

SmokeLoader Triage

🔗 research.openanalysis.net/smoke/smokeloader/loader/config/yara/triage/2022/08/25/smokeloader.html

OALABS Research

August 25, 2022

Overview

Samples

- Packed parent
[cef4f5f561b5c481c67e0a9a3dd751d18d696b61c7a5dab5ebb29535093741b4](#)
- Unpacked SmokeLoader
[041a05dd902a55029449bf412cedbe59a593f8d4e67d4ae37cf7a928c92f22ca](#)

SmokeLoader Background

This SmokeLoader sample is on [MalwareBazaar](#) and through sandbox runs we know that it was used to download Vidar. From [JoeSandbox public report](#) we know we should find the following config in this loader

```
{
  "c2 list": [
    "http://piratia.su/tmp/",
    "http://piratia-life.ru/tmp/",
    "http://diwebseite.at/tmp/",
    "http://faktync.com/tmp/",
    "http://mupsin.ru/tmp/",
    "http://aingular.com/tmp/",
    "http://mordo.ru/tmp/"
  ]
}
```

References

- [Deep Analysis of SmokeLoader](#)
- [Historical Changes and Trends \(Marcos Alvares\)](#)

Stage 2

Opaque predicate deobfuscation

From this [blog](#) we have a simple jmp fix script.

```

import idc

ea = 0
while True:
    ea = min(idc.find_binary(ea, idc.SEARCH_NEXT | idc.SEARCH_DOWN, "74 ? 75 ?"), #
    JZ / JNZ
        idc.find_binary(ea, idc.SEARCH_NEXT | idc.SEARCH_DOWN, "75 ? 74 ?")) #
    JNZ / JZ
    if ea == idc.BADADDR:
        break
    idc.patch_byte(ea, 0xEB) # JMP
    idc.patch_byte(ea+2, 0x90) # NOP
    idc.patch_byte(ea+3, 0x90) # NOP
..

```

Once we fix the jmps we need to nop out the junk code between the code to allow IDA to convert this into a function

```

```python
import idaapi

start = 0x00402DDD
end = 0x00402EBF
ptr = start
while ptr <= end:
 next_ptr = next_head(ptr)
 junk_bytes = next_ptr - ptr
 if ida_bytes.get_bytes(ptr, 1) == b'\xeb':
 idaapi.patch_bytes(ptr, junk_bytes * b'\x90')
 ptr = next_ptr

```

Or, we could use this excellent script from [@anthonyprintup](#)

```

import ida_ua
import ida_name
import ida_bytes

def decode_instruction(ea: int) -> ida_ua.insn_t:
 instruction: ida_ua.insn_t = ida_ua.insn_t()
 instruction_length = ida_ua.decode_insn(instruction, ea)
 if not instruction_length:
 return None
 return instruction

def main():
 begin: int = ida_name.get_name_ea(idaapi.BADADDR, "start")
 end: int = begin + 0xE2

 instructions: dict[int, ida_ua.insn_t] = {}

 # Undefine the current code
 ida_bytes.del_items(begin, 0, end)

 # Follow the control flow and create instructions
 instruction_ea: int = begin
 while instruction_ea <= end:
 if instruction_ea not in instructions.keys():
 instruction: ida_ua.insn_t = ida_ua.insn_t()
 instruction_length: int = ida_ua.create_insn(instruction_ea, instruction)
 else:
 instruction: ida_ua.insn_t = decode_instruction(instruction_ea)
 instruction_length: int = instruction.size
 if not instruction_length:
 print(f"Failed to create an instruction at address {instruction_ea:#x}")
 return

 # Append the current instruction address to the list
 instructions[instruction.ip] = instruction

 # Handle unconditional jumps
 current_instruction_mnemonic: str = instruction.get_canon_mnem()
 next_instruction: ida_ua.insn_t | None = decode_instruction(instruction_ea +
instruction.size)
 if next_instruction is not None:
 next_instruction_mnemonic: str = next_instruction.get_canon_mnem()
 if (current_instruction_mnemonic == "jnz" and next_instruction_mnemonic
== "jz") or \
 (current_instruction_mnemonic == "jz" and
next_instruction_mnemonic == "jnz"):
 # Unconditional jump detected
 assert instruction.ops[0].type == ida_ua.o_near
 instruction_ea = instruction.ops[0].addr

```

```

 ida_ua.create_insn(next_instruction.ip)
 instructions[next_instruction.ip] = next_instruction
 continue

 if current_instruction_mnemonic == "jmp":
 assert instruction.ops[0].type == ida_ua.o_near
 instruction_ea = instruction.ops[0].addr
 else:
 instruction_ea += instruction.size

NOP the remaining instructions
for ea in range(begin, end):
 skip: bool = False
 for _, instruction in instructions.items():
 if ea in range(instruction.ip, instruction.ip + instruction.size):
 skip = True
 break
 if skip:
 continue

Patch the address
ida_bytes.patch_bytes(ea, b"\x90")

if __name__ == "__main__":
 main()

```

After this we can see that the next function address is built using some stack/ret manipulation.

## Function Decryption

---

Some functions are encrypted. We can find the first one by following the obfuscated control flow until the first `call`. This call calls into a function which then calls the decryption function. The decryption function takes a size and a offset to the function that needs to be decrypted. The size is placed in the `ecx` register, and the function offset follows the call.

```
mov ecx, 0E7h ; 'ç'
```

The decryption itself is a single byte xor but the decryption key is moved into the `edx` register as a full DWORD (we only used the LSB).

```
mov edx, 76186250h
```

From this [blog](#) we have a simple deobfuscation script updated for our sample. This script didn't perform well for some reason so we ended up manually decrypting the functions!

```

import idc
import idutils

def xor_chunk(offset, n):
 ea = 0x400000 + offset
 for i in range(n):
 byte = ord(idc.get_bytes(ea+i, 1))
 byte ^= 0x50
 idc.patch_byte(ea+i, byte)

def decrypt(xref):
 call_xref = list(idutils.CodeRefsTo(xref, 0))[0]
 while True:
 if idc.print_insn_mnem(call_xref) == 'push' and
idc.get_operand_type(call_xref, 0) == idaapi.o_imm:
 n = idc.get_operand_value(call_xref, 0)
 break
 if idc.print_insn_mnem(call_xref) == 'mov' and
idc.get_operand_type(call_xref, 1) == idaapi.o_imm:
 n = idc.get_operand_value(call_xref, 1)
 break
 call_xref = prev_head(call_xref)
 n = idc.get_operand_value(call_xref, 0)
 offset = (xref + 5) - 0x400000
 xor_chunk(offset, n)
 idc.create_insn(offset+0x400000)
 ida_funcs.add_func(offset+0x400000)

xor_chunk_addr = 0x00401118 # address of the xoring function
decrypt_xref_list = idutils.CodeRefsTo(xor_chunk_addr, 0)

for xref in decrypt_xref_list:
 decrypt(xref)

```

## API Hashing

---

According to this [blog](#) we are expecting to see some API hashing using the **djb2** algorithm. We can try to find this function by searching for the constant 0x1505.

Though the djb2 algorithm is used for the API hashing the API hashes have some further obfuscation that prevents them from being resolved directly .... **TODO**

```
import requests

hash = 812437175

def hash_djb2(s):
 hash = 5381
 for x in s:
 hash = ((hash << 5) + hash) + x
 return hash & 0xFFFFFFFF

TEST_1 = 1555243728
test = b'ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz0123456789'

print(hash_djb2(test))

api_name = b'CopyFileW\x00'
api_hash = 0x306cceb7

print(hex(hash_djb2(api_name)))

205455184
0x306cceb7
```

## Extract Stage 3

---

There is a 32-bit and a 64-bit version of stage 3 stored consecutively in the binary. In older versions of the loader these stages were simply compressed with LZN1 but in this version there appears to be another step that is missing... **TODO**

```

data_string =
'5058187647a21889fa5d103a51605c8eb0621a575b63149f64673087b5753166bb323066f2465570508c1

data_string =
'5036187647a21889fa5d1e12d6615c8ea0623a565b60149f1a8e1e471479b89d409f9877506c1a7a7446c

def unhex(hex_string):
 import binascii
 if type(hex_string) == str:
 return binascii.unhexlify(hex_string.encode('utf-8'))
 else:
 return binascii.unhexlify(hex_string)

data = unhex(data_string)

import struct
def decrypt_dw(data, dw):
 out = b''
 for i in range(0, (len(data)//4)*4, 4):
 tmp = struct.unpack('<I', data[i:i+4])[0]
 out += struct.pack('<I', tmp ^ dw)
 return out

out = decrypt_dw(data, 0x76186250)

out2 = []
for c in out:
 out2.append(c^ 0x50)

import malduck
ptxt_data = malduck.lznt1(bytes(out2))
print(ptxt_data)
#open('/tmp/out.bin', 'wb').write(ptxt_data)

b''

```

## Stage 3

---

**@doomedraven**

i need to leave in 3rd stage there is a cnc table, just search in 3rd stage for immediate value of 1000 and you will be pretty close to the table