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Reverse Engineering the Mustang Panda PlugX RAT – Extracting the Config

Hello everyone! This is a continuation on the series of blog posts focused on reverse engineering a new-ish variant of PlugX malware gaining traction around the Asia Pacific region.

On my previous [post](#), we reverse engineered the loader to determine how it decrypts, load, and execute the actual RAT component of PlugX. For this post, we will continue from where we left off and focus on one thing – extracting the malware configuration.

Introduction

Extracting malware configuration is one of the sub-tasks we can focus on when reverse engineering malware. By accomplishing this, we can obtain information used by the malware during execution such as hardcoded C2 addresses, bot IDs, mutexes, file paths, and registry keys among others.

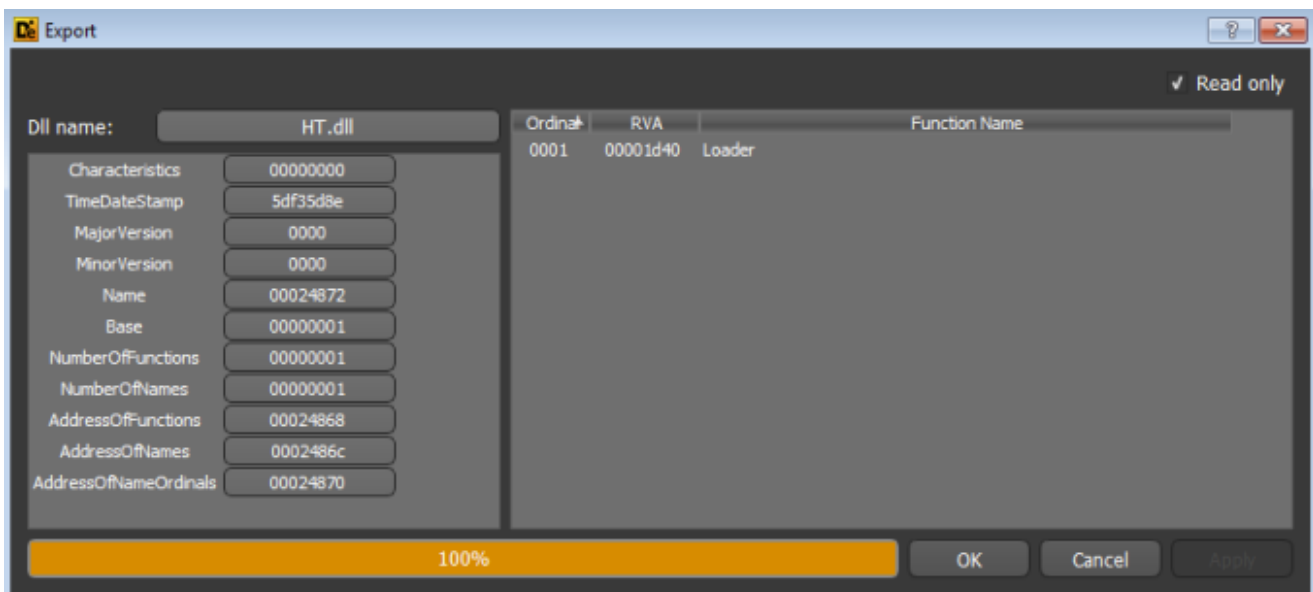
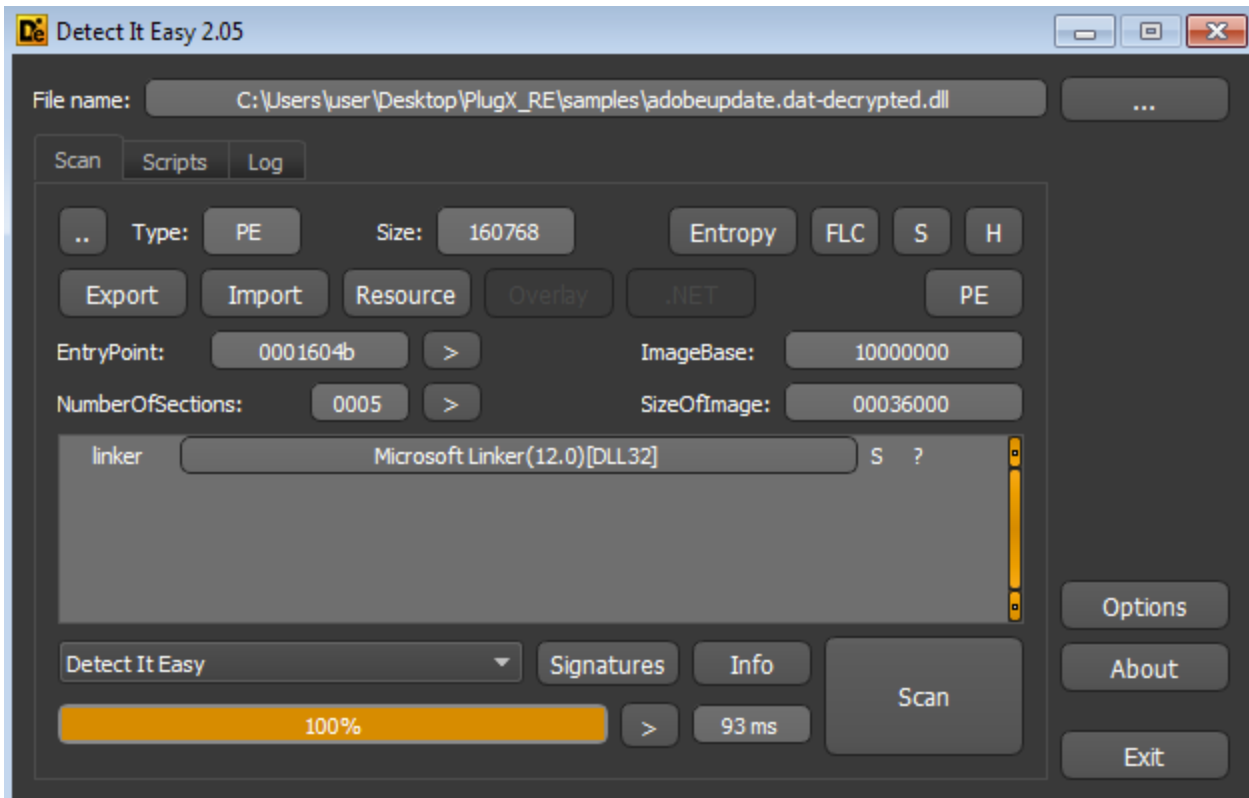
More often than not, a malware configuration is stored within the malware binary itself as an encoded or encrypted blob to prevent easy detection and analysis. We have to do some reverse engineering in order to figure out where it is located, how it is encoded or encrypted, and how to reverse the algorithm to see the plain text configuration. After successfully extracting the configuration, we can create an automation scrip to facilitate extraction to a huge number of samples. This is extremely useful for a number of use cases such as tracking C2 addresses for a specific malware family, blacklisting C2 addresses on network security devices, hunting for indicators-of-compromise within an environment, assisting incident responders during an investigation, and so on.

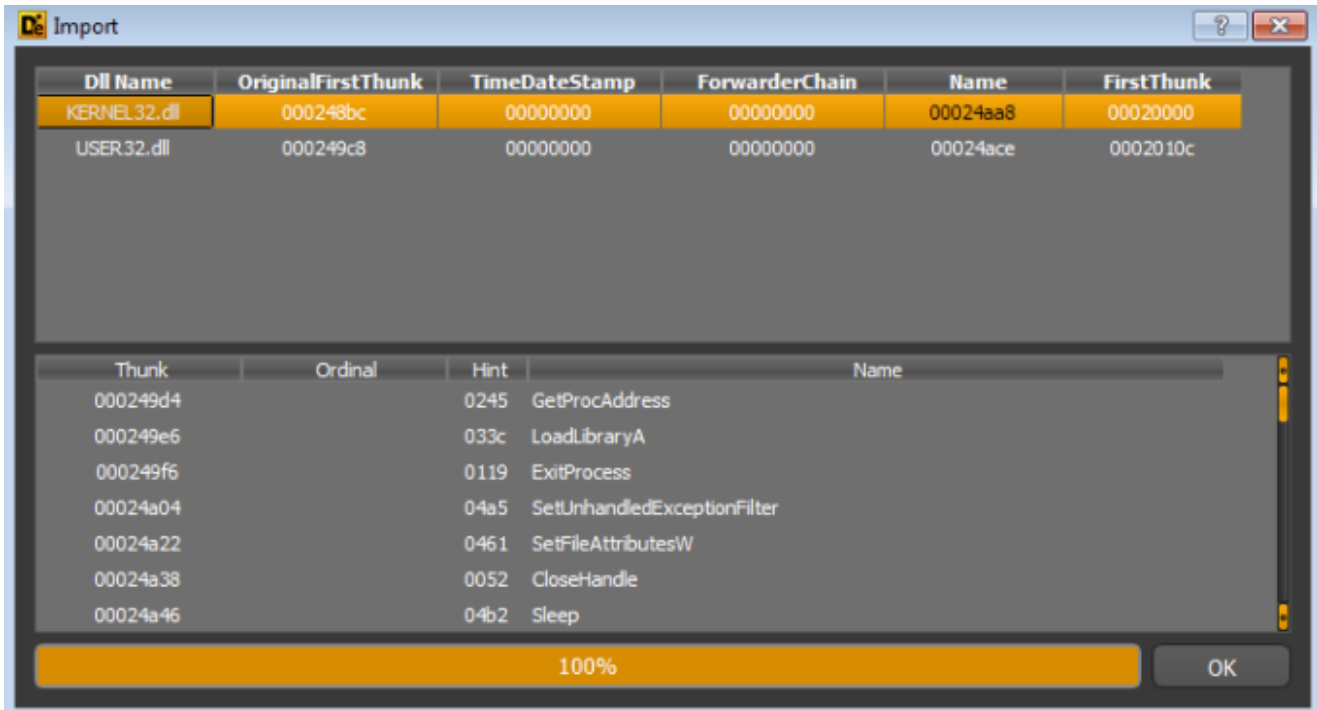
Let's continue with the analysis.

Extracting the configuration

On my previous [post](#), we were able to figure out that the payload is decrypted by the loader using XOR with multi-byte key. We also created an automation script to decrypt the payload and save it to disk. We'll use this script to save a copy of the decrypted payload to disk and start from there.

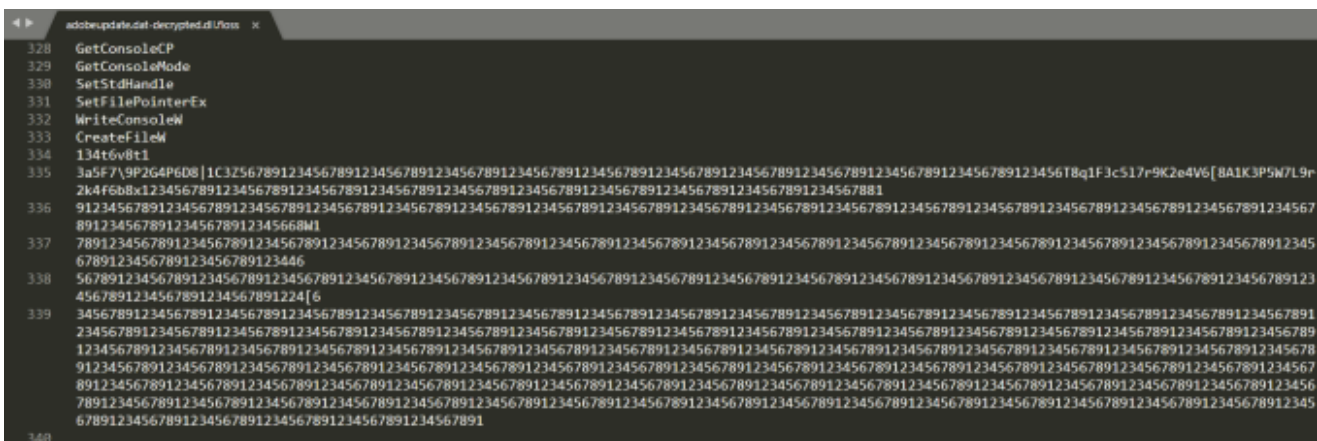
Looking at the payload in [DiE](#), we can see that it is a DLL file with a DLL name of **HT.dll**, 1 Export function named **Loader**, and a bunch of Import functions from **kernel32.dll** and **user32.dll**.





Interestingly enough, the Import functions for this sample are not that many and doesn't include typical functions you see for a RAT. Furthermore, the presence of **GetProcAddress** and **LoadLibraryA** potentially indicates that this sample dynamically resolves Win32 API functions at run time just like the loader.

Running **FLOSS** at the sample will give us a bunch of interesting strings to pivot from at IDA for a deeper analysis, but one interesting thing that immediately struck me was this blob with repeating pattern of "**123456789**".



By this time, some of you may already have an idea of what this is potentially and why the repeating pattern of **123456789** is so interesting for us. But to be sure, let's take a closer look at the disassembly in IDA.

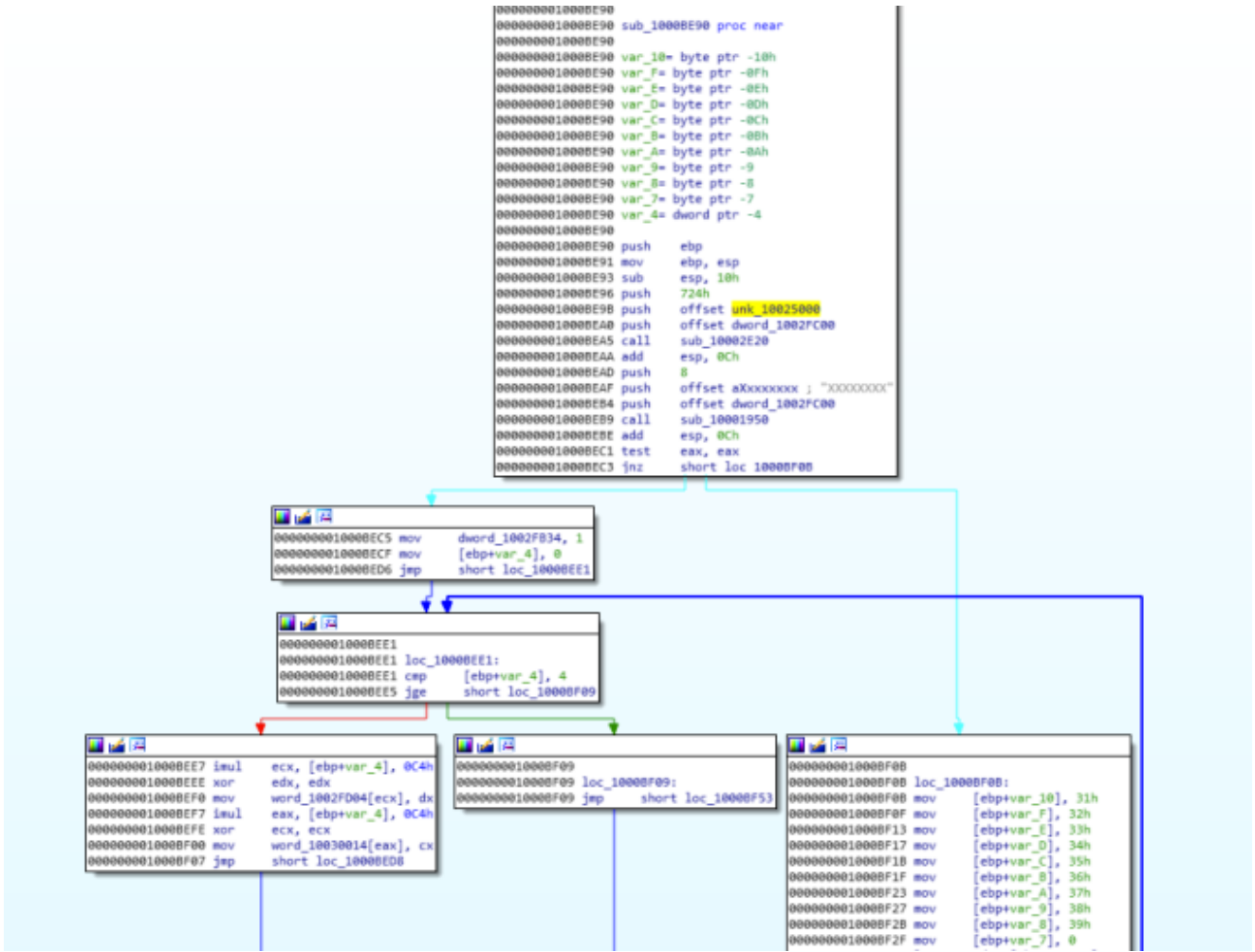
Clicking this on the **Strings window** of IDA will show us that this blob starts at offset 0 of the data section.

```

.data:10025000 ; Section 3. (virtual address 00025000)
.data:10025000 ; Virtual size           : 00000960 ( 55648.)
.data:10025000 ; Section size in file          : 00001800 ( 6144.)
.data:10025000 ; Offset to raw data for section: 00024400
.data:10025000 ; Flags C0000040: Data Readable Writable
.data:10025000 ; Alignment           : default
.data:10025000 ; =====
.data:10025000
.data:10025000 ; Segment type: Pure data
.data:10025000 ; Segment permissions: Read/Write
.data:10025000 _data          segment para public 'DATA' use32
.data:10025000          assume cs:_data
.data:10025000          ;org 10025000h
.data:10025000 unk_10025000 db 009h ; Û ; DATA XREF: sub_1000BE90+Bf0
.data:10025001          db 31h ; 1
.data:10025002          db 33h ; 3
.data:10025003          db 34h ; 4
.data:10025004          db 74h ; t
.data:10025005          db 36h ; 6
.data:10025006          db 76h ; v
.data:10025007          db 38h ; 8
.data:10025008          db 74h ; t
.data:10025009          db 31h ; 1
.data:1002500A          db 12h
.data:1002500B          db 33h ; 3
.data:1002500C          db 61h ; a
.data:1002500D          db 35h ; 5
.data:1002500E          db 46h ; F
.data:1002500F          db 37h ; 7
.data:10025010          db 5Ch ; \
.data:10025011          db 39h ; 9
.data:10025012          db 50h ; P
.data:10025013          db 32h ; 2
.data:10025014          db 47h ; G
.data:10025015          db 34h ; 4
.data:10025016          db 50h ; P
.data:10025017          db 36h ; 6
.data:10025018          db 44h ; D
.data:10025019          db 38h ; 8
.data:1002501A          db 7Ch ; |
.data:1002501B          db 31h ; 1
.data:1002501C          db 43h ; C
.data:1002501D          db 33h ; 3
.data:1002501E          db 5Ah ; Z
.data:1002501F          db 35h ; 5
.data:10025020          db 36h ; 6

```

Checking the x-refs of unk_10025000, will bring us to a function at **sub_1000BE90**.



As you can see, **unk_10025000** is pushed onto the stack along with **724h** and another memory offset **dword_1002FC00** as parameters for a function located at **sub_1000E20**.

```

000000001000BE90 push    ebp
000000001000BE91 mov     ebp, esp
000000001000BE93 sub     esp, 10h
000000001000BE96 push    724h
000000001000BE98 push    offset unk_10025000
000000001000BEA0 push    offset dword_1002FC00
000000001000BEA5 call   sub_1000E20

```

Looking closer at **sub_1000E20**, we can immediately figure out that this is a wrapper function that dynamically resolves the address of **memcpy** from **msvcrt** (**sub_100018B0**) via **GetProcAddress** and **LoadLibrary** (**sub_100018B0**) and then calls it with the parameters passed to the function.

```
0000000010002E20
0000000010002E20
0000000010002E20 ; Attributes: bp-based frame
0000000010002E20 sub_10002E20 proc near
0000000010002E20 ProcName= byte ptr -8
0000000010002E20 var_7= byte ptr -7
0000000010002E20 var_6= byte ptr -6
0000000010002E20 var_5= byte ptr -5
0000000010002E20 var_4= byte ptr -4
0000000010002E20 var_3= byte ptr -3
0000000010002E20 var_2= byte ptr -2
0000000010002E20 arg_0= dword ptr 8
0000000010002E20 arg_4= dword ptr 0Ch
0000000010002E20 arg_8= dword ptr 10h
0000000010002E20
0000000010002E20 push    ebp
0000000010002E21 mov     ebp, esp
0000000010002E23 sub     esp, 8
0000000010002E26 mov     [ebp+ProcName], 'm'
0000000010002E2A mov     [ebp+var_7], 'e'
0000000010002E2E mov     [ebp+var_6], 'm'
0000000010002E32 mov     [ebp+var_5], 'c'
0000000010002E36 mov     [ebp+var_4], 'p'
0000000010002E3A mov     [ebp+var_3], 'y'
0000000010002E3E mov     [ebp+var_2], 0
0000000010002E42 cmp     dword_10026850, 0
0000000010002E49 jnz     short loc_10002E60
```

```
0000000010002E4B lea    eax, [ebp+ProcName]
0000000010002E4E push   eax ; lpProcName
0000000010002E4F call  sub_10001800
0000000010002E54 push   eax ; hModule
0000000010002E55 call  ds:GetProcAddress
0000000010002E5B mov    dword_10026850, eax
```

```
0000000010002E60
0000000010002E60 loc_10002E60:
0000000010002E60 mov    ecx, [ebp+arg_8]
0000000010002E63 push   ecx
0000000010002E64 mov    edx, [ebp+arg_4]
0000000010002E67 push   edx
0000000010002E68 mov    eax, [ebp+arg_0]
0000000010002E6B push   eax
0000000010002E6C call  dword_10026850
0000000010002E72 add    esp, 0Ch
0000000010002E75 mov    esp, ebp
0000000010002E77 pop    ebp
0000000010002E78 retn
0000000010002E78 sub_10002E20 endp
0000000010002E78
```

```
00000000100018B0
00000000100018B0
00000000100018B0 ; Attributes: bp-based frame
00000000100018B0
00000000100018B0 sub_100018B0 proc near
00000000100018B0
00000000100018B0 LibFileName= byte ptr -8
00000000100018B0 var_7= byte ptr -7
00000000100018B0 var_6= byte ptr -6
00000000100018B0 var_5= byte ptr -5
00000000100018B0 var_4= byte ptr -4
00000000100018B0 var_3= byte ptr -3
00000000100018B0 var_2= byte ptr -2
00000000100018B0
00000000100018B0 push    ebp |
00000000100018B1 mov     ebp, esp
00000000100018B3 sub     esp, 8
00000000100018B6 mov     [ebp+LibFileName], 'm'
00000000100018BA mov     [ebp+var_7], 's'
00000000100018BE mov     [ebp+var_6], 'v'
00000000100018C2 mov     [ebp+var_5], 'c'
00000000100018C6 mov     [ebp+var_4], 'r'
00000000100018CA mov     [ebp+var_3], 't'
00000000100018CE mov     [ebp+var_2], 0
00000000100018D2 cmp     dword_10026814, 0
00000000100018D9 jnz     short loc_100018EA
```

```
00000000100018DB lea    eax, [ebp+LibFileName]
00000000100018DE push   eax ; lpLibFileName
00000000100018DF call  ds:LoadLibraryA
00000000100018E5 mov    dword_10026814, eax
```

```
00000000100018EA
00000000100018EA loc_100018EA:
00000000100018EA mov    eax, dword_10026814
00000000100018EF mov    esp, ebp
00000000100018F1 pop    ebp
00000000100018F2 retn
00000000100018F2 sub_100018B0 endp
00000000100018F2
```

In simpler terms, this function copies the first **1828** bytes of the data section (**unk_10025000**) to another memory location (**dword_1002FC00**).

After that, **dword_1002FC00** is pushed onto the stack along with **8** and another memory offset **aXXXXXXX** as parameters for a function located at **sub_10001950**. This function is quite similar with the previous one but dynamically resolves the address of **memcmp** from **msvcrt** and then calls it with the parameters passed to the function.

```

0000000010001950 push    ebp
0000000010001951 mov     ebp, esp
0000000010001953 sub     esp, 8
0000000010001956 mov     [ebp+ProcName], 'm'
000000001000195A mov     [ebp+var_7], 'e'
000000001000195E mov     [ebp+var_6], 'm'
0000000010001962 mov     [ebp+var_5], 'c'
0000000010001966 mov     [ebp+var_4], 'm'
000000001000196A mov     [ebp+var_3], 'p'
000000001000196E mov     [ebp+var_2], 0
0000000010001972 cmp     dword_10026830, 0
0000000010001979 jnz     short loc_10001990

```

```

000000001000197B lea    eax, [ebp+ProcName]
000000001000197E push   eax           ; lpProcName
000000001000197F call   sub_100018B0
0000000010001984 push   eax           ; hModule
0000000010001985 call   ds:GetProcAddress
000000001000198B mov     dword_10026830, eax

```

```

0000000010001990
0000000010001990 loc_10001990:
0000000010001990 mov     ecx, [ebp+arg_8]
0000000010001993 push   ecx
0000000010001994 mov     edx, [ebp+arg_4]
0000000010001997 push   edx
0000000010001998 mov     eax, [ebp+arg_0]
000000001000199B push   eax
000000001000199C call   dword_10026830
00000000100019A2 add     esp, 0Ch
00000000100019A5 mov     esp, ebp
00000000100019A7 pop     ebp
00000000100019A8 retn
00000000100019A8 sub_10001950 endp
00000000100019A8

```

Again in simpler terms, this function compares the first 8 bytes of **dword_1002FC00** to “XXXXXXXX”. The result of the comparison determines the execution path of the malware.

Note: The malware uses a lot of wrapper functions like this to dynamically resolve the address of various Win32 API functions from different DLLs via **GetProcAddress** and **LoadLibrary** and then execute it. This is one of the anti-detection measures implemented by this specific PlugX variant and is common through out the whole binary.

Let’s focus on the execution path where it doesn’t match “XXXXXXXX”. As you can see, it pushes the string “123456789” onto the stack and calls a function at **sub_10002DC0**.


```

000000001000BF0B
000000001000BF0B loc_1000BF0B:
000000001000BF0B mov     [ebp+var_10], '1'
000000001000BF0F mov     [ebp+var_F], '2'
000000001000BF13 mov     [ebp+var_E], '3'
000000001000BF17 mov     [ebp+var_D], '4'
000000001000BF1B mov     [ebp+var_C], '5'
000000001000BF1F mov     [ebp+var_B], '6'
000000001000BF23 mov     [ebp+var_A], '7'
000000001000BF27 mov     [ebp+var_9], '8'
000000001000BF2B mov     [ebp+var_8], '9'
000000001000BF2F mov     [ebp+var_7], 0
000000001000BF33 lea    edx, [ebp+var_10]
000000001000BF36 push   edx
000000001000BF37 call   sub_10002DC0
000000001000BF3C push   eax
000000001000BF3D lea    eax, [ebp+var_10]
000000001000BF40 push   eax
000000001000BF41 push   724h
000000001000BF46 push   offset dword_1002FC00
000000001000BF4B call   sub_1000B840
000000001000BF50 add    esp, 10h

```

This is another wrapper function but this time, for **lstrlenA** to find the length of the string passed as a parameter. The resulting string length is stored in **EAX** and is pushed onto the stack along with the string “**123456789**“, **724h**, and **dword_1002FC00** which contains the blob we noted earlier before a call to **sub_1000B840** is made.

```

000000001000BF0B
000000001000BF0B loc_1000BF0B:
000000001000BF0B mov     [ebp+var_10], '1'
000000001000BF0F mov     [ebp+var_F], '2'
000000001000BF13 mov     [ebp+var_E], '3'
000000001000BF17 mov     [ebp+var_D], '4'
000000001000BF1B mov     [ebp+var_C], '5'
000000001000BF1F mov     [ebp+var_B], '6'
000000001000BF23 mov     [ebp+var_A], '7'
000000001000BF27 mov     [ebp+var_9], '8'
000000001000BF2B mov     [ebp+var_8], '9'
000000001000BF2F mov     [ebp+var_7], 0
000000001000BF33 lea    edx, [ebp+var_10]
000000001000BF36 push   edx
000000001000BF37 call   sub_10002DC0
000000001000BF3C push   eax
000000001000BF3D lea    eax, [ebp+var_10]
000000001000BF40 push   eax
000000001000BF41 push   724h
000000001000BF46 push   offset dword_1002FC00
000000001000BF4B call   sub_1000B840
000000001000BF50 add    esp, 10h

```

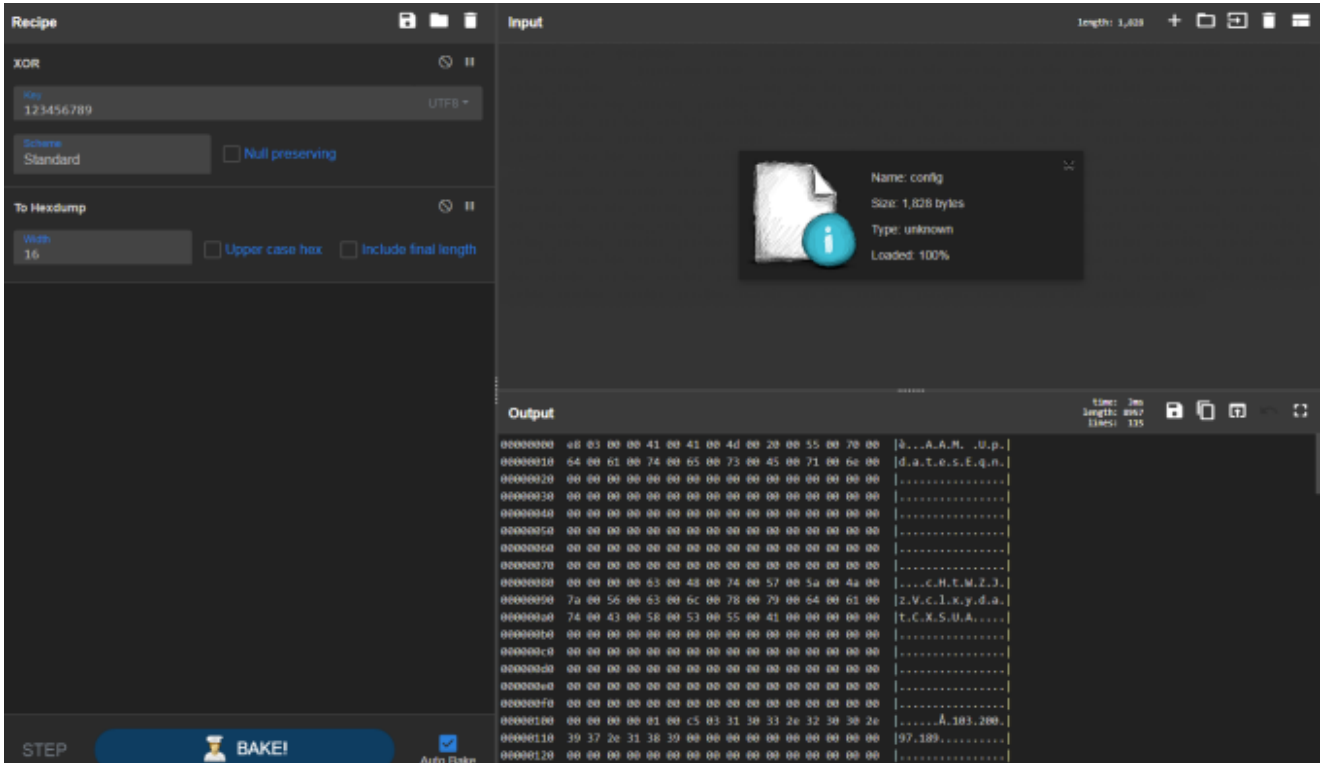
Looks familiar right? Yes, this is the same format of parameters we noted on the loader before the call to the decryption function – the key length, the key, the length of the encrypted data, and the address of the encrypted data.

Looking closer at **sub_1000B840** and due to the fact that there are repeating patterns of **123456789** on the encrypted blob, we can immediately recognize that the algorithm used to encrypt it is XOR with multi-byte key as well.



So by extracting the first **1828** bytes of the data section and performing a multi-byte XOR on it with the key **123456789**, we can decrypt the blob and see if it really contains the malware configuration.

Let's try it with a quick Cyberchef recipe:



As you can see, there are a few interesting information in here such as:

- * An Adobe themed unicode string “**AAM UpdatesEqn**”
- * A random looking unicode string “**cHtWZJzVclxydatCX SUA**”
- * IP addresses that are likely the C2 addresses

Each IP address entry looks like it starts with **01 00**, then followed by two bytes which are likely port numbers in hex, and then finally the IP address itself.

If we convert the port numbers in hex to decimal (little endian) and the IP addresses in hex to ascii we will get the following addresses:

- 103.200.97[.]189:965
- 103.200.97[.]189:110
- 185.239.226[.]17:965
- 185.239.226[.]17:110

Let’s try to confirm our findings by running the sample on a VM and monitoring for any interesting events using some dynamic analysis tools.

File, registry, and process events captured by Procmon:

Time of Day	Process Name	PID	Operation	Path	Result
12:25:47.8104818 PM	AdobeInstall.exe	3464	RegCreateKey	HKCR\ms-pu	SUCCESS
12:25:47.8194847 PM	AdobeInstall.exe	3464	WriteFile	C:\ProgramData\AAM UpdatesEqn\AAM Updates.exe	SUCCESS
12:25:47.8195258 PM	AdobeInstall.exe	3464	WriteFile	C:\ProgramData\AAM UpdatesEqn\AAM Updates.exe	SUCCESS
12:25:47.8195747 PM	AdobeInstall.exe	3464	WriteFile	C:\ProgramData\AAM UpdatesEqn\AAM Updates.exe	SUCCESS
12:25:47.8227567 PM	AdobeInstall.exe	3464	WriteFile	C:\ProgramData\AAM UpdatesEqn\hex.dll	SUCCESS
12:25:47.8258334 PM	AdobeInstall.exe	3464	WriteFile	C:\ProgramData\AAM UpdatesEqn\adobeupdate.dat	SUCCESS
12:25:47.8258685 PM	AdobeInstall.exe	3464	WriteFile	C:\ProgramData\AAM UpdatesEqn\adobeupdate.dat	SUCCESS
12:25:47.8259203 PM	AdobeInstall.exe	3464	WriteFile	C:\ProgramData\AAM UpdatesEqn\adobeupdate.dat	SUCCESS
12:25:47.8263681 PM	AdobeInstall.exe	3464	RegCreateKey	HKLM\Software\Wow6432Node\Microsoft\Windows\CurrentVersion\Run	ACCESS DENIED
12:25:47.8264076 PM	AdobeInstall.exe	3464	RegCreateKey	HKLM\Software\Wow6432Node	SUCCESS
12:25:47.8264242 PM	AdobeInstall.exe	3464	RegCreateKey	HKLM\SOFTWARE\Wow6432Node\Microsoft	SUCCESS
12:25:47.8264495 PM	AdobeInstall.exe	3464	RegCreateKey	HKLM\SOFTWARE\Wow6432Node\Microsoft\Windows	SUCCESS
12:25:47.8264720 PM	AdobeInstall.exe	3464	RegCreateKey	HKLM\SOFTWARE\Wow6432Node\Microsoft\Windows\CurrentVersion	SUCCESS
12:25:47.8264900 PM	AdobeInstall.exe	3464	RegCreateKey	HKLM\SOFTWARE\Wow6432Node\Microsoft\Windows\CurrentVersion\Run	ACCESS DENIED
12:25:47.8265332 PM	AdobeInstall.exe	3464	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Run	SUCCESS
12:25:47.8265583 PM	AdobeInstall.exe	3464	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Run\AAM UpdatesEqn	SUCCESS
12:25:47.8272756 PM	AdobeInstall.exe	3464	WriteFile	C:\ProgramData\AAM UpdatesEqn\AAM Updates.exe	SUCCESS
12:25:47.8330298 PM	AdobeInstall.exe	3464	Process Create	C:\ProgramData\AAM UpdatesEqn\AAM Updates.exe	SUCCESS

From the Procmon output, we can see that the malware:

- * Created a new folder named **AAM UpdatesEqn** on the Program Data directory and copied the PlugX components to it
- * Attempted to access **HKLM\SOFTWARE\Wow6432Node\Microsoft\Windows\Current Version\Run** for registry write operation but failed
- * Succeeded in accessing **HKCU\SOFTWARE\Microsoft\Windows\Current Version\Run** for registry write operation
- * Created a new value named **AAM UpdatesEqn**. This is a persistence mechanism which will execute “**C:\ProgramData\AAM UpdatesEqn\AAM Updates.exe**” **701** across system reboots
- * Created a new process for **C:\ProgramData\AAM UpdatesEqn\AAM Updates.exe** passing the value **701** as parameter

HTTP request captured by fakenet:

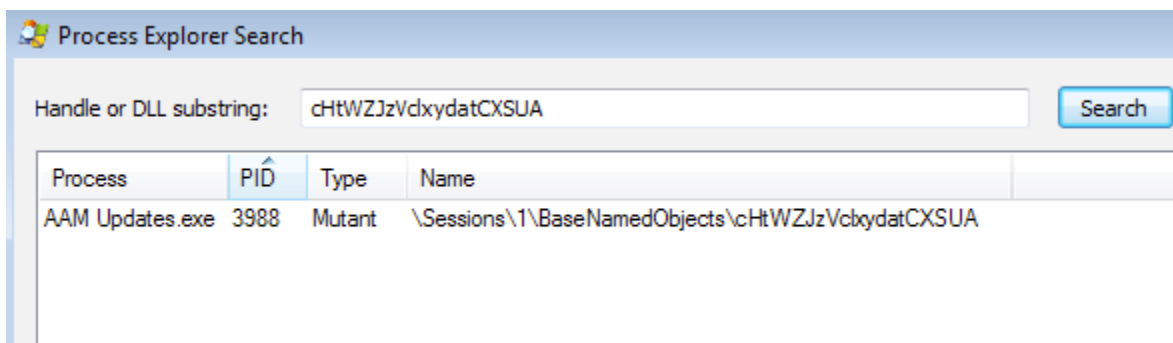
```

06/28/20 12:28:50 PM [ HTTPListener80]
06/28/20 12:28:50 PM [ HTTPListener80] Storing HTTP POST headers and data to http_20200628_122850.txt.
06/28/20 12:28:52 PM [ HTTPListener80] POST /update?ud=06304bb6 HTTP/1.1
06/28/20 12:28:52 PM [ HTTPListener80] Accept: */*
06/28/20 12:28:52 PM [ HTTPListener80] x-debug: 0
06/28/20 12:28:52 PM [ HTTPListener80] x-request: 0
06/28/20 12:28:52 PM [ HTTPListener80] x-content: 61456
06/28/20 12:28:52 PM [ HTTPListener80] x-storage: 1
06/28/20 12:28:52 PM [ HTTPListener80] User-Agent: Mozilla/4.0 (compatible; MSIE 6.0; Windows NT 5.1;SUI;
06/28/20 12:28:52 PM [ HTTPListener80] Host: 103.200.97.189:965
06/28/20 12:28:52 PM [ HTTPListener80] Content-Length: 0
06/28/20 12:28:52 PM [ HTTPListener80] Connection: Keep-Alive
06/28/20 12:28:52 PM [ HTTPListener80] Cache-Control: no-cache
06/28/20 12:28:52 PM [ HTTPListener80]

```

From the Fakenet output, we can see that the malware communicated to one of the IP address and port combinations we noted earlier.

Searching for cHtWZJzVclxydatCXSUA handle in Procexp:



Lastly, we can see that **cHtWZJzVclxydatCXSUA** is a mutex object used by the malware.

These observations confirm that we were successful in extracting the configuration of the malware.

Automating the config extraction

To make our lives easier, I created this quick-and-dirty python script to automatically extract the configuration information for this variant:

I've tested this script on a bunch of samples and it seems to work fine. However, there can be instances where it won't work if the config is structured differently.

That's it guys! I really hope you learned something new today and as always, thank you for reading my blog.

Tags: [Malware](#), [Mustang Panda](#), [PlugX](#), [RAT](#), [Reverse Engineering](#)