# **Malware-Analysis/SmoothOperator.md at main · dodosec/Malware-Analysis · GitHub**

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main

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## **[Malware-Analysis/](https://github.com/dodo-sec/Malware-Analysis)[SmoothOperator](https://github.com/dodo-sec/Malware-Analysis/tree/main/SmoothOperator)/SmoothOperator.md**

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## **SmoothOperator**

This analysis is focused on the SmoothOperator payloads from Sentinel One. They were obtained via [vx-underground](https://github.com/dodo-sec/Malware-Analysis/blob/main/SmoothOperator/share.vx-underground.org) and comprise two DLLs. The first stage has the hash **bf939c9c261d27ee7bb92325cc588624fca75429**.

## **First stage**

This DLL is a straightforward PE loader, with no obfuscation or encryption present. A good first step is looking for references to VirtualProtect there are two.



First one looks promising, given the ERW flag being passed to it. Checking the function called afterwards (\_\_guard\_dispatch\_icall\_fptr) leads us to an offset, which in turn leads to jmp rax. This is probably a jump to unpacked code or the next stage. Let's circle back to the start of the function where those calls to VirtualProtect are and see what exactly we're marking as executable and then jumping to.



This looks promising. A DLL named d3dcompiler\_47.dll and a call to CreateFileW, followed by memory allocation of the same size as that file. Moving on, we'll see some obvious parsing of a PE file.



Finally, we see a loop that starts looking for the sequence  $0 \times FE$   $0 \times ED$   $0 \times FA$ 0xCE at the Security directory of d3dcompiler\_47.dll and moves forward. If we can find that sequence of bytes in a DLL file, we probably have d3dcompiler\_47.dll - it just so happens that sequence in present in the second DLL from Sentinel One,

**20d554a80d759c50d6537dd7097fed84dd258b3e**. Going forward there are several arithmetic operations followed by the aforementioned VirtualProtect and jmp rax. Instead of worrying about those, just pop the DLL into a debugger, rename

20d554a80d759c50d6537dd7097fed84dd258b3e to d3dcompiler\_47.dll and run until the jmp rax. First stage is done.

## **Second stage**

A quick glance at the debugger following the jmp to rax shows we land at some shellcode at allocated memory.



The dump window also shows the same memory region. One should be careful when dumping it though, since there's plenty of random data preceding the shellcode and d3dcompiler\_47.dll; throwing it in Ida before getting rid of that data will make for an annoying time.

On that note, even though Ida Home supports shellcode analysis, I decided to convert this stage to a PE file. The reason is twofold: first, it means I won't have to import local types manually; second, it means I can keep the dump as is, which is advantageous because we'll be able to follow direct references to the DLL that follows the shellcode. For that end, I do a simple hack with **FASM**:

```
include '..\..\fasmw17330\include\win64ax.inc'
.code
start:
  file 'stage2.bin'
  invoke ExitProcess, 0
.end start
```
The start of the shellcode features basic position independent code (call  $$+5$  followed by pop rcx), which is used to get the address of the start of the DLL read into memory by the first stage into rcx. Another displacement is applied to get a pointer to what appears to be an User-Agent string into r8:

1200 2400 "Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) 3CXDesktopApp/18.11.1197 Chrome/102.0.5005.167 Electron/19.1.9 Safari/537.36"



The next call is to a function that will be tasked with mapping d3dcompiler\_47.dll. Although it's already in memory, it has not been mapped as an executable needs to be before it's able to run. Here's the start of it, after renaming and adjusting types for the arguments to match what was placed in the registers preceding the call.



In another common practice with shellcode, API hashes are present. [HashDB](https://hashdb.openanalysis.net/) identifies the algorithm employed here as one used by Metasploit. If one decides to look into the mw\_import\_by\_hash function, it's important to remember that this code deals with PEB64 and TEB64, structs that I couldn't find in Ida. I recommend [this resource from BITE\\*](http://bytepointer.com/resources/tebpeb64.htm) to create your own struct for both. Doing this will solve you a couple hours of confused cursing at the 32 bit structures.

Next up the actual mapping of the DLL into memory takes place. This is made evident by several snippets of code that parse PE Section Headers relevant to the mapping process. The one below checks to see if the PE

being processed is indeed for a 64-bit architecture; other lines deal with the PE sections and the entry point address:



A bit further down, memory is allocated to match the size of the DLL (according to the value in



IMAGE\_NT\_HEADERS64.OptionalHeader.SizeOfImage):

A very interesting sequence follows. It's responsible for resolving all imports of the third stage DLL by using LoadLibraryA and

GetProcAddress. Taking note of which fields of the PE are being parsed and watching a few loops of it running will help you grasp how an import table is built when an executable is mapped.



A lot more code follows this, mapping sections and using VirtualProtect to assign the correct protections to each one. We're almost done now!

There's then a call rbx instruction that leads to a rabbit hole of shellcode functions. Unfortunately what follows next is something no one likes to read in an analysis like this, but *I have no idea what those do*. My educated guess is some combination of anti-emulation/anti-sandbox, since there are multiple uses of the cpuid instruction in there and a test following that call will skip the jump to the next stage and instead just return. If anyone is curious, feel free to give it a look.



After the return from the mysterious shellcode rabbit hole, we have only a few steps left. The code ensures it has mapped the DLL correctly by checking the size of its Data Directory and the exported functions (there is only one, DllGetClassObject); it then maps the address of said name to r8. Then the name of the export itself is checked by a simple ROR 13 ADD hash function, another callback to metasploit:



Finally, the arguments (remember those from ages ago??) are put back into the relevant registers and there is a jump to r8, which now holds the address of exported function of the third stage DLL. Its command line arguments are the User-Agent string from earlier and the constant 0xAA (thanks to the [Sentinel One Report](https://www.sentinelone.com/blog/smoothoperator-ongoing-campaign-trojanizes-3cx-software-in-software-supply-chain-attack/) for pointing out that this constant is the size of the User-Agent string).

## **Important time-saving tip:**

It's only as I wrap up this write-up that I realized there is no decryption of the third stage DLL done by the shellcode, only mapping and *maybe* some anti-emulation shenanigans. As such, one can really speed up their analysis by extracting the full stage 2 payload and getting rid of everything before the MZ header of the third stage DLL.