been investigated (especially for operating systems such as Linux, because WinDBG, that launched by LiveCloudKd, doesn't work with them). But finally partitions of virtual machines were divided into two categories: container's partitions and Full VM partitions.

The hvmm.sys!VidGetFullVmMemoryBlock function at the input receives a section descriptor, a buffer in which to write the received data, the size of the buffer in bytes and the GPA of the virtual machine.

BOOLEAN VidGetFullVmMemoryBlock(PVM_PROCESS_CONTEXT pPartitionHandle, PCHAR pBuffer, ULONG len, ULONG64 GPA)
GPA – it is page number, which is calculated: GPA = GpaInfo.StartAddress / PAGE_SIZE;

The start address should be aligned on the page boundary, if the hvmm driver function is called directly (LiveCloudKdSdk prepared usermode buffer for that).

Next, we need to find GPAR object, that describes the requested GPA. Each GPA is included in the memory block, previously allocated by the hypervisor, and this memory block is described by the GPAR object. Fields GpaIndexStart and GpaIndexEnd are located, respectively, at the offsets 0x100 and 0x108 of the GPAR objects. You can understand whether the GPAR object describes the GPA or not, by the value of these fields. For example:

This GPAR object control GPA from 0 to 0x8fbff.

GPAR objects count in Full VM are much smaller than in containers. For example, Generation 2 Full VM has 3-4 GPAR objects, containers have about 780. Then guest OS has more memory, then more blocks it allocates with HvMapGpaPages* hypercalls and, correspondingly, there are greater numbers of GPAR objects. The maximum range of GPAs, described by GPAR object, that I met, was 0x96000 pages.

Let's get back to our driver. We can find GPAR object using hvmm.sys!VidGetGparObjectForGpa function. Partition handle and GPA are passed to the function. How does it work? As described above, each partition handle has a pointer to a GPA block descriptor. This is a structure, which, among other things, contains a pointer to the partition handle itself, a pointer to array with pointers to GPAR objects, and the count of elements in the array of GPAR objects (see the diagram of the relationship of structures above).

```

1: kd> dc fffffd808`20ad2960
ffffd808`20ad2960 72151947 00000000 00000001 00000000 Gpar.....
ffffd808`20ad2970 00650001 00000000 20ad2978 fffffd808 x).....
ffffd808`20ad2980 20ad2978 fffffd808 00000113 00000000 x).....
ffffd808`20ad2990 00000000 00000000 00000000 00000000 .....
ffffd808`20ad29a0 e954b210 ffffff02 20ad2968 fffffd808 ..T.....h).....
ffffd808`20ad29b0 00000000 00000000 00000000 00000000 .....
ffffd808`20ad29c0 00000000 00000000 e9571c00 ffffff02 .....W.....
ffffd808`20ad29d0 20ad2960 fffffd808 00000113 00000000 .....).
1: kd> dc fffffd808`20ad2960+0x100
ffffd808`20ad2a60 00000000 00000000 008fbfff 00000000 .....
ffffd808`20ad2a70 008fbfff 00000000 00000005 00000000 .....
ffffd808`20ad2a80 00000000 00001611 00000000 00000000 .....
ffffd808`20ad2a90 00000000 00000000 00000000 00000000 .....
ffffd808`20ad2aa0 00000000 00000000 00000000 00000000 .....
ffffd808`20ad2ab0 00000000 00000000 00000000 00000000 .....
ffffd808`20ad2ac0 00000000 00000000 00000000 00000000 .....
ffffd808`20ad2ad0 23a49cf0 fffffd808 00000000 00000000 ...#.

```

```

typedef struct _GPAR_BLOCK_HANDLE {
    PVOID PartitionHandle;
    PGPAR_OBJECT GparArray;
    UINT32 Unknown01;
    UINT32 CountInGparArray;
    GPAR_BLOCK_HANDLE *PGPAR_BLOCK_HANDLE;
}

```

```

typedef struct _GPAR_OBJECT {
    CHAR cGparSignature[0x8]; // "GPAR" signature - eq GPA Range
    CHAR Unknown01[0xF8];
    UINT64 GpaIndexStart; //offset +0x100, size 0x8
    UINT64 GpaIndexEnd; //offset +0x108, size 0x8
    UINT64 UnknownParam01;
    UINT64 UnknownParam02;
    UINT32 KernelMemoryBlockGpaRangeFlags; //offset +0x120, size 0x4
    CHAR Unknown02[0x4C];
    PHMEMORY_BLOCK objMBlock; //offset +0x170, size 0x8 //in Windows 10 20H1 up to 0x8 bytes
    ULONG64 SomeGpaOffset; //offset +0x178, size 0x8
    ULONG64 VmmMemGpaOffset; //offset +0x180, size 0x8
} GPAR_OBJECT, *PGPAR_OBJECT;

```

```

pGparBlockHandle = pPartitionHandle->pGparBlockHandle;

if (pGparBlockHandle == 0)
{
    pGparBlockHandle = pPartitionHandle->pGparBlockHandle20H1;
    if (pGparBlockHandle == 0) {
        KDbgPrintString("\tSomething wrong with offset of GparBlockHandle");
        return NULL;
    }
}

```

When we got this information, we can run cycle through the GPAR objects and find 1 GPAR the object, that is responsible for the GPA. Code is quite simple, as you can see. This is a simplified implementation of VsmmLookupMemoryBlockByHandle function of vid.sys driver. Vid.sys driver also has additional procedure for encrypted memory reading - VsmmpSecureReadMemoryBlockPageRangeInternal. It uses AES XTS through BCryptEncrypt\BCryptDecrypt functions from ksecdd.sys driver. I can't find in what cases they are used, because even for Shielded VMs with TPM enabled, memory is not encrypted. Perhaps some special areas are encrypted, but they haven't been found still. But if you try use vid.dll! VidRead\WriteMemoryBlockPageRange functions vid.sys starts analyze second bit in 0x18 byte of Prtn object (test byte ptr [Prtn_obj+18h], 2), and if that bit is not zero crypto-memory functions will be executed. But for standard OS regions they will return fails. It means, for reading Shielded VM memory using vid.dll functions, Prtn object must be patched (2nd bit in 18h byte must be zeroed). Obviously, guest OS directly make reading/writing operations to the already allocated memory area without calling any functions from vid.sys. All exceptions must be caught and handled by the hypervisor. Accordingly, if the root OS encrypts some parts of the memory, then the guest OS will not be able to transparently access them.

Go back to the hvmm code. When we found a suitable GPAR object, we exit from cycle.

There are GPAR objects exist, that don't describe the GPA, but instead of the necessary data, contain a pointer to a certain usermode structure inside the vmwp.exe process. They are tied to the memory allocated for virtual Hyper-V devices. Usually, there is 1 such GPAR object per partition (see content of that memory later in Docker part of that article).

We don't need in that objects during memory reading operations.

What data is contained in the GPAR object and will help to read the data from the guest OS? This is another data type - an MBlock object (MEMORY_BLOCK). It contains guest PFN data and other useful information. A fairly large structure, at the beginning contains the signature "Mb ".

From all the fields, we need only a pointer to the GPA array. Size of the array element is 16 bytes. One 8-byte part contains the GPA (in guest OS), and other 8-byte part contains the SPA information (in root OS).

We can calculate SPA by following formula:

```

Index = pGparBlockHandle->CountInGparArray;
pGparArray = (PVOID64)pGparBlockHandle->GparArray;

if (pGparArray == 0)
{
    KDbgPrintString("\tSomething wrong with offset of Gpar array");
    return NULL;
}

for (LONG i = Index - 1; i >= 0; i--)
{
    uElement = *((PVOID64)pGparArray + i);
    if (uElement != 0)
    {
        objGpar = (PGPAR_OBJECT)uElement;
        KDbgLog("pGparElement->GpaIndexStart", objGpar->GpaIndexStart);
        KDbgLog("pGparElement->GpaIndexEnd", objGpar->GpaIndexEnd);

        if ((GPA >= objGpar->GpaIndexStart) && (GPA <= objGpar->GpaIndexEnd))
        {
            return objGpar;
        }
    }
    else
    {
        KDbgLog("\tGpar Element is NULL, i = ", i);
    }
} // end for

```

```

if (objGpar->GpaIndexStart == objGpar->GpaIndexEnd) {
    KDbgPrintString("MBlock in GPAR object is vmwp.exe descriptor");
    return FALSE;
}

```

```

typedef struct _MEMORY_BLOCK {
    CHAR cMbLockString[0x8]; // "Mb " signature
    PVOID PartitionHandle; // size 0x8
    CHAR Unknown0[0x8];
    ULONG MbHandle;
    CHAR Unknown2[0x1C];
    ULONG64 BitMapSize0; // offset 0x38, size 0x8
    ULONG64 BitMapSize02; // offset 0x40, size 0x8
    CHAR Unknown3[0xA8];
    PULONG64 pGuestGPAArray; //offset 0xF0, size 0x8
} MEMORY_BLOCK, *PMemoryBlock;

```

```

1: kd> dq 0xffffcc8150400000
ffffcc81 50400000 00000000 00078c00 00f20800 00000000
ffffcc81 50400010 00000000 00078c01 00f20800 00000000
ffffcc81 50400020 00000000 00078c02 00f20800 00000000
ffffcc81 50400030 00000000 00078c03 00f20800 00000000
ffffcc81 50400040 00000000 00078c04 00f20800 00000000
ffffcc81 50400050 00000000 00078c05 00f20800 00000000
ffffcc81 50400060 00000000 00078c06 00f20800 00000000
ffffcc81 50400070 00000000 00078c07 00f20800 00000000

```

```

objMBlock = objGpar->objMBlock;
HostSPA = *(PULONG)((PCHAR)objMBlock->pGuestGPAArray + 0x10 * (GPA - objGpar->GpaIndexStart + i));

```

For SPA reading, we need mapped it to root OS virtual address space. Use MDL structure for this:

There is an array of PFN at the end of each MDL structure. A pointer to it can be obtained using MmGetMdlPfnArray macro. When we received the pointer, we had wrote HostSPA index to it. Of course, it is possible to put in MDL more than one PFN at one time. But there is a chance to get to the border of GPAR blocks, therefore memory reading is done page by page. For Full VM, this is not very profitable, since the size of each block is large enough, but speed is still good.

```

pMDL = IoAllocateMdl(VirtualAddress, PAGE_SIZE, FALSE, FALSE, NULL);

```

```

HdIPfnArray = HmGetHdIPfnArray(pMDL);
*HdIPfnArray = HostSPA;

try
{
    SourceAddress = MmMapLockedPagesSpecifyCache(pMDL, KernelMode, MmCached, NULL, FALSE, NormalPagePriority);
    if (!SourceAddress)
    {
        !IfFreeHdI(pMDL);
        return FALSE;
    }

    RtlCopyMemory(pBuffer + 1 * PAGE_SIZE, SourceAddress, PAGE_SIZE);
    MmUnmapLockedPages(SourceAddress, pMDL);
}
except (EXCEPTION_EXECUTE_HANDLER)
{
    KDbgLog(" RtlCopyMemory failed", GetExceptionCode());
}

!IfFreeHdI(pMDL);

```

Next, we get virtual address using the nt!MmMapLockedPagesSpecifyCache function and use it to copy guest OS memory block using nt!RtlCopyMemory. Accordingly, reading is performed in a loop. 1 memory page is copied on 1 iteration. During copying, it is recommended to pause the virtual machine in order to avoid memory modification during reading. In LiveCloudKdSdk, the SdkControlVmState function is implemented for this. It suspends the execution of the virtual machine either by the usual powershell-cmdlets Suspend-VM/Resume-VM, or works with the special register of each virtual processor calling HvWriteVpRegister hypercall and set the HvRegisterExplicitSuspend register to 0 (resume) or 1 (suspend).

Container memory reading

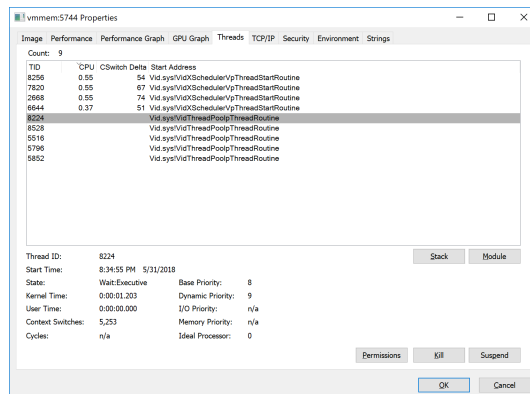
Consider reading the container's memory on Windows Defender Application Guard example (to use it, it's need install same name component in Windows 10. It has been present since the 1803 build). Access to memory of Windows Sandbox and docker container in Hyper-V isolation mode is same.

It made by next function of hvmm.sys driver:

BOOLEAN VidGetContainerMemoryBlock(PVM_PROCESS_CONTEXT pPartitionHandle, PCHAR pBuffer, ULONG len, ULONG64 GPA)

Before executing it, as for Full VM, we must get partition handle first. Then, we will additionally need vmmem process handle. This process is created, when containers work, and works in kernel mode only.

We can see it's threads, when launched container on a 4-processor PC (there are no user mode threads):



The vmmem process descriptor is present in the partition handle. We can find it, using 'scrP' signature (see the hvmm!VidFindVmmemHandle function for details).

We get a pointer to the GPAR object, as same way for reading memory in Full VM. Next we see differences - other fields of the GPAR structure are used to read blocks of memory. VmmMemGpaOffset - the main offset, which allows us convert GPA to SPA for a specific memory block. There is additional offset present (SomeGpaOffset), which can influence to final result, but during my experiments it was always 0.

```

typedef struct _GPAR_OBJECT {
    CHAR cGparSignature[0x8]; // "GPAR" signature - eq GPA Range
    CHAR Unknown01[0xF8];
    UINT64 GpaIndexStart; //offset +0x100, size 0x8
    UINT64 GpaIndexEnd; //offset +0x108, size 0x8
    UINT64 UnknownParam01;
    UINT64 UnknownParam02;
    UINT32 KernelMemoryBlockGpaRangeFlags; //offset +0x120, size 0x4
    CHAR Unknown02[0x4C];
    PHYSICAL_ADDRESS ObjsBlock; //offset +0x170, size 0x8 //in Windows 10 20H1 up to 0x8 bytes
    ULONG64 SomeGpaOffset; //offset +0x178, size 0x8
    ULONG64 VmmMemGpaOffset; //offset +0x180, size 0x8
} GPAR_OBJECT, *PGPAR_OBJECT;

```

Next, we calculate source address, using the following formula and copy data block directly from the address space of vmmem process:


```

for (i = 0; i < uBlocks; i++)
{
    objGpa = VidGetObjectForGpa(pPartitionHandle, Gpa+i);
    if (objGpa == NULL) {
        return FALSE;
    }
    SourceAddress = (Gpa + i - objGpa->GpaIndexStart - objGpa->SomeGpaOffset) * PAGE_SIZE + objGpa->VmmMemGpaOffset;
    if (objGpa->SomeGpaOffset != 0) {
        KDbgLog("  objGpa->SomeGpaOffset", objGpa->SomeGpaOffset);
    }
    try {
        KeReadProcessMemory((PEPROCESS)g_vmmHandle, (PVOID)SourceAddress, pBuffer + i * PAGE_SIZE, PAGE_SIZE);
        Ret = TRUE;
    }
    except (EXCEPTION_EXECUTE_HANDLER)
    {
        KDbgLog("  KeReadProcessMemory exception", GetExceptionCode());
        Ret = FALSE;
    }
}

```

Now we can see key difference between reading container memory from reading Full VM memory: we need copy data from virtual memory of the vmmem process. There is no need for memory mapping using MDL.

Hyper-V memory API

Direct access to memory without corresponding exported Windows functions is interesting, but a more reliable method is to use some of APIs, which is provided by Microsoft. But for reliability you will have to pay the restrictions imposed by Microsoft on these APIs. In particular, for hypercalls they work only with Full VM and for containers they always return FALSE, additionally they read/write no more than 0x10 bytes at one time. The vid.dll function API is generally forbidden to be called from any module other than the vmwp.exe process in latest versions of Windows.

Vid.dll has next functions for reading/writing memory:

- VidTranslateGvaToGpa
- VidReadMemoryBlockPageRange (wrapper on vid.sys!VidReadWriteMemoryBlockPageRange)
- VidWriteMemoryBlockPageRange (wrapper on vid.sys!VidReadWriteMemoryBlockPageRange)

And hypercalls (it must be called from ring 0):

- HvTranslateVirtualAddress
- HvWriteGPA
- HvReadGPA

See it in more detailed.

Reading/writing memory using hypercalls

HvReadGpa using is quite simple, if you don't take, that memory block shouldn't fall on the page boundary. Otherwise, the reading operation will be broken and end of block, that must be read from the second page, will contain zero bytes. Blocking separation is implemented in the usermode part of LiveCloudKdSdk. Driver hvmm calls WinHvReadGPA - HvReadGpa wrapper from winhvr.sys driver. You can call HvReadGpa directly through vmcall, but before you will have to additionally perform operations to prepare hypercall parameters.

```

for (i = 0; i < uBlocks; i++)
{
    Status = WinHvReadGpa(GpaInfo.PartitionId, VpIndex, GpaInfo.StartPage+VID_READ_WRITE_GPA_BUFFER_SIZE, VID_READ_WRITE_GPA_BUFFER_SIZE, ControlFlags, AccessResult, pBuffer + i * VID_READ_WRITE_GPA_BUFFER_SIZE);
    KDbgLog("Status of WinHvReadGpa", Status);
    KDbgLog("AccessResult", AccessResult.ResultCode);
}

```

Boundary checking for writing operation was made in hvmm.sys driver.

```

if (PageBoundaryCheckLowerBorder == PageBoundaryCheckHighBorder)
{
    Status = WinHvWriteGpa(GpaInfo.PartitionId, VpIndex, GpaInfo.StartPage + 1 * VID_READ_WRITE_GPA_BUFFER_SIZE, VID_READ_WRITE_GPA_BUFFER_SIZE, ControlFlags, (PVOID)uPosition, AccessResult);
}
else
{
    PageBoundaryCheckWriteLockSize = (PAGE_SIZE - ((GpaInfo.StartPage + 1 * VID_READ_WRITE_GPA_BUFFER_SIZE) & 0xFFF));
    PageBoundaryCheckWriteLockSize = VID_READ_WRITE_GPA_BUFFER_SIZE - PageBoundaryCheckWriteLockSize;
    Status = WinHvWriteGpa(GpaInfo.PartitionId, VpIndex, GpaInfo.StartPage + 1 * VID_READ_WRITE_GPA_BUFFER_SIZE, PageBoundaryCheckWriteLockSize, ControlFlags, (PVOID)uPosition, AccessResult);
    Status = WinHvWriteGpa(GpaInfo.PartitionId, VpIndex, GpaInfo.StartPage + 1 * VID_READ_WRITE_GPA_BUFFER_SIZE + PageBoundaryCheckWriteLockSize, PageBoundaryCheckWriteLockSize, ControlFlags,
}

```

An additional check is performed before reading virtual address space using winhvr.sys!WinHvTranslateVirtualAddress. The function converts a virtual address into a physical one, using the current context of the CPU (and accordingly, CR3 register).

Possible validation options (LiveCloudKd uses only HV_TRANSLATE_GVA_VALIDATE_READ and HV_TRANSLATE_GVA_VALIDATE_WRITE).

```

#define HV_TRANSLATE_GVA_VALIDATE_READ    (0x0001)
#define HV_TRANSLATE_GVA_VALIDATE_WRITE  (0x0002)
#define HV_TRANSLATE_GVA_VALIDATE_EXECUTE (0x0004)
#define HV_TRANSLATE_GVA_PRIVILEGE_EXEMPT (0x0008)
#define HV_TRANSLATE_GVA_SET_PAGE_TABLE_BITS (0x0010)
#define HV_TRANSLATE_GVA_TLB_FLUSH_INHIBIT (0x0020)
#define HV_TRANSLATE_GVA_CONTROL_MASK    (0x003F)

```

WinDBG in memory dump mode works with physical addresses only (for debugger it is file offsets). Accordingly, it makes all the work for converting virtual address to physical, therefore we don't need to do additional hypercall for checking memory address.

Microsoft Hyper-V Virtualization Infrastructure Driver Library (vid.dll) API
 First, see vid.dll!VidReadMemoryBlockPageRange

```
VIDDLLAPI
BOOL
WINAPI
VidReadMemoryBlockPageRange(
    __in PT_HANDLE Partition,
    __in MB_HANDLE MemoryBlock,
    __in MB_PAGE_INDEX StartMbp,
    __in UINT64 MbpCount,
    __out_bcount(BufferSize)
    PVOID ClientBuffer,
    __in UINT64 BufferSize
);
```

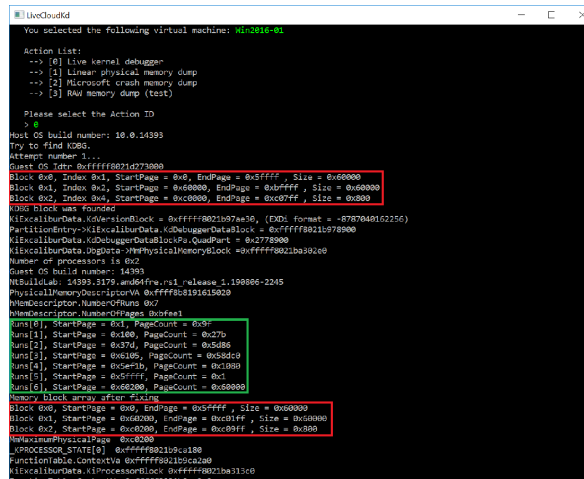
Partition parameter – it is user mode partition handle;
 ClientBuffer – pointer to memory region, where result will be stored;
 BufferSize – yes, buffer size, and nothing more;

Two parameters can cause some questions: MemoryBlock and StartMbp. MemoryBlock is number of the MBlock object from which data will be read. In Windows Server 2008 R2 kernel-mode handle must be pointed as that parameter (yes, the user mode application contained kernel mode descriptor addresses - the original version of LiveCloudKd was built on this logic):

<https://github.com/comaeio/LiveCloudKd/blob/07ac5901ff5cac5258033f1dd95cfc2bd0e06815/hvdd/memoryblock.c#L159>, (buffer contains memory of vmwp.exe)

StartMbp is index, which is equal to physical memory page number. We just need to get the GPA and divide it into PAGE_SIZE (0x1000). The page size in this case is virtual. For example, when ntoskrnl.exe image memory page is usually 2 Mb LARGE_PAGE, but the page numbers will still be 4 Kb granular for that region. Buffer can be specified less, then less data will be written to it. Everything is clear, with one exception - this index is relative to the beginning of MB_HANDLE MemoryBlock. For example, for the first memory block, index will match with physical memory page number. If blocks are placed continuously, index of second block will be equal to page number minus first block size. Index of third block will be equal to page number minus the size of the first block and minus the size of the second block. Everything seems to be clear. The main problem is that physical memory blocks are not continuous. Moreover, these boundaries cannot be easily determined from the user mode. Microsoft didn't provide such APIs even from the time of Windows Server 2008 R2.

```
if ((Buffer[i] >= MemonPagePoolStart) && (Buffer[i] < MemonPagePoolEnd))
{
    for (j = 0; j < BlockIndex; j++)
    {
        if (Blocks[j].MemoryHandle == (MB_HANDLE)Buffer[i])
        {
            Blocks[j].Hits += 1;
            break;
        }
    }
}
```



Matt used a separate function for searching descriptors in memory, but Microsoft closed this opportunity by replacing the descriptors with their indexes in the table, located in kernel mode, and therefore I used vid.dll! VidReadMemoryBlockPageRange function.

```
for (i = 0; i < MAX_INDEX_BLOCK_NUMBER; i++)
{
    Ret = g_VidDll.VidReadMemoryBlockPageRange(PartitionEntry->PartitionHandle, (MB_HANDLE)i, MemoryBlockPageIndex, IULL, Buffer, Size);
    if (Ret == TRUE) {
        //printf("Valid MB_HANDLE: %d\n", (ULONG)i);
        MBlockCount++;
        IndexArray[i] = 1;
    }
}
```

First, we can get the HANDLE numbers by doing a simple search, reading first memory page of each block. If function returns TRUE – it means, that block exists, if FALSE - block doesn't exist. Based on practical experience, I determined the maximum size of the index to be

The first block isn't need for adjustment. Memory is mapping 1 in 1, which allows us to read data from the first block, where ntoskrnl.exe is located, in order to calculate the values of the `_PHYSICAL_MEMORY_DESCRIPTOR` structure later. After calculation, we can perform the offset correction. I described in driver code the case, when one guest block can consist of several blocks, allocated by the hypervisor, but I haven't encountered such case in my stand. The last of the blocks with a size of 0x800 pages is used for video memory, as was explained above. In our case, in a virtual machine, the maximum physical address available for reading is greater than maximum address, specified in `PHYSICAL_MEMORY_DESCRIPTOR`. This block is not specified in `PHYSICAL_MEMORY_DESCRIPTOR`, so we just assume, that it goes sequentially after the last guest OS block. Offset of this block can't be determined without a driver in the host OS. We can assume, that this is memory used by the device, and it can be read, for example, by LiveCloudKd.

After correction, we can read all physical guest OS memory without the driver, excepting pages. Which was paged in pagefile.sys. I complete code description on that point. The remaining details can be found in sources of hvmm driver.

Additional details

I wrote PyKD script `ParsePrtnStructure.py` for better visualization of GPAR objects and Mblock objects (link is given at the beginning of the article). For using it, you have to find partition handle first. To do this, run `hvmmsys driver`, which outputs the value of this descriptor to the debugger and then inserted this value into the script.

Script output for Windows Server 2019 guest OS:

```
0: kd> !py @"F:\ida_files\ParsePrtnStructure.py"
Partition signature: Prtn
Partition name: Win2019-02
Partition id: 2
MBlocks table address: 0xfffffa8024f32bec0L
MBlocks table element count: 2
Gpar block handle address: 0xfffffa8024d426780L
Gpar Element Count: 3
pGparArray address: 0xfffffa8024f0fd550L

GPAR Array content:
-----
Index Sign StartPageNum EndPageNum UmFlag MBlock SomeGPA VMEM GPA
-----
0 Gpar 0x0 0x8f00 0 0xfffffa80248f67d20L 0x0 0x0
1 Gpar 0xfec00 0xfec00 1 0x141f5c3c460L 0x0 0x0
2 Gpar 0xffff800 0xfffffff 0 0xfffffa8024f80f920L 0x0 0x0

MBlock Array content:
-----
Index Sign MCHandle BitmapSize01 BitmapSize02 GPA Array
-----
0 Mb 1 0x8fc00 0x8fc00 0xfffffa8024fc00000L
1 Mb 2 0x800 0x800 0xfffffa8024f8f7000L
0: kd> g
```

Count of GPAR and memory blocks for containers is much more:

```
0: kd> !py @"F:\ida_files\ParsePrtnStructure.py"
Partition signature: Prtn
Partition name: Virtual Machine
Partition id: 3
MBlocks table address: 0xffff958b84cf000L
MBlocks table element count: 958
Gpar block handle address: 0xffff958b794c2c10L
Gpar Element Count: 958
pGparArray address: 0xffff958b70c49000L

GPAR Array content:
-----
Index Sign StartPageNum EndPageNum MemoryBlockGpaRangeFlag MBlock SomeGPA offset VMemGPA offset
-----
0 Gpar 0x0 0x3fff 0 0xffff958b712f4a0L 0x0 0x19435040000L
1 Gpar 0x40000 0x7fff 0 0xffff958b74c2010L 0x0 0x1912c300000L
2 Gpar 0x20000 0x2ffff 0 0xffff958b7234400L 0x0 0x19141300000L
3 Gpar 0xfec00 0xfec00 1 0x23e91c0a0L 0x0 0x0
4 Gpar 0x10000 0x1000c 0 0xffff958b71c0620L 0x0 0x19475040000L
5 Gpar 0x10010 0x1003c 0 0xffff958b71c0840L 0x0 0x19475080000L
6 Gpar 0x10020 0x1003a 0 0xffff958b7346860L 0x0 0x19475200000L
7 Gpar 0x10037 0x10037 0 0xffff958b71c166a0L 0x0 0x19475300000L
8 Gpar 0x1003d 0x10047 0 0xffff958b71d866a0L 0x0 0x19475400000L
9 Gpar 0x10040 0x1004c 0 0xffff958b71c086a0L 0x0 0x19475440000L
10 Gpar 0x1004d 0x1004a 0 0xffff958b720c0620L 0x0 0x19475480000L
11 Gpar 0x10042b 0x100434 0 0xffff958b72027c20L 0x0 0x19475490000L
12 Gpar 0x100435 0x10043a 0 0xffff958b80c3100L 0x0 0x194754a0000L
13 Gpar 0x1004f 0x10071 0 0xffff958b7380760L 0x0 0x194754b0000L
14 Gpar 0x10072 0x1007c 0 0xffff958b71c1700L 0x0 0x19475800000L
15 Gpar 0x1007c 0x1007d 0 0xffff958b615970L 0x0 0x19475850000L
16 Gpar 0x10076 0x1007a 0 0xffff958b74f1c0L 0x0 0x19475870000L
17 Gpar 0x1007e 0x10000 0 0xffff958b74f1c0L 0x0 0x19475870000L
```

In Hyper-V containers all Mblock objects contains zero. Like this:

```
0: kd> dc 0xffff958b7f0d14d0
ffff958b`7f0d14d0 00000000 00000000 00000000 00000000 .....
ffff958b`7f0d14e0 00000000 00000000 00000000 00000000 .....
ffff958b`7f0d14f0 00000000 00000000 00000000 00000000 .....
```

there is additional type of block inside vid.sys driver: reserve bucket block (VSM Reserve Bucket)

But it is not need for reading guest OS memory in standard case. We see that address is pointing to themselves (0x10 alignment).

Docker container with Hyper-V isolation mode

Docker container in Hyper-V isolation mode creates quite a lot of processes (processes for 1 Windows Server 2019 nanoserver 1809 container):

```
MBlock Array content:
-----
Index Sign MCHandle BitmapSize01 BitmapSize02 GPA Array
-----
0 Mb 1 0x40000 0x40000 0xffff958bd010000L
1 Mb 2 0x1d 0xffff958b73fdcd70L
2 Mb 3 0x1f0 0xffff958b7350e00L
3 Mb 4 0x15a 0xffff958b751ef00L
4 Mb 5 0x31 0xffff958b72fd6d20L
5 Mb 6 0x30 0xffff958b73468620L
6 Mb 7 0x15 0xffff958b74afcd0L
7 Mb 8 0xa 0xffff958b6c57300L
8 Mb 9 0xa 0xffff958b80ef90L
9 Mb 10 0xa 0xffff958b80ef820L
10 Mb 11 0x373 0xffff958b7385400L
11 Mb 12 0x16 0xffff958b701350L
12 Mb 13 0x16 0xffff958b870920L
13 Mb 14 0x11 0xffff958b72ab6e0L
14 Mb 15 0x1c 0xffff958b72f760L
15 Mb 16 0x11 0xffff958b72a5170L
16 Mb 17 0x26 0xffff958b71ca560L
17 Mb 18 0x59 0xffff958b6f8010L
18 Mb 19 0x13 0xffff958b74e250L
19 Mb 20 0x1a 0xffff958b72af90L
20 Mb 21 0x2b 0xffff958b83fc5620L
21 Mb 22 0x14 0xffff958b730f880L
22 Mb 23 0x40 0xffff958b756e1300L
```

```
0: kd> dc 0xffff958b70c49000
ffff958b`70c49000 00000000 00000000 .....
ffff958b`4f18300 4f18300 4f18300 4f18300 4f18300 .....
ffff958b`4f18300 4f18300 4f18300 4f18300 4f18300 .....
ffff958b`4f18300 4f18300 4f18300 4f18300 4f18300 .....
ffff958b`4f18300 4f18300 4f18300 4f18300 4f18300 .....
ffff958b`4f18300 4f18300 4f18300 4f18300 4f18300 .....
ffff958b`4f18300 4f18300 4f18300 4f18300 4f18300 .....
```

Process Name	Private Bytes	Working Set	Commitment Bytes	Private Bytes	Working Set	Commitment Bytes	Process Name
vmcompute.exe	3904	0.01	54 B/s	4.28 MB	NT AUTHORITY\SYSTEM	Hyper-V Host Compute Service	
vmmap.exe	4516	0.02	120 B/s	7.8 MB	NT VIRTUAL MACHINE\BAG74342-711D-484A-8087-F19997025858	Virtual Machine Worker Process	
vmtoolsd.exe	6012			5.11 MB	NT VIRTUAL MACHINE\60DC1987-2EA7-4C11-862F-609C8259A7A1	Virtual Machine Worker Process	
vmtoolsd.exe	6224			1 GB	NT VIRTUAL MACHINE\60DC1987-2EA7-4C11-862F-609C8259A7A1	Virtual Machine Worker Process	
vmtoolsd.exe	3308			1 GB	NT VIRTUAL MACHINE\60DC1987-2EA7-4C11-862F-609C8259A7A1	Virtual Machine Worker Process	
vmtoolsd.exe	2324			1 GB	NT VIRTUAL MACHINE\6A95AC6-D266-4887-8E2F-12180F70AE39	Virtual Machine Worker Process	
vmtoolsd.exe	6352			7.34 MB	NT VIRTUAL MACHINE\6A95AC6-D266-4887-8E2F-12180F70AE39	Virtual Machine Worker Process	

We see 2 partition handles (by the count of vmwp.exe processes). The name of 1st of them matches the name of the user in the context of which the process is running.

```

i: kd |py @"F:\lida_files\ParsePrintStructure.py"
Partition signature: Pntn
Partition name: 600C19E7-2A47-4C11-8027-609C82947A1
Partition id: 1
MBlocks table address: 0xfffff8906977a010.
MBlocks table element count: 161
Gpar block handle address: 0xfffff890699e9310.
Gpar Element Count: 1
pGparArray address: 0xfffff89069047800.

GPA Array content:
-----
Index Signature StartPageNum EndPageNum BlockSize MemoryBlockOffsetRangeFlag MBlock SomeDPA offset VmmemDPA offset
-----
0 Gpar 0x0 0x3fff 0x4000 0 0xfffff890694d6500 0x0 0x1280000000 0x1280000000
-----

MBlock Array content:
-----
Index Signature MBlock Address MBlockSize MBlockElementCount MBlockElementOffset GPA Array
-----
0 MBlock 0xfffff890694d6500 1 0x4000 1 0xfffff89069356000
-----

```

However, this partition has irrelevant table of MBlock objects:

```

1: kd> dc 0xfffff8906977a010
fffff8906`977a4010 0000008e 00000000 9948b660 ffff8906 .....H....
fffff8906`977a4020 00000090 00000000 00000002 00000000 .....
fffff8906`977a4030 00000003 00000000 00000004 00000000 .....
fffff8906`977a4040 00000005 00000000 00000006 00000000 .....
fffff8906`977a4050 00000007 00000000 00000008 00000000 .....
fffff8906`977a4060 00000009 00000000 0000000a 00000000 .....

```

Elements count is 0x8e, but the MBlock object itself is only one, and it is empty.

Name of 2nd partition coincides with the identifier, created for container, and contains necessary Nt-kernel data, that can be used to access the memory of the container using WinDBG.

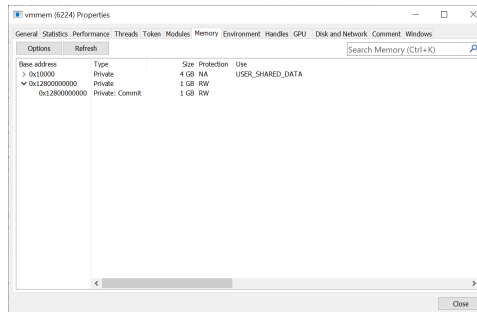
```

i: kd |py @"F:\lida_files\ParsePrintStructure.py"
Partition signature: Pntn
Partition name: c01967089544f675d711e0bc7ffc39f589df9fed48df171729a2a61c85a
Partition id: 4
MBlocks table address: 0xfffff890697f4010.
MBlocks table element count: 159
Gpar block handle address: 0xfffff890699e93a0.
Gpar Element Count: 154
pGparArray address: 0xfffff89069496d10.

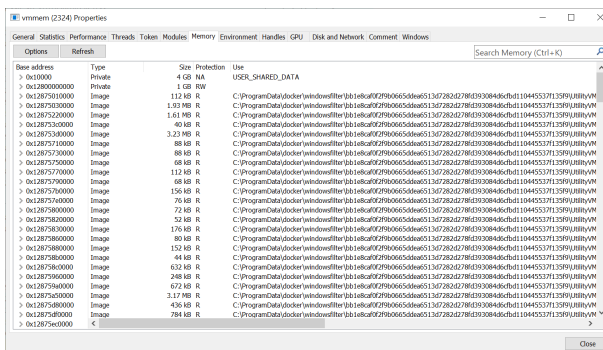
GPA Array content:
-----
Index Signature StartPageNum EndPageNum BlockSize MemoryBlockOffsetRangeFlag MBlock SomeDPA offset VmmemDPA offset
-----
0 Gpar 0x0 0x3fff 0x4000 0 0xfffff890697e4d6500 0x0 0x1280000000 0x1280000000
1 Gpar 0x4000 0x400b 0xc 0 0xfffff890697e4d6500 0x0 0x12875930000 0x12875930000
2 Gpar 0x400c 0x4028 0x1d 0 0xfffff890691ad6500 0x0 0x12875930000 0x12875930000
3 Gpar 0x4029 0x4034 0x1c 0 0xfffff8906975e4d6500 0x0 0x1287520000 0x1287520000
4 Gpar 0x4035 0x403a 0x6 0 0xfffff890691a6500 0x0 0x1287530000 0x1287530000
5 Gpar 0x403f 0x406a 0x33b 0 0xfffff890697e4d6500 0x0 0x1287530000 0x1287530000
-----

```

Base address is the same as the Vmmem GPA Offset parameter, which is used for reading memory block from the context of the vmmem process.



The offset of file mapping region in another vmmem instance are the same as VmmemGPA offset, using by hvmm.sys driver.



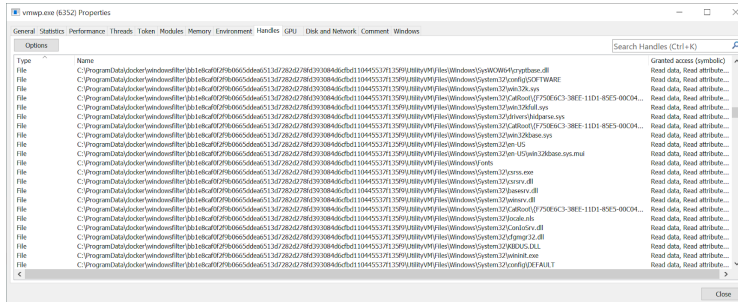
Different vmmem processes load different executables. But in the process, where there are fewer files, the number of active threads is 0.


```

1: kd> dps 00007fff6 1466a418
00007fff6 1466a418 00007fff6 14d99750 vmwp!VIND_HANDLER_CONTEXT: AddReference
00007fff6 1466a420 00007fff6 14d99750 vmwp!VIND_HANDLER_CONTEXT: RemoveReference
00007fff6 1466a428 00007fff6 14d999e0 vmwp!VIND_HANDLER_CONTEXT: GetCallBackBatch
00007fff6 1466a430 00007fff6 14d99a00 vmwp!Vml::VMContext23HostStream: GetBufferOffset
00007fff6 1466a438 00007fff6 14d999f0 vmwp!ProcessorManager::GetVirtualProcessorCount
00007fff6 1466a440 00007fff6 14d99b00 vmwp!ProcessorManager::GetProcessorOvercommitAllowed
00007fff6 1466a448 00007fff6 14d99a00 vmwp!ProcessorManager::GetCpuGroupID
00007fff6 1466a450 00007fff6 14d99a50 vmwp!VIND_HANDLER_CONTEXT:: RTTI Complete Object Locator'
00007fff6 1466a458 00007fff6 14d99a10 vmwp!VIND_HANDLER_CONTEXT:: vector deleting destructor'
00007fff6 1466a460 00007fff6 14d99b70 vmwp!VIND_HANDLER_CONTEXT:: PrepareSelf
00007fff6 1466a468 00007fff6 14d99a00 vmwp!VIND_HANDLER_CONTEXT:: PrepareSelf
00007fff6 1466a470 00007fff6 14d3db00 vmwp!Vml::VMShareableObject: QuiesceSelf
00007fff6 1466a478 00007fff6 14d0b110 vmwp!Vml::VMConnectionPointContainer<ComVirtualMachine>::~VMConnectionPointContainer<ComVirtualMachine>
00007fff6 1466a488 00007fff6 14e9a5d0 vmwp!Vml::VMAutoLock: RTTI Complete Object Locator'
00007fff6 1466a490 00007fff6 14d7ade0 vmwp!Vml::VMAutoLock: vector deleting destructor'

```

Vmwp.exe process of docker container contain descriptor of files, that used inside container:



More information about docker containers internals you can see in video from Microsoft Ignite conference: https://www.youtube.com/watch?time_continue=2291&v=tG8R5SQGPck (OS internals: Technical deep-dive into operating system innovations - BRK3365, starting from 38:11).

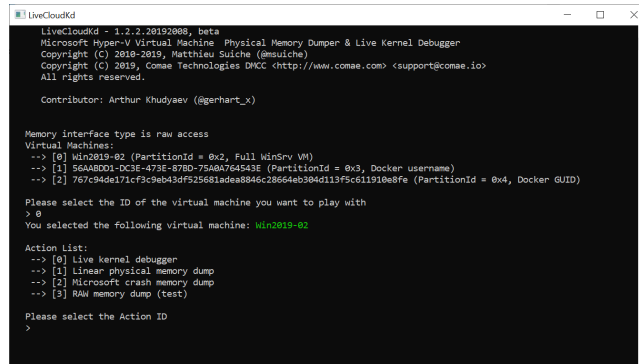
Usage examples

In which programs can we use the ability to read/write memory to the guest OS?

LiveCloudKd (as an alternative to Sysinternals LiveKd in the -hvl option part).

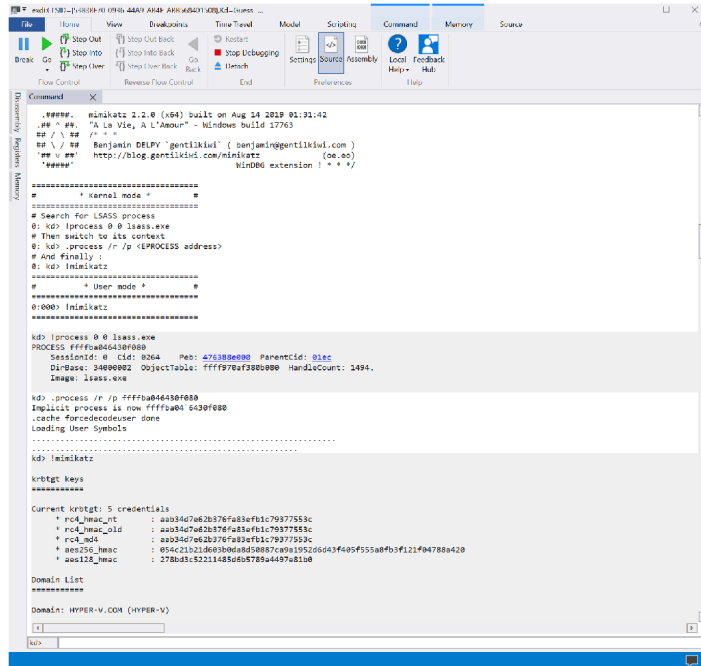
On screenshot, one Full VM with Windows Server 2019 and 1 Docker container in Hyper-V isolation mode are running on Hyper-V host server.

<https://github.com/gerhart01/LiveCloudKd/releases>



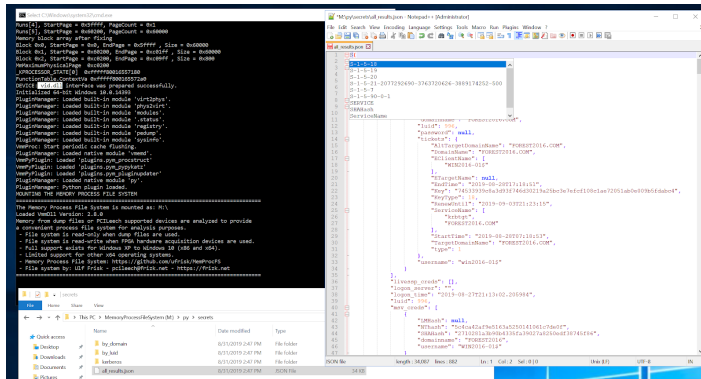
EXDi-plugin for WinDBG - the options are the same, but allows you to use legal functions for WinDBG integration (LiveCloudKd uses hooks of some functions inside WinDBG). It even works with WinDBG Preview, which itself runs in a separate container (UWP application). At the time of writing, EXDi-plugin only works with Windows Server 2019/Windows 10 with the hvmm.sys driver loaded, since it requires a write operation to the guest OS. The screenshot shows the operation of WinDBG Preview in EXDi mode and the mimilib.dll plugin, which is part of the mimikatz utility.

<https://github.com/gerhart01/LiveCloudKd/tree/master/ExdiKdSample>



The plugin for the MemProcFs program (<https://github.com/ufrisk/MemProcFS>), which is integrated with pypykatz (<https://github.com/skelsec/pypykatz>) also allows you to scan the guest OS for hashes (in the screenshot, guest OS - domain controller, based on Windows Server 2016).

<https://github.com/gerharto1/LiveCloudKd/tree/master/LeechCore>



It is clear, that for using this method you need get access the host server with administrator rights. So, first of all, I position the utility as an opportunity to dig inside the OS when the debugger is long configured\too lazy or unable to connect (for example, the Secure Boot option is active).

Conclusion

The article described various ways to accessing memory of Hyper-V guest partitions, created in a variety of cases. I hope that working with Hyper-V memory has become a little more understandable. Hyper-V evolves very quickly and integrates more and more actively into the Windows kernel, while remaining virtually undocumented.

The information may be useful to those who want to understand the internal structure of Hyper-V, and possibly get transparent access to the guest OS memory, as well as make its modification. For LiveCloudKd usage it is necessary to have access to the root OS, where the virtual machines are located, and I don't think that it carries any security risk. However, for Windows Server 2016 such access can be obtained using only the user mode API, which is rather problematic to control. For protection, it is recommended to enable either the Shielded VM option (then, to bypass it, you will need to load the driver), or use Windows Server 2019, where Microsoft blocked the API call from vid.dll for third-party processes and turned on for vmwp.exe the prohibition of injecting libraries, that not signed by Microsoft. However, the latest work on introducing code into third-party processes, demonstrated in August 2019 at Blackhat in Las Vegas (report by Process Injection Techniques - Gotta Catch Them All from Itzik Kotler and Amit Klein from SafeBreach Labs), shows that there are ways to get around these restrictions from user mode (of course, this requires local administrator rights). The only reliable protection against such access to guest OS is Microsoft's Code Integrity in conjunction with the Shielded VM.

