

# Breaking the B0 ransomware: Investigation & Decryption

[porthas.com/blog/b0-ransomware-decryption/](https://porthas.com/blog/b0-ransomware-decryption/)

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## Breaking the unbreakable: Story of B0 software group ransomware

This is how deep technical curiosity, DFIR expertise, and persistence led to the downfall of B0 Software Group's ransomware. In this article, we walk through the full investigation from start to finish: from the initial response and behavioral analysis to reverse engineering and decryptor development. With annotated code snippets, screenshots, and step-by-step breakdowns, we reveal not just what was done, but why it worked.

If you're looking for a real-world look at how threat intelligence and malware analysis can turn the tide in ransomware recovery, this is it.

## Known Tactics, Techniques, and Procedures

Since ransomware groups can often sell their malware to multiple affiliates. These affiliates can have different flavors and attack methodologies for how they breach a targeted network. As such, there is no single set of TTPs that can generally cover how B0 software group ransomware can end up on your network. In this case, however, the DFIR investigation highly suggests that brute-forcing an internet-facing RDP (RDP brute-forcing) was the initial entry point.

### The assessment was based on the following factors:

- Evidence of a brute force attack. Found repeated failed logins to the Administrator account on the compromised server. The brute force attempts originated from external IP addresses.

- The default RDP port 3389 was directly exposed to the internet at the time of the incident. The default RDP port is frequently an attractive attack vector and can be easily discovered by automated port scanners such as Nmap.

## DFIR Investigation Conclusion

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As for activities within the network and lateral movement, the DFIR investigation concluded the following points:

- Network discovery/reconnaissance. Upon gaining access to the server, the threat actor used the network reconnaissance tools Advance IP Scanner and kportscan 3.0 to identify additional devices to which to move laterally. However, because the server was the only device within a segmented, isolated environment, further lateral movement was unsuccessful.
- Account usage and privilege escalation. The threat actor gained access to the privileged Administrator account on the server via a brute force attack. Upon successfully logging into the system, they also created the account B0Us. The threat actor also executed the credential harvesting tool, Mimikatz, on the server. This is a credential dumping tool, typically used by threat actors to obtain any available credentials stored within the memory of a target system.
- Ransomware deployment. The threat actor transferred the compressed file 'pth.rar' to the system, which contained their toolkit, including the ransomware executable 'b0.exe'. The compressed file was initially placed in the directory 'C:\users\b0Us\desktop\' and then subsequently extracted using the file compression tool, WinRAR. The threat actor placed their ransomware binary within the folder 'C:\users\public', manually executing this on the server leading to the encryption of files on the system and creation of ransom notes.

## Exfiltration Analysis

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We assess that the threat actor **did not exfiltrate data** from the impacted server. This assessment is based on the following factors:

- No evidence of file transfer tooling. S-RM found no evidence of data transfer tooling that could have been used by the threat actor to exfiltrate data from the server.
- No evidence of data access or staging activities. S-RM reviewed data extracted from the system and did not identify any conclusive evidence typically preceding exfiltration, such as data reconnaissance and compression. While we found evidence of interaction with the server's data drives, it is likely that this occurred in the context of pre-encryption activities.
- Threat actor claims. The threat actor did not claim to have exfiltrated any data from the server in their ransom note. Further, S-RM found no evidence of a leak site maintained by the threat actor to publish victim data as part of data extortion tactics.

## Limitations

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- Due to the hosted nature of the server, S-RM could not obtain access to network or firewall logs, typically providing insight into data traffic and volumes. We were therefore unable to determine whether there any significant spikes in outbound data during the timeframe of the incident consistent with exfiltration activities.
- The evidence indicates that a limited number of files were accessed by the threat actor, however there is no accompanying evidence that would show that this data was transferred out of the system.

## Technical analysis

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The ransomware was written in Golang, making it more challenging to reverse engineer compared to other ransomware variants which are in C/C++.

The compilation timestamp of the analyzed sample was set to zero, which might have been a tactic used by the authors to confuse analysts and make it more difficult to determine when the sample was compiled. Notably, the **ransomware includes its own kill switch**. It calculates the time difference between a hardcoded UTC timestamp of **March 24, 2025, at 00:00:00**, and the time when the sample is executed. The ransomware checks if this time difference falls within a 3-day range. If it does, the ransomware continues its normal execution. If not, it immediately restarts the system by executing the command “**shutdown /r /t 0**”.

```
start = time.Date(2025, 3, 24, 0, 0, 0, 0, time.UTC[0]); // March 24 2025 0:0:0 UTC
v23 = v0;
current_time = time.Now();
time_difference = time.Time.Sub(current_time, 3, v2, start, 3, v23, v3, v4, v5);
if ( time_difference <= 0xEBBDB3ED0000LL ) // check if time_difference is within 3 day range
{
    main_HavePassword(time_difference, 3LL, v7, start, 3); // proceed with execution
}
else
{
    v24[0] = &RTYPE_string;
    v24[1] = &off_10FFAE0;
    v12 = os.Stdout;
    v13 = fmt.Fprintln(
        (unsigned int)go_itab__os_File_io_Writer,
        os.Stdout,
        (unsigned int)v24,
        1,
        1,
        v8,
        v9,
        v10,
        v11,
        v19,
        v21);
    ((void ( __golang *) ( __int64 ))main_restartSystem)(v13); // shutdown /r /t 0 (force immediate restart of the system)
    time.Sleep(0x540BE400, v12, v14, 1, 1, v15, v16, v17, v18, v20);
}
```

After confirming that the process runs within the allowed time range, it moves on to another check. During this check, it prompts the user for a password and a network ID. The provided password is then hashed using SHA-256 and compared against a hardcoded

hash. If the hashes match, the execution will proceed; otherwise, it will prompt for the password again, up to three times. If all three attempts fail, the process will exit.

```
hashed_password = main_hashPasswordLogin(v77, v74, v49, 1, 1, v58, v51, v52, v53);
if ( hash_length == (char *)0x40 )
{
    hash_length = "de2b8de1ca70f7d905bb8e4e7717a42af266e0f35a2c3cbf0a45061c9a700615f1ate: invalid compression level %d: want value in range [-2, 9]json: invalid number :
    if ( (unsigned __int8)runtime_memequal(
        hashed_password,
        (QWORD)"de2b8de1ca70f7d905bb8e4e7717a42af266e0f35a2c3cbf0a45061c9a700615f1ate: invalid compression level %d: want value in range [-2, 9]json
    {
        v82[0] = &RTYPE_string;
        v82[1] = &off_10FFB10;
        fmt_Fprintln(
            (unsigned int)go_itab__os_File_io_Writer,
            os_Stdout,
            (unsigned int)v82,
            1,
            1,
            v55,
            v56,
            v57,
            v58,
            v67,
            v68);
        return main_HaveRun(v78, v75);        // proceed with execution
    }
}
--
```

Upon successfully completing all the previously mentioned checks, the process moves on to configure the setup. First, it initializes a pseudo-random number generator using the current system time obtained by calling “time.Now.UnixNano().” Next, it generates a random nine-character string that will later be used for the key derivation function. This key will be crucial for the decryption and recovery of ransomware, which we will discuss later.

```
current_time = time.Now();        // obtain current system time
v97 = current_time;
v98 = a2;
if ( current_time < 0 )
{
    v98 = ((unsigned __int64)(2 * current_time) >> 31) + 0xD07B17F80LL;
    v97 = current_time & 0x3FFFFFFF;
}
v3 = v98;
if ( v97 < 0 )
{
    v3 = ((unsigned __int64)(2 * v97) >> 31) + 0xD07B17F80LL;
    math_rand_ptr_Rand_Seed(math_rand_globalRand, (v97 & 0x3FFFFFFF) + 1000000000 * v3 - 0x5E4DFC14C2E6000LL); // Seed or initialize the pseudo-random number generator
    RandomString_RandomString = B0Client_GetRandomString_RandomString( // Generate 9 character key string
        9,
        (v97 & 0x3FFFFFFF) + 1000000000 * v3 + 1025114112,
        v4,
        v97 & 0x3FFFFFFF,
        v3,
        v5,
        v6,
        v7,
        v8);
}
```

Next, it will verify the existence of the JSON file “telemetry.ASM-WindowsDefault.json” located in “C:\Windows\System32”. This JSON file contains the ransomware configuration, including the initial key (a random nine-character string), the extension, and the Network ID.

```
{ "configure": [ { "extension": "B0-3e72d", "key": " ", "network": " " } ] }
```

*JSON config file containing the extension, key and the network ID.*

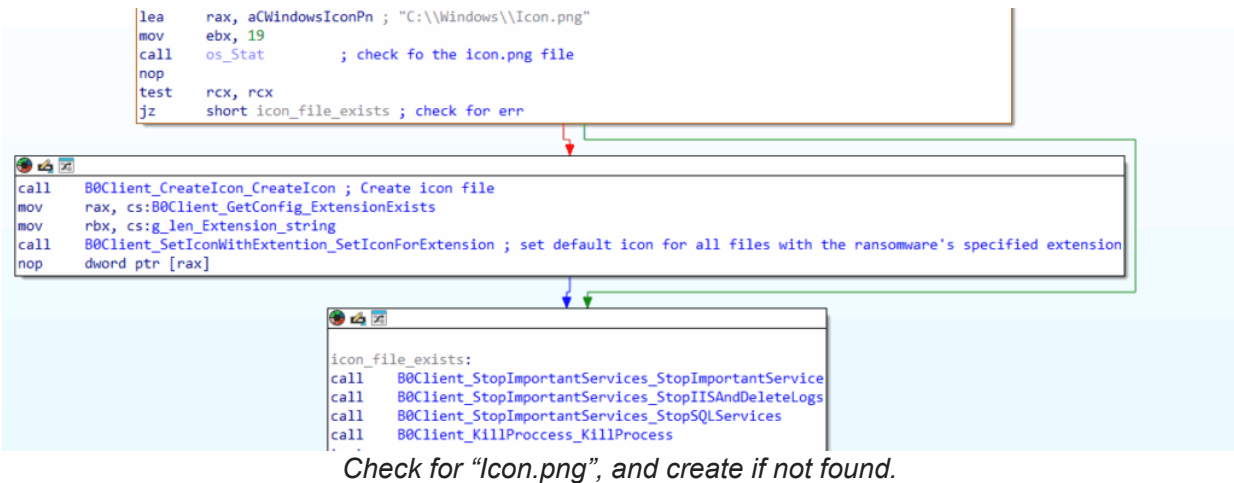
If the file is found, it continues by reading the file and setting the corresponding global variables for ransomware configuration.

```
v9 = os.Stat((unsigned int)&aUseOfClosedNet[14772], 53, a3, a4, a5, a6, a7, a8, a9, v17); // check the existence of key file
if ( !v10 )
    return B0Client_GetConfig_ReadKey(v9, 53LL, 0LL, a4, a5, v11, v12, v13, v14); // Read key file and set the corresponding global variables for ransomware configuration
v15 = qword_1108198;
result = os_underlyingErrorIs(v10, a4, os_ErrNotExist, qword_1108198, qword_1108198, v11, v12, v13, v14, v18);
if ( (_BYTE)result )
    return B0Client_GetConfig_CreateKeyFile(a1, a2, a3, a4, v15); // Create key file and set ransomware configuration globals
```

*Configuration check and setup.*

Similar to the key file, it will check for the existence of a file “Icon.png”, located in the Windows directory “C:\Windows”. Same as before, if the file is not found, it proceeds to create it. The icon file contents are stored as base64 encoded text that will be decoded and written to the file at run-time.

If the file was not found, it then creates the file “Icon.png” and sets the default icon for files with the ransomware’s extension (as specified in the ransomware configuration) to this “Icon.png.”



## Process and Service targeting

Before proceeding with key derivation and encryption, the ransomware will attempt to stop related services for Windows IIS, MSSQL, and MySQL by executing the commands listed below.

### BlockListed services

- cmd /C iisreset stop
- cmd /C NET STOP IISADMIN
- cmd /C net stop WAS
- cmd /C NET stop MSSQLSERVER
- cmd /C NET stop \"SQL Server (MSSQLSERVER)\"
- cmd /C net stop MSSQL\$SQLEXPRESS
- cmd /C net stop SQLSERVERAGENT
- cmd /C net stop mysql
- sc stop W3SVC
- sc stop WAS
- sc stop IISADMIN



It's important to note that not only it targets specific services by name but also attempts to obtain a list of running services and processes using the commands "sc query" and "tasklist." Additionally, it tries to terminate any service or process that contains the word "SQL" in its name.

## Anti-forensics measures

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The ransomware deploys a variety of anti-forensic measures, to avoid leaving behind any leads that can be used in a forensic investigation. The following is a list of the anti-forensics' measures taken by the ransomware to wipe any evidence.

1. Delete the Windows IIS logs located in the directory: **C:\inetpub\logs**.
2. Empty the Recycle Bin.
3. Constantly clears Windows Event Logs.
4. Constantly Clear the contents of the Windows Temp folder by running the following command:

```
Del /S /F /Q %windir%\Temp
```

5. Create a registry entry under the Run key to clear Windows Event Logs upon user login. Use the following command:

```
powershell.exe -Command New-ItemProperty -Path  
HKCU:\Software\Microsoft\Windows\CurrentVersion\Run -Name DeleteLogWithStartUp -  
Value 'powershell wevtutil el | ForEach-Object { wevtutil cl $_ }' -PropertyType  
String
```

This will ensure that the Windows Event Logs are cleared each time the user logs in.

Furthermore, the following commands are written to a file called "Del.cmd," which is saved in the system directory C:\Windows\System32\ and executed. This script will delete the ransomware executable as well as the Windows Event logs.

## Contents of Del.cmd, that will be executed after encryption to wipe any traces of its activity

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```
timeout /t 10  
del /f ransomware_executable_path  
powershell "wevtutil el | Foreach-Object {wevtutil cl "$_"}
```

It is noteworthy to mention that the ransomware does not delete the Windows Event Logs; instead, it continuously clears them during execution to erase any traces of its activity. The same behavior applies to the Windows Temp folder.

```

call    main_FindFile
call    B0Client_RunCmdExecutable_DelTemp ; wipe the Windows Temp folder
call    B0Client_RunCmdExecutable_ClearLog ; clear event logs
call    B0Client_RunCmdExecutable_DeleteEndedFiles ; execute contents of Del.cmd to clear the event logs and delete the ransomware executable after encryption

```

## Key Derivation

The initially generated random nine-character string is used to generate the AES-GCM key, which is then used during encryption. The ransomware first uses the hashids library to generate a hashed ID using the random string as the salt and 1,000,000,000(0x3B9ACA00). This library can generate hashed/obfuscated IDs from numbers. The hashed ID's size is 7 bytes, and its character set is limited to 'abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ1234567890'.

```

hd[0] = &abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ1234567890;
hd[1] = 0x3ELL;
hd[4] = saltLen;
hd[3] = Key; // salt
hd[2] = 1LL;
v9 = github_com_speps_go_hashids_NewWithData(
    (unsigned int)hd,
    saltLen,
    (unsigned int)&abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ1234567890);
id = 1000000000LL;
return github_com_speps_go_hashids_ptr_HashID_Encode(v9, (__int64)&id);

```

The hashed ID is then hashed in the following order: SHA1 → MD5 → SHA512 → SHA1 → MD5 → MD5 → SHA1 to generate a password that will be used in the password based key derivation function. The hashing functions also encode the hashed data before returning it.

```

__int64 __golang B0Client_Hashing_GetCode(int hashedId)
{
    int Sha1Hash; // eax
    _OWORD *MD5Hash; // rax
    int sha512Hash; // eax
    int Sha1Hash_1; // eax
    int MD5Hash_1; // eax
    int MD5Hash_2; // eax

    Sha1Hash = B0Client_Hashing_Sha1Generator(hashedId);
    MD5Hash = B0Client_Hashing_GetMD5Hash(Sha1Hash);
    sha512Hash = B0Client_Hashing_sha512String(MD5Hash);
    Sha1Hash_1 = B0Client_Hashing_Sha1Generator(sha512Hash);
    MD5Hash_1 = B0Client_Hashing_GetMD5Hash(Sha1Hash_1);
    MD5Hash_2 = B0Client_Hashing_GetMD5Hash(MD5Hash_1);
    return B0Client_Hashing_Sha1Generator(MD5Hash_2);
}

```

After this, PBKDF2 is used to derive the AES-GCM key from the password. The salt is the same as the password; 100,000 iterations are used, and the key length is 32 bytes. SHA512 is used as the hash function for the password based key derivation function. Notably, if a null value is passed as the salt(that indicates for some reason it failed to generate and hash the random nine-character string), the ransomware generates a random salt using the crypto/rand package; however, if this occurs, it is impossible to decrypt the file



even with the threat actor's decryptor, since the ransomware does not store the salt somewhere. Regardless, this scenario would not happen, as the third argument is always the password converted to bytes in the current version of the ransomware.

```
if ( !salt )
{
    n32_1 = *password2;
    salt_2 = runtime_makeslice(&RTYPE_uint8, 32, 32, n32_2, n32, v5, v6, v7, v8);
    crypto_rand_Read(salt_2, 32, 32, n32_2, n32, v9, v10, v11, v12);
    *password2 = n32_1;
    n32 = 32LL;
    n32_2 = 32LL;
    salt = salt_2;
}
salt = salt;
password2_1 = password2;
password = runtime_stringtoslicebyte(0LL, *password2, *a2, n32_2, n32, a2, v6, v7, v8);
return golang_org_x_crypto_pbkdf2_Key(password, password2_1, v16, salt, n32_2, n32, 100000LL, 32uLL, &sha512_New);
```

The resulting key is used to encrypt the files with AES-GCM, and the same key is used for each file.

## Ransom Note


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As is common to ransomware groups, a note with instructions on how to contact the actor is dropped in each driver's root directory. The victim can negotiate payment options via a communication channel of the actor's choosing. The note is dropped in two forms : a text file, `How_To_Recovery_Files.txt`, and an HTML file also named `How_To_Recovery_Files.html` with slightly different content.


It is worth mentioning that the Network ID(serves a unique identifier to the victim) provided earlier in the early of its execution, is added to the contents of the ransom note before it's written to disk.

```
os_Stat(&C_DriveRansomnotePath, 28);
if ( err )
    B0Client_MessageSoft_MessageRan(          // drops ransomnotes in the root directory for each drive letter.
        B0Client_GetConfig_NetWorkId,
        n35,
        err,
        password,
        n32_3,
        v68,
        v69,
        v70,
        v71);
File = main_FindFile(initial_key, n32, aes_key, size, a5);
v73 = B0Client_RunCmdEexecutable_DelTemp(File);
v74 = B0Client_RunCmdEexecutable_ClearLog(v73);
return B0Client_RunCmdEexecutable_DeleteEndedFiles(v74);
```


*How the ransom notes are dropped in the root directory of each drive letter.*




Your unique system id : REDACTED

**What happened** 


Your computers and servers are encrypted in the B0 group exclusive algorithm, backups are deleted. We use strong encryption algorithms, so you cannot decrypt your data. Each system in the network is encrypted with a unique identifier. Our algorithm is based on international security standards and uses advanced encryption methods that keep facial information secure and without the possibility of unauthorized access. But you can restore everything by purchasing a special program from us - universal decryptor.

**Risks caused by non-cooperation** Sure, here are the 18 potential consequences of data loss in your business 

1- Loss of trust and satisfaction of customers 2- Damage to the company's reputation and trust 3- Incurrence of legal and compensatory costs 4- Loss of competitive advantage 5- Increased security risks and data protection concerns 6- Damage to stock value and business value 7- Negative impact on relationships with regulatory bodies and legal authorities 8- Increased risk of contractual breaches 9- Harm to company communications and public relations 10- Delays in operations and business processes 11- Reduced employee productivity and mental health issues 12- Damage to the company's brand and market recognition 13- Breach of contractual agreements and obligations 14- Issues related to transportation and distribution 15- Delays in programs and projects 16- Extra expenses for transparency and public disclosure 17- Increased risk of fines and legal problems 18- Growing employee concerns and stress

**What guarantees you have** 

We are an international group without political or religious affiliations and members from different parts of the world join us. In this group, our goal is to earn money and achieve financial success without discrimination based on political opinions or ideology, and we care about nothing but your money. You can send some encrypted files from your network, and we will decrypt it for free. We value our reputation. If we do not fulfill our work and obligations, no one will pay us. It's not in our interest. All of our decryption software is perfectly tested and will decrypt your data.

**Attention!** 

Do not rename encrypted files. Do not try to decrypt your data using third party software, it may cause permanent data loss. Decryption of your files with the help of third parties may cause increased price (they add their fee to our) or you can become a victim of a scam.

*HTML version of the ransomware note.*

```
#B0 Group

Your unique System Id: REDACTED
How to Contact with us?
1. Download and install Tox (https://tox.chat/download.html).
2. Tox ID Support: E3B88B09D68CF0AE576BC843D000300ED939068E473D952FF7F1C1B2D246507D082117729A3B

Tox is a decentralized, open-source messaging platform that uses peer-to-peer (P2P) protocols and end-to-end encryption to ensure user security and privacy. This platform allows for sending text messages, making voice and video calls, and transferring files without the need for centralized servers or registration. However, due to its peer-to-peer nature, the quality of communication on Tox heavily depends on users having a stable and fast internet connection. Without reliable internet access, users may experience issues such as call drops or delays in message delivery.

We have published all the steps and decryption protocols, along with a comprehensive video to verify the process, on our website to ensure transparency and build your trust.

Tor Website Address: http://bzsn5crutff7eig5mlohn2n2lkaqcafp6tsr6g5b3y4rxq4g32bg3uid.onion

To open Tor websites, you need to download and install the Tor browser. You can download it from the official Tor website at https://www.torproject.org.

If your system password has been changed, it is now REDACTED
```

*Text version of the ransomware note.*

## File Encryption

The *main\_FindFile* function retrieves a list of the root directory paths for each available drive letter, and then calls *main\_run* on each.

```

initial_key_size = max;
drive = B0Client_GetAllDrivesList_GetDrives();
n32_2 = max;
idx = 0LL;
while ( max > idx )
{
    n32_4 = idx;
    drive_1 = drive;
    main_run(drive->path.ptr, drive->path.len, initial_key, initial_key_size, aes_key, size, a5, v6, v7);
    drive = &drive_1->drive2;
    idx = n32_4 + 1;
    max = n32_2;
}
return drive;

```

The *main\_run* uses *filepath.Walk()* to recursively walk through all the files, and directories. The callback function passed as an argument to *filepath.Walk()* performs some checks on the file path to decide whether the file can be encrypted or not. It first checks whether any of the following blacklisted paths are in the filepath, and if they are, the callback function returns without encrypting the file.

```

if ( !result )
{
    (loc_774464)(v58, &off_902C20); // \WindowsPowerShell\Modules\
    :Windows = runtime_concatstring2(&v54, a10, a11, &unk_8618BE, 9); // : \Windows
    v60 = a10;
    :ProgramDataCertOpenStore = runtime_concatstring2(&v53, a10, a11, &unk_862775, 13); // : \ProgramData
    v62 = a10;
    :Program_FilesReference_Assemblies = runtime_concatstring2(0, a10, a11, &unk_869AF2, 37); // : \Program Files\Reference Assemblies\
    v64 = a10;
    :Program_FilesCommon_Files = runtime_concatstring2(&v52, a10, a11, &unk_86765C, 29); // : \Program Files\Common Files\
    v66 = a10;
    :Program_FilesInternet_Explorer = runtime_concatstring2(0, a10, a11, &unk_868F06, 34); // : \Program Files\Internet Explorer
    v68 = a10;
    :Program_Files_x86Reference_Assemblies = runtime_concatstring2(0, a10, a11, &unk_86AE57, 43); // : \Program Files (x86)\Reference Assemblies\
    v70 = a10;
    :Program_Files_x86Common_Files = runtime_concatstring2(0, a10, a11, &unk_869327, 35); // : \Program Files (x86)\Common Files\
    v72 = a10;
    :Program_Files_x86Internet_Explorer = runtime_concatstring2(0, a10, a11, &unk_86A493, 40); // : \Program Files (x86)\Internet Explorer\
    v74 = a10;
    :Program_FilesWindowsApps = runtime_concatstring2(&v51, a10, a11, &unk_8671C2, 28); // \Program Files\WindowsApps\
    v76 = a10;
    :Program_FilesEmbedded_Lockdown_Manager = runtime_concatstring2(0, a10, a11, &unk_86AB8D, 42); // : \Program Files\Embedded Lockdown Manager\
    v78 = a10;
    v79 = runtime_concatstring2(&v50, a10, a11, ":\Program Files\VMware\\"", 23);
    v80 = a10;
    :ProgramDataMicrosoft = runtime_concatstring2(&v49, a10, a11, &unk_865D25, 24); // : \ProgramData\Microsoft\
    v82 = a10;
    n15 = 15;
    v83 = runtime_concatstring2(&v48, a10, a11, ":\$Recycle.Bin\\"", 15);
    v84 = a10;
    v16 = v58;
    filePath 1 = filePath;
}

```

Below we summarize the list of paths excluded by the ransomware during encryption.

## BlackListed Filepaths:

- \WindowsPowerShell\Modules\
- \Windows
- \ProgramData
- \Program Files\Reference Assemblies\
- \Program Files\Common Files\
- \Program Files\Internet Explorer
- \Program Files (x86)\Reference Assemblies\
- \Program Files (x86)\Common Files\
- \Program Files (x86)\Internet Explorer\
- \Program Files\WindowsApps\
- \Program Files\Embedded Lockdown Manager\
- \Program Files\VMware\
- \ProgramData\Microsoft\

- \\\$Recycle.Bin\

The next check determines whether the ransomware extension is at the end of the file, and whether the extension is *.ini* or *.sys*. If it is either of these, the file is skipped. After this, it checks whether the extension is *.bak* or *.log*, and if it is, it proceeds with deleting the file and returning. If a file doesn't pass any of these checks, the encryption function is called on it. Since these checks do not check whether the file is the ransomware executable itself, the ransomware ends up encrypting its own executable. However, it fails to remove the plaintext executable after the encryption process using *os.remove*.

```
if ( result )
{
    v37 = j_2 == 4;
    if ( j_2 == 4 )
    {
        if ( *fileExtension == 'ini.' || *fileExtension == 'sys.' )// skip .ini and .sys files
            return result;
        v37 = 1;
    }
    if ( v37 && (*fileExtension == 'bak.' || j_2 == 4 && *fileExtension == 'gol.') )// remove files with .bak or .log in their name
    {
        cap_7 = cap;
        result = os_Remove(filePath, cap, j, cap_3, n15, v28, v29, fileExtension, i_2);
        if ( result )
            return B0Client_Error_CheckError(result, cap_7, v39, cap_3, n15, v40, v41, v42, v43);
    }
    else
    {
        cap_7 = cap;
        LODWORD(cap_3) = edi0;
        n15 = n15_2;
        result = B0Client_Ecryptor_Encrypt(filePath, cap, ecx0, edi0, n15_2, r8_0, r9d0, fileExtension, i_2);
        if ( result )
            return B0Client_Error_CheckError(result, cap_7, v39, cap_3, n15, v40, v41, v42, v43);
    }
}
```

The file encryption function first prints the filepath. After which, it attempts to open the file. If that fails, the ransomware just skips the file instead of trying to kill the process that's using the file. The ransomware also checks that the file is not a directory after opening it.

```
fmt_Fprintln(&go_itab_os_File_io_Writer, os_Stdout, &v34);
v30 = os_OpenFile(filePath, cap);
err.cap = cap;
v28 = v11;
if ( B0Client_Ecryptor_IsFile(filePath, cap) )
{
    if ( err.cap )
        // check if err is not nil
    {
        LODWORD(error) = (*(err.cap + 24))(v28, cap, err.cap, 0LL, 1LL);
        if ( strings_Index(error, cap, &sunk_868FD2, 34, 1) >= 0 )// The process cannot access the file
        {
            v35 = v9;
            v15 = runtime_convTstring(filePath, cap);
            *v35 = &RTYPE_string;
            *(&v35 + 1) = v15;
            *v32 = fmt_Errorf(&dword_86739E, 28, &v35);// file is locked or in use: %s
            *(&v32 + 1) = 28LL;
        }
        else
        {
            v36 = v9;
            v37 = v9;
            v13 = runtime_convTstring(filePath, cap);
            *v36 = &RTYPE_string;
            *(&v36 + 1) = v13;
            *v37 = *(err.cap + 8);
            *(&v37 + 1) = v28;
            *v32 = fmt_Errorf(&dword_86905A, 34, &v36);
            *(&v32 + 1) = 34LL;
        }
    }
}
```

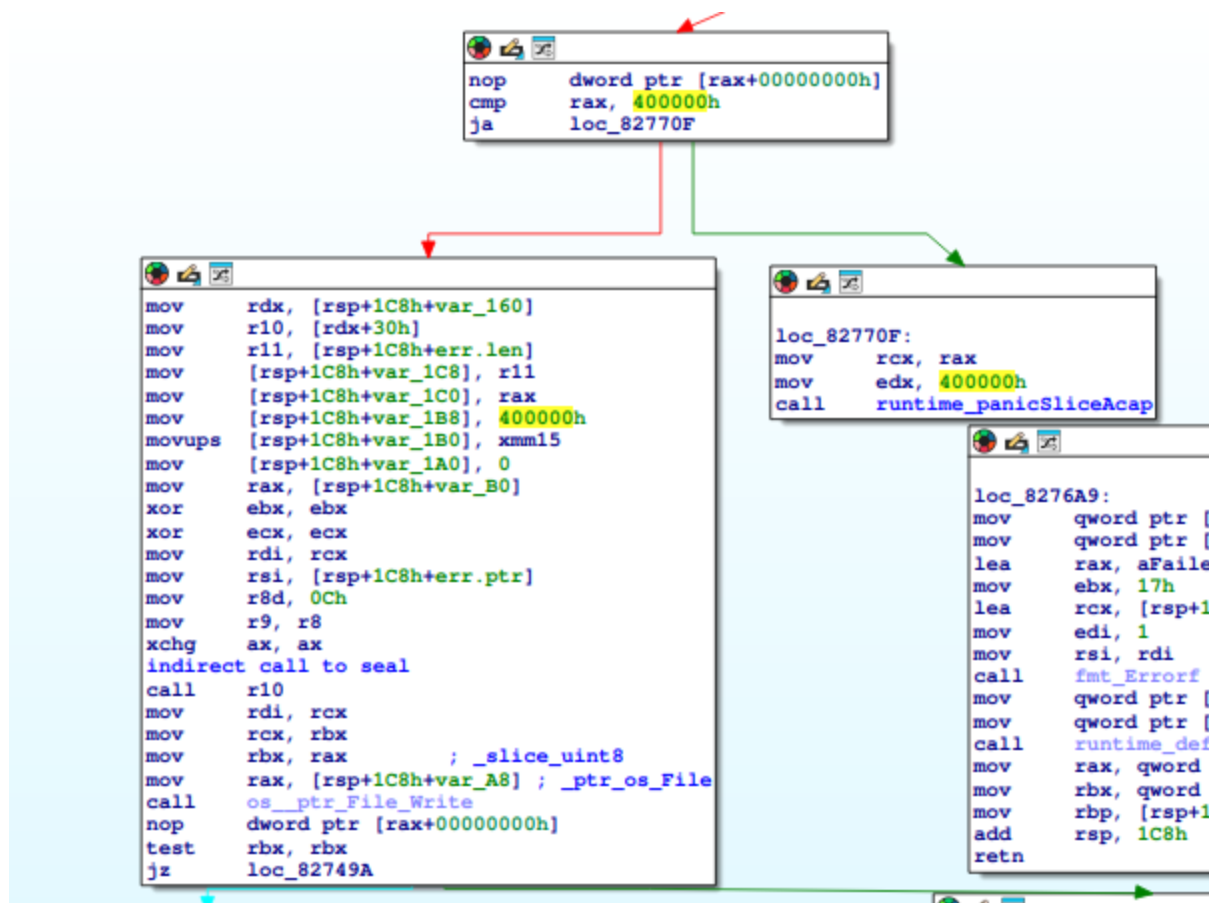
The ransomware initiates the encryption process by generating an AES cipher object and initializing it with the previously generated key. It uses GCM mode with a 16-byte tag size and a 12-byte nonce size. The nonce is randomly generated using the *crypto/rand* package.

```

tab = crypto_aes_NewCipher(n15, data, r9d0);
if ( v18 )
{
    v31 = v9;
    v31 = *(v18 + 8);
    *v28 = fmt_Errorf(&dword_8681D8, 31, &v31);
    *(&v28 + 1) = 31LL;
    runtime_deferreturn(v28);
    return v28;
}
else
{
    crypto_cipher_newGCMWithNonceAndTagSize(tab, data, 12LL, 16LL);
    if ( v19 )
    {
        v31 = v9;
        *v31 = *(v19 + 8);
        *(&v31 + 1) = 16LL;
        *v28 = fmt_Errorf(&dword_8681F7, 31, &v31);
        *(&v28 + 1) = 31LL;
    }
    else
    {
        data_1 = data;
        err.ptr = runtime_makeslice(&RTYPE_uint8, 12);
        crypto_rand_Read(err.ptr, 12);
        *v31 = MEMORY[0x14];
        *(&v31 + 1) = v20;
        *v28 = fmt_Errorf(&dword_86734A, 28, &v31);
        *(&v28 + 1) = 28LL;
    }
    runtime_deferreturn(v28);
}

```

After this, the ransomware creates a new file named by appending *.id-* to the original filename. The initially generated key is stored both in the configuration file and within each encrypted file's name. The ransomware then writes a 16-byte marker, consisting of null bytes, followed by the random nonce, to the new file. After this, it reads data from the original file in chunks of 0x400000 bytes, encrypts each chunk using the *Seal()* function, which handles both encryption and authentication. The resulting ciphertext, along with the 16-byte authentication tag, are both written to the new file. Once all chunks have been encrypted and written, the original file is deleted using *os.remove*.



## Encrypted file structure

The structure of the encrypted file is quite straightforward. At the beginning of the encryption process, a 16-byte marker (highlighted in purple) consisting of NULL bytes is added to the start of the file to indicate that it is encrypted. Following this marker, there is a 12-byte nonce (highlighted in Red in the screenshot below) value that remains consistent throughout the entire file. This nonce is stored within the file for use during decryption later. The rest of the data is just the file's encrypted data (highlighted in light green in the screenshot below).



Offset(h)	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F	Decoded text
00000000	B0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	.....
00000010	B1	BD	41	84	2E	6A	A9	29	0C	C1	D1	A7	CA	9F	25	C2	±A,,.j@).ÁÑŞÊÝ%Â
00000020	1E	1E	E6	49	02	EA	DB	90	73	80	38	E5	33	CB	4E	B9	..æI.êÛ.sε8â3ËN²
00000030	90	08	C8	FD	34	66	37	78	C2	B1	4B	90	09	31	5F	09	..ËÝ4f7xÂ+K..l_.
00000040	EF	0D	E5	87	BD	5D	FB	43	2A	34	5D	60	FF	3C	9A	11	i.â+½]ûC*4]`ÿ<š.
00000050	D0	16	DF	AD	CB	15	0D	BA	91	03	DC	A8	91	59	95	19	Ð.ß.Ë...°`·Û`·Y·.
00000060	9E	96	1D	09	01	05	E8	CA	E5	08	82	9C	9A	29	C1	19	ž-...èÊâ.,œš)Á.
00000070	CC	DB	59	9B	D8	BA	27	9A	A3	86	11	0A	51	30	12	87	ÏÛY>ø°'š£+..Q0.+
00000080	23	50	0E	08	01	AF	B2	BB	D6	2E	F8	94	81	3A	F1	04	#P...¯°»Ö.ø".:ñ.
00000090	A4	64	EC	43	EA	3B	AC	EE	E0	22	7F	78	C2	E5	27	50	µdiCê;-îà".xÂâ'P
000000A0	08	65	DF	7C	4B	D4	AA	EB	30	70	10	CC	5B	27	01	F7	.eß KÔ²ëOp.Î['.÷
000000B0	0C	9C	AC	59	35	9A	DF	F2	88	98	0B	EC	52	7C	2F	99	.œ-Y5šßò`~.îR /²µ
000000C0	C5	05	B7	51	38	CA	28	9D	68	23	49	EB	D4	AD	0E	5B	Å..Q8Ê(.h#IëÖ..[
000000D0	06	40	FD	0E	53	5D	DD	48	B3	0C	8E	2C	C8	42	9E	04	.@ý.S]YH³.Ž.ËBž.
000000E0	01	CF	ED	4A	84	06	39	25	28	36	19	93	99	D4	ED	3A	.İiJ,,.9% (6.´²µÖi:
000000F0	F4	03	68	74	16	45	C2	0F	3A	41	CA	C2	19	EF	66	E1	ô.ht.EÂ.:AÊÂ.îfá
00000100	60	F8	DA	4D	67	2F	6C	0A	1F	48	58	FC	B4	71	9D	13	`øÛMg/l..HXü'q..
00000110	70	D4	2D	70	CF	EE	44	AD	18	DB	8C	98	17	26	98	B3	pÔ-pİiD..ÛÊ~.ε~²
00000120	A5	A1	A7	3A	C7	BB	6C	BA	F2	5C	87	C4	F8	24	F3	2F	≡j\$;Ç»l°ò\+Ãø\$ó/
00000130	92	0E	85	07	41	4C	A4	43	A3	6D	D5	12	A1	41	CD	CD	'...ALµC£mÖ..jAÍÍ
00000140	95	B7	35	77	3F	9C	23	91	63	38	6F	77	33	AE	FE	63	*·5w?œ#`c8ow3@)pc
00000150	FE	75	61	5F	5E	81	AD	2C	E7	DB	1C	3F	CE	0A	F1	4E	pua_^..,çÛ.²Î.ñN
00000160	0D	86	33	44	09	D7	85	21	03	D4	1D	AC	3D	A7	B6	9C	.+3D.x...!..Ô.¬=\$qœ
00000170	31	0E	3A	93	8E	C8	FF	26	F6	B3	24	12	59	DB	DC	FF	l.:`ŽÈÿ&ö³\$.YÛÛÿ
00000180	F4	E9	B7	FD	22	F8	BB	0C	BF	80	0B	D9	C0	C8	75	8F	ôé·ý"ø».¿€..ÛÀEu.
00000190	8E	3C	D6	9C	26	1F	CC	39	3F	5B	70	1F	9E	2B	60	BB	Ž<Öœ&.î9?[p.ž+»
000001A0	93	67	57	9D	2D	E2	FB	B7	00	07	30	2A	D6	AB	68	82	`gW.-âû...0*Ö«h,
000001B0	54	B0	80	D2	0C	F0	53	63	AE	3B	EC	B7	CE	5D	2F	2A	T°ëÒ.ðSc@;î·Î]/*
000001C0	E4	63	EE	39	39	69	11	D1	58	09	50	E8	1F	88	79	70	äcî99i.ÑX.Pè.^ÿp
000001D0	3B	95	DA	01	1D	DA	FF	99	8E	A0	E9	A4	EE	E8	69	B9	;·Û..Ûÿ²² éµîèi²
000001E0	5F	8C	2F	FA	88	DC	1F	33	D4	03	BD	92	46	21	77	44	_E/ú^Û.3Ö.½'F!wD
000001F0	BA	DA	C9	12	DF	5D	49	5F	06	5E	E3	95	F7	4E	C4	43	°ÚÉ.ß]I_.^ä·-NÄC
00000200	61	8B	DD	6D	86	D2	12	99	71	E9	7F	07	98	4A	8D	63	a<ÿm+Ö.²µqé...~J.c

Encrypted file structure highlighting marker, nonce, and encrypted bytes.

## Weakness

There are two major weaknesses in the encryption. The first, and most obvious, is that the AES-GCM key can be easily derived from the initial key (randomly generated nine-character string), which is unique to each victim and generated at run-time, then saved to disk later.

The second weakness is less apparent. Although the ransomware uses AES-GCM to encrypt blocks (which combines AES-CTR for encryption and GMAC for authentication), it doesn't generate a new nonce for each block. As a result, the same keystream is reused for each block within a file. If we know the plaintext for one block, we can recover the keystream by xoring it with the corresponding ciphertext and then use that keystream to decrypt all subsequent blocks. However, exploiting this vulnerability is not straightforward, since it requires a significant amount of known plaintext (0x400000) bytes.

## Decryptor

We have created a [b0 ransomware decryptor](#) for those affected by this ransomware. Errors are logged in .log.txt. Below is the usage guide.

Usage of decryptor:

- dirpath string: Path to the directory to recursively decrypt files
- key string: Key to use for decryption, in the following filename test.pdf.id-QVKGBICKS.B0-3e72d, QVKGBICKS is the key.
- path string: Path to the encrypted file

## Indicators Of Compromise

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### Known hashes

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cea74b854b6cd161884e5ced5c8b7ba44c1064ff22703667cc913d9a67f1767b (SHA-256)

### File Paths

---

Ransomware executable	B0.exe
Ransomware configuration written to disk at run-time	C:\Windows\System32\ telemetry.ASM-WindowsDefault.json
Ransomware icon	C:\Windows\Icon.png
Batch script	C:\Windows\System32\Del.cmd
Ransom note in text and HTML versions	C:\How_To_Recovery_Files.txt C:\How_To_Recovery_Files.html
Registry entry created to clear Windows Event logs on user login	HKCU:\Software\Microsoft\Windows\CurrentVersion\Run\DeleteLogWithStartUp

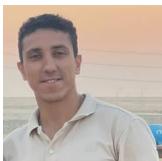
## Conclusion

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This investigation is a powerful reminder that even in an ecosystem dominated by rapidly evolving ransomware threats, there's still room for skilled analysis, persistence, and creative problem-solving to make a real impact. While not every case ends in a decryptor, this one proves that deep malware analysis and DFIR expertise can tip the scales.

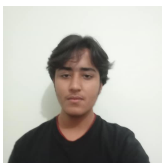
## Authors

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[Mohamed Talaat](#)

Mohamed Talaat is a self-taught security researcher specializing in malware analysis, threat intelligence, and tracking APT groups and crimeware. With a background in engineering and over three years of independent research, his focus is on ransomware recovery, evaluating recovery paths, analyzing encryption schemes, and occasionally developing decryptors. Passionate about reverse engineering and low-level analysis, he also writes YARA rules to help detect and hunt emerging threats and campaigns.



[Hassan Faraz](#)

Specializing in malware reverse engineering and data recovery, Hassan Faraz develops decryption solutions for vulnerable ransomware and creates custom tools to extract data from various file formats such as VHD, VMDK, MDF, backup formats, etc. He has a proven track record of recovering critical file formats with minimal data loss, helping organizations mitigate the impact of cyberattacks.